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L. AMMERAAL
AN INTERPRETER FOR SIMPLE ALGOL 68 PROGRAMS

2e boerhaavestraat 49 amsterdam

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Abstract

An interpretative implementation for a sublanguage of ALGOL 68 is informally described in this report. Since the implementation has been written in ALGOL 60 and there is no translation from ALGOL 68 source text to an object program, a high degree of machine independence has been achieved. The report includes both the text of the interpreter and some typical ALGOL 68 programs, along with the output produced by the interpreter.

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1. Introduction

An implementer of a high-level programming language will have to pay attention to a number of general implementation aspects which may include:

- . efficiency with respect to speed and to use of (core) memory,
- . completeness,
- . error diagnostics or even error correction,
- . machine independence,
- . flexibility with respect to future language revisions,
- . readability.

As many implementers are concerned about the first three of these points in particular, and because of the ambitious nature of the language, the implementation of ALGOL 68 [1] is generally considered to be a very big software project. Although there is no doubt that these three points are very important, it may also be worthwhile to gain some experience with a simpler type of ALGOL 68 implementation, with less emphasis on machine efficiency, no error correction facilities and perhaps some language restrictions, and to see whether other aspects, such as machine independence, flexibility and readability, may be improved as a consequence of these concessions.

A tolerant attitude towards the requirements of machine efficiency leads the implementer to the idea of using a high-level language like ALGOL 60 which has, of course, well-known advantages over the traditional use of assembler languages. Choosing ALGOL 60 for this purpose, at least two kinds of highly machine independent implementations can be thought of.

The first is an ALGOL 68 compiler, written in ALGOL 60, and producing an object program in ALGOL 60.

The second is to use ALGOL 60 for the implementation, but to produce no object program at all, i.e. to let the implementation consist of an ALGOL 68 interpreter, written in ALGOL 60. Since the latter approach involves no object program, it seems to be conceptually simpler than the former. As to machine efficiency, it is clear that we must not expect too much of such an interpreter, at least not if we are interested in the execution speed of

production jobs.

From a more theoretical point of view, this interpretative type of implementation might have some value in the following sense. Semantic actions are described in the Report [1] in a very concise and almost algorithmic way, but not in a form that can be presented to a machine in order to perform these actions. In an interpreter, however, the semantic actions are described in such a form. This description is more direct in an interpreter than in a compiler where the actions are only prepared.

In view of these considerations an interpreter was written in ALGOL 60. It will be discussed in the next chapters of this report.

2. Elaboration as a side-effect of syntactic analysis

The top-down method for syntactic analysis has well-known attractive properties. The method is presented in a tutorial manner by Knuth [4]. It has often been illustrated by parsing simple pieces of ALGOL 60 programs. Top-down parsing is less often described using ALGOL 60 as a tool, however. Yet this can be done very well by means of a set of Boolean procedures. A method for generating such a set automatically, was suggested by Koster [2].

Although this idea is very attractive, the Boolean procedures mentioned in the sequel of this report were written by hand. They are based on the assumption that the program to be analysed is stored in an integer array X , and that at any moment a global integer p , the so-called program pointer, has such a value that $X[p]$ is the point from which the analysis has to proceed at that moment. Initially, $p=1$. Now consider the production rule for some notion N of a given context-free grammar. Then there is a Boolean procedure N with the task to determine whether or not the array X contains a terminal production of N starting at $X[p]$ for the current value of p . If there is not such a terminal production the value false is delivered. If, however, there is, the value true is delivered, and the procedure has to perform at least one additional task, namely to set the program pointer p to the new value p' where $X[p']$ is the point in the program immediately after the found terminal production.

The two tasks mentioned, of which the second is a typical side effect, suffice for the syntactic analysis; the following hypothetic example may show how this can be done:

Example 1

A : B; p,q; C,D .
B : r,A; u .
C : t,u .
D : v; w,D,z .

In contrast to the ALGOL 68 report, here capital and small letters mean notions and (terminal) symbols, respectively.

To analyse in ALGOL 60 whether a given integer array X contains a terminal production of A, and assuming that p has a correct initial value, say 1,

we need just call the Boolean procedure A, given below, which will yield the desired answer true or false:

Boolean procedure INP (s); integer s;

if X[p] =s then begin INP:= true; p:=p+1 end else INP:=false;

Boolean procedure AND (U,V); Boolean U,V;

begin integer pold; pold:=p;

if if U then V else false then AND:=true else

begin AND:=false;

p:=pold

end

end;

Boolean procedure A;

A:= if B then true else

if AND (INP(p), INP(q)) then true else

AND(C,D);

Boolean procedure B;

B:= if AND(INP(r),A) then true else INP(u);

Boolean procedure C;

C:= AND(INP(t), INP(u));

Boolean procedure D;

D:= if INP(v) then true else

AND(INP(w), AND(D,INP(z)));

It may be a useful exercise to verify by means of these procedures that

rrtuwwvvzzz

is a terminal production of the notion A.

Of course, the Boolean expressions

if E then F else false

and

if G then true else H,

must not be replaced by

$E \wedge F$

and

$G \vee H$,

respectively, since in the former expressions, it depends on E and G whether F and H , respectively, will be evaluated, but in the latter expressions both operands will be evaluated in any case.

Although this undesirable behaviour of the \wedge - and \vee - operators has often been mentioned by others, it may be observed that top-down parsing is another example where it would be particularly convenient if \wedge and \vee had the property that the second operand is only evaluated if the evaluation of the first operand does not suffice to determine the result. If they had this property, both expressions

if if A then B else false then C else false

and

if if A then (if B then C else false) else false,

e.g., could be replaced by

$A \wedge B \wedge C$.

The transformation of grammar rules into Boolean procedures, as sketched above, is not allowed for each context-free grammar. If it is for some CF-grammar, care has to be taken in choosing the order of the alternatives of a given notion.

For these reasons the following considerations may be useful.

Suppose that α is a string of n terminal symbols, β is the string composed of the first m symbols of α ($0 \leq m < n$) in the same order, and both α and β are produced by a notion N . Now if a terminal production of N has to be recognized, starting at $X[p]$, and it happens that

$\alpha = X[p] \dots X[p+n-1]$ and

$\beta = X[p] \dots X[p+m-1]$,

we want to make a definitive choice.

An example of such a situation is the choice between the real denotation 12.3 and the integral denotation 12 , if a denotation has to be recognized and the four symbols 12.3 are available in $X[p] \dots X[p+3]$. In this example the second alternative must not be taken, because otherwise the two symbols $.3$ in $X[p+1]$ and $X[p+2]$ cannot be parsed afterwards. This knowledge cannot be derived from the rule for the notion "denotation", but other rules have to be considered as well.

Similarly, in the general case above we will reject the choice of the shorter string β if the longer string α is also available as a candidate. By this strategy, a grammar with the rules

$N : P, Q .$
 $P : u, v ; u .$
 $Q : v, w ,$

cannot be used in this form to recognize the string uvw as a terminal production of N , so not all context free grammars are suited for our method.

For the grammar rules

$N : P, Q .$
 $P : u, v ; u .$
 $Q : w .$

uvw can be recognized, however. The procedure body for P is

$P := \underline{if} \text{ AND } (INP(u), INP(v)) \underline{then} \underline{true} \underline{else} INP(u)$

For reasons of efficiency, however, the rule for P is preferably replaced by the rules

$P : u, R .$
 $R : v ; .$

Now the same considerations as above apply to R , with $\alpha = v$, $n = 1$, $\beta = \epsilon$ (the null string), $m = 0$. The procedure body for R might be

$R := \underline{if} INP(v) \underline{then} \underline{true} \underline{else} \underline{true} .$

It is well-known that for top-down parsing, left-recursive rules have to be excluded. This phenomenon and many other interesting topics concerning syntactic analysis are extensively discussed by Knuth [4].

Syntactic analysis is only part of an implementation. If the implementation had the task to produce an object program, this, or some preparation for it, would have been accomplished by another side-effect of the Boolean procedures. In an interpreter, however, the semantic actions themselves may be executed; in other words, the elaboration of a notion may take place as a side-effect of the Boolean procedure corresponding to that notion. If the notion was e.g. formula, and the found terminal production was $3+4$ it may be our intention that the sum 7 is immediately yielded. On the other hand, it has sometimes to be decided whether there can be found a terminal production of a given notion, starting at $X[p]$, and if so, p should be advanced, but the elaboration side-effect must not take place. Still another possibility exists, namely the requirement that only defining occurrences of either labels or modes, if present, are dealt with. This will be discussed for labels in chapter 8.

So there are three cases with respect to the desired elaboration side-effect. Information about this is supplied to the procedure by means of a formal parameter q of which the corresponding actual parameter may have one out of three values. These three actual parameter values are written here as the letters f , g and h , which might be associated with the speed of the scanning process:

- $q = f$ (fly) : no elaboration required
- $q = h$ (hasten) : only defining occurrences of labels and modes
are to be dealt with
- $q = g$ (go) : all other elaboration actions required.

3. The internal representation of the ALGOL 68 text

The input for the interpreter is an ALGOL 68 program whose smallest compounding parts, at this level, are characters.

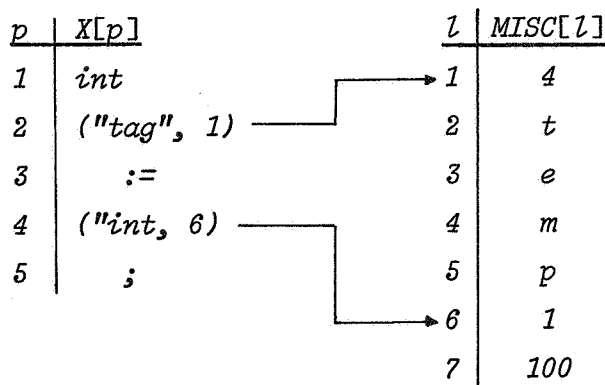
Thus the piece of ALGOL 68 text

```
'INT' TEMP:=100;
```

is a sequence of 16 characters.

The first step after the program has been read character by character, is to form symbols from small sequences of characters, e.g. the five character sequence 'INT' is replaced by the integral symbol int and the two character sequence := is replaced by the becomes symbol :=. Blank spaces that are not inside string denotations are omitted, of course. During this lexical scan the integer array X, mentioned before, is filled. For a tag, such as TEMP, only one element of X is used, containing both a reference to a point in an auxiliary integer array MISC, where the letters of the tag are stored if they were not already, and the mark "tag", for recognition purposes later on. Likewise an integral denotation like 100, is stored into the auxiliary array MISC and a reference to this is put into array X, together with the mark "int".

Now the example above may be represented internally as follows



The array MISC is used for miscellaneous purposes. It contains variable length items, which are all mutually different. Immediately preceding each

item, its length is written. Now if e.g. the tag *TEMP* is encountered later on, perhaps with a completely different meaning like a label or a selector, a search for *TEMP* is performed in *MISC* and when it is found at the point l' , this value l' is written into array *X*, together with the mark "*tag*". This mark is simply a unique small integer, also called a mnemonic constant. A number pair like ("*tag*", l) is stored as the integer "*tag*" * $C + l$. Here C is some appropriate large integer, e.g. 100 000. Besides the mentioned marks "*tag*" and "*int*", there are also the marks "*real*", "*ind*", "*char*" and "*string*" for real denotations, indications, character denotations and string denotations, respectively. They are all coupled with a reference l which, with the exception of "*real*", points to the integer array *MISC*, where the characters composing the given denotation or indication are stored. To store real denotations a real array *RNR*, instead of the integer array *MISC*, is used. Here each item consists of only one element, so the item need not be preceded by the item length. A pleasant property of the sketched method of internal representation is the lack of limitations with respect to the lengths of tags and indications other than those which are implied by the finite size of array *MISC*.

4. The internal representation of values

In ALGOL 68 the elaboration of a notion like a serial clause, a unit, a tertiary, a secondary or a primary, delivers a value. Such a value may disappear shortly after it is created, like the value of $(a+b)$ in the elaboration of the formula

$$(a+b) * c .$$

On the other hand, there may be a mode identifier (or an operator) that possesses the given value, which relationship should hold for some time. Consider, as an example, the identity declaration

$$\underline{int} \ i = 5$$

As we are now using the output produced by the lexical scan, we have the number pairs $(\text{"tag"}, l^i)$ for i , and $(\text{"int"}, l^5)$ for 5 in array X , with

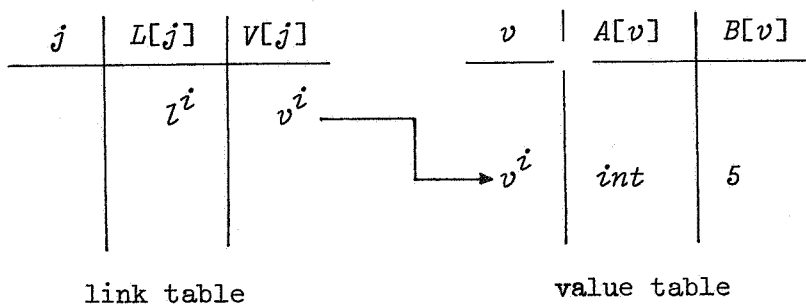
$$\begin{aligned} MISC [l^i] &= 1 \\ MISC [l^i+1] &= \text{the integral representation for the character } i \\ MISC [l^5] &= 1 \\ MISC [l^5+1] &= 5 . \end{aligned}$$

Now as a result of the elaboration of the given identity declaration, l^i is written in the first free element $L[j]$ of the integer array L . This array L will be regarded as the first column of a two column table, called link table, the second column being integer array V . Since the identifier i is to possess the value 5, it looks attractive to set $V[j]$ to 5 by which the relationship to "possess" between an external and an internal object would simply be visualised by the contents of a row of our link table. Unfortunately this cannot be done this way in general because not every value in ALGOL 68 will fit in a single integer array element $V[j]$. Moreover, not only the value itself, but also information about the mode of this value has to be stored. A third reason for introducing something else than the link table to store values is the superfluity of column L for

"anonymous" values like $a+b$ in the formula $(a+b)*c$. Here we meet something characteristic for ALGOL 68. In most other languages it is not necessary to bother about the way in which such an intermediate result $a+b$ is stored since this value appears only very temporarily in a single machine word or even in a register, whereas in ALGOL 68 this value $a+b$ may have quite a complicated structure. It would probably be efficient with respect to storage space, to store values of some simple mode in a table, different from a table used for values of a more complicated mode, or to use some method for mapping a given value onto a one-dimensional array. For reasons of simplicity, however, but at the cost of a less efficient use of storage, a very easy and general method to represent values is chosen in our interpreter. Values, simple or complicated, anonymous or possessed by either an identifier or an operator, are all stored in a so called value table. This consists of the five columns A , B , C , D , E . Each of the columns A , B , C and D is an integer array. E is a real array. Now as a result of the elaboration of

int $i = 5$

we obtain the following picture



In column V of the link table, a reference to the v^i -th row of the value table is stored, where v^i is obtained using a free-list technique. In the value table, columns A and B contain the mark " int " and the number 5, respectively. Columns C and E have no function for this simple case; column D has, as will appear in the sequel.

For a boolean value the mark "bool" will appear in column A and the integer 1 or 0 in column B for true or false, respectively.

Likewise for a character, the mark "char" is put into column A, and an integer representation of a character into column B. A real number will be put into column E, in which case column B is not used and the mark "real" is in column A.

Now suppose that after the identity declaration

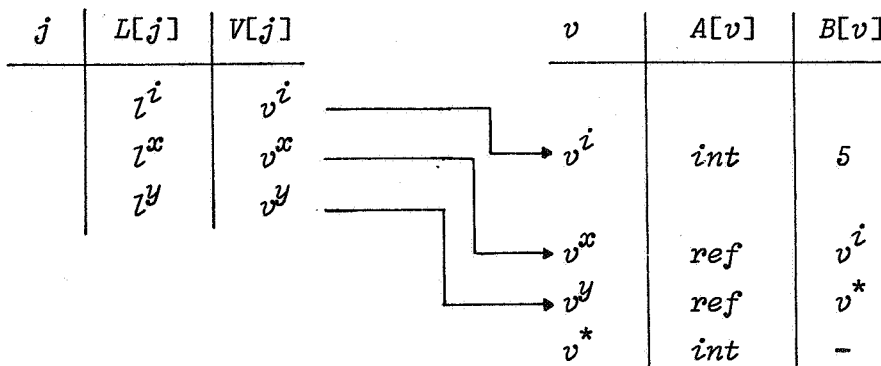
int $i = 5$

with the resulting picture above, the variable declaration

int $x := i$, int y

follows.

Then the given picture will be supplemented as follows.



A line in the link table appears for x and for y . The references v^x and v^y are put into column V of these lines, pointing to new lines in the value table with the mark "ref" in column A, since x and y have modes beginning with "reference to". Now x has to refer to the integral value i , so v^i is put into $B[v^x]$. Likewise, y has the mode "reference to integral" but no integral object for y to refer to is given by the declaration. Yet the given information about the mode of y has to be stored somehow. This is done here by creating a new integral value on a line v^* and by assigning v^* to $B[v^y]$. Note that $v^* \neq v^y + 1$ may hold. It will be clear that a mode beginning

with more than one "reference to" will easily be represented in the value table by additional lines with "ref" in column A and values v in column B.

The representation of a multiple value consists of three parts in the value table. The first part is the so-called "initial line" with the mark "row", a reference v^d to the descriptor, and the number (say n) of subscripts in columns A, B and C, respectively. The second part is the descriptor part, starting with line v^d , with

$A[v^d] = 1$ or 0 , for the state (strict or flex),

$B[v^d] = v^{fe}$, a reference to the first element,

$C[v^d] = c$, the offset, being 1 in most cases.

This line is immediately followed by n lines, one for each subscript. For subscript i ($i=1, \dots, n$) we have

$A[v^d+i] = l_i$, the lower bound of the i -th subscript,

$B[v^d+i] = u_i$, the upper bound of the i -th subscript,

$C[v^d+i] = d_i$, the stride of the i -th subscript.

The third part is filled as follows. If $u_i \geq l_i$ for all i ($1 \leq i \leq n$), the number of elements is $m = \prod_{i=1}^n (u_i - l_i + 1)$, and the elements, or at least their initial lines, are stored at the lines $v^{fe}, \dots, v^{fe} + m - 1$, in the usual order; otherwise one fictitious element is stored in the same way, using line v^{fe} .

Thus the example

```
[5 : 6 , 21 : 23] int a ;
```

will result in the following scheme:

j	$L[j]$	$V[j]$	v	$A[v]$	$B[v]$	$C[v]$
	:		v^a	<i>ref</i>	v'	
	l^a	v^a	v'	<i>row</i>	v^d	2
	:	:	v^d	1	v^{fe}	1
			v^{d+1}	5	6	3
			v^{d+2}	21	23	1
			v^{fe}	<i>int</i>	-	
			v^{fe+1}	<i>int</i>	-	
			v^{fe+2}	<i>int</i>	-	
			v^{fe+3}	<i>int</i>	-	
			v^{fe+4}	<i>int</i>	-	
			v^{fe+5}	<i>int</i>	-	

In the chosen method of representation, it is essential that, if an initial line of some value is complemented by other lines, these other lines need not immediately follow the initial line, but may appear at any point in the value table.

Another application of this principle is the representation of a structured value. Like a multiple value, a structured value is represented in the value table by several parts, the first of which is its initial line.

In the first three columns, this line contains the mark "*struct*", a reference v^d to a "descriptor", and the number (n) of fields.

The descriptor, being the second part, consists of the lines

$v^d, v^{d+1}, \dots, v^{d+n-1}$, with

$B[v^d + i - 1] = v^i$, a reference to (the initial line of) the i -th field

$C[v^d + i - 1] = l^i$, the unique representation, see chapter 3, of the tag used for the i -th selector

($i=1, \dots, n$).

Finally, each field of the structured value is represented in the general way in which values are stored, i.e. the i -th field or its initial line is v^i . In contrast to a multiple value, the initial lines v^i of the elements of a structured value need not be consecutive.

The next example will show that the value table is not only used to store values. It will not only illustrate the representation of a structured value, but show the treatment of a mode declaration and a collateral clause as well.

Consider the declarations

```
mode vec = struct (int x,y,z);  
vec a = (1,2,3)
```

For the present version of the interpreter, these two declarations must not be interchanged.

The ALGOL 68 report states that the elaboration of a mode declaration involves no action. The interpreter, indeed, performs no substantial action in elaborating the mode declaration; it only makes a note where the mode declaration occurs in the program, to be able to use it afterwards. For this note, both the link table and the value table are used. In chapter 3, it was mentioned that an indication is represented as a number pair ("ind", l) in a single element of array X. So both occurrences of the mode indication vec appear as the pair ("ind", l^{vec}). Now, regardless of the text following the equals symbol, the effect of

```
mode vec = ... .
```

will be both the addition of a line j , with $L[j] = l^{vec}$, $V[j] = v^{vec}$, to the link table and the addition of a line v^{vec} to the value table, where $A[v^{vec}]$ is set to the mark "newmode", and $C[v^{vec}]$ is set to the value p' , pointing to the beginning of the actual mode declarer in array X.

On encountering the second occurrence of vec, v^{vec} is looked up in the link table with the given l^{vec} , in a way to be described in chapter 6. Using v^{vec} , the value of $p' = C[v^{vec}]$ is found and now the current value p of the program pointer is temporarily replaced by p' , in order to use the normal syntactic analysis and elaboration procedures, resulting in a new, fictitious value of the mode struct (int x,y,z). This value is represented in the value table as described above. The next elementary step is to

elaborate the given collateral clause. Since a collateral clause may also be used in connection with a multiple value, the interpreter creates first an intermediate value, using in its initial line the mark "collat", a reference v' and the number (say n) of elements; v' points to the consecutive lines v' , $v'+1, \dots, v'+n-1$ where the elements or at least their initial lines are stored. Afterwards, this intermediate value is transformed into either a structured value or a multiple value and its space in the value table is released. The ultimate picture resulting from the two given declarations is

j	L[j]	V[j]		v	A[v]	B[v]	C[v]
	z^{vec}	v^{vec}	→	v^{vec}	new mode	-	p'
	z^a	v^a	→	v^a	struct	v^d	3
				v^d		v^x	z^x
				v^{d+1}		v^y	z^y
				v^{d+2}		v^z	z^z
				v^x	int	1	
				v^y	int	2	
				v^z	int	3	

Since a routine is a value in ALGOL 68 which may be possessed by an identifier or an operator, it belongs in the value table like the values met before. It is more appropriate, however, to store only part of it in the value table, together with a reference p , pointing to the beginning p of the routine text $X[p], X[p+1], \dots$.

Consider the example

proc $f = (\underline{int} \ i, \underline{real} \ r) \underline{bool} : 5 * i \uparrow 2 > r ;$

op $s = (\underline{int} \ i) \underline{int} : i * i + i ;$

The resulting picture will be given first and then explained.

j	$L[j]$	$V[j]$	v	$A[v]$	$B[v]$	$C[v]$
	z^f	v^f	v^b	$bool$		
	z^s	v^s	v^i	int		
			v^r	$real$		
			v^f	$proc$	$v^{f'}$	
			$v^{f'+1}$		2	p^f
			$v^{f'+2}$		v^b	
			$v^{f'+3}$		v^i	
			v^s	$proc$	$v^{s'}$	
			$v^{s'+1}$		-10	p^s
			$v^{s'+2}$		v^i	
			$v^{s'+3}$		v^i	

The first three lines of the value table are always filled as shown in the picture, so this is not a result of the given declarations.

A routine starts in the value table with an initial line v containing the mark "proc" and a reference v' to a second part of the table. If the routine is not possessed by an operator, $B[v']$ is the number (n) of formal parameters and $C[v']$ is a reference pointing to the beginning of the routine text in array X (if this routine text is contained in the particular program under consideration). The next $n+1$ lines contain the mode that the routine is to deliver and the modes of the formal parameters. Here a mode is given by a reference to a value having this mode. If the routine is possessed by an operator its internal representation is similar, except for $B[v']$. Here the priority number of the operator is stored negatively,

where -10 means that the operator is monadic. For routines whose routine texts are not contained in the particular program, but are available by virtue of the standard prelude, the value $C[v']$ is not a (positive) program pointer value, but a negative number corresponding uniquely to such a routine.

An identifier with some mode united from other modes is stored in the link table as a pair (l, v) like other identifiers. In the value table, the mark "union" is put into $A[v]$. If $v', v'', \dots, v^{(n)}$ are the initial lines of values whose modes are exactly the modes from which the given mode is united, and if v^0 is the value possessed by the identifier at a given moment, then this is represented internally by

$$\begin{aligned} B[v] &= v^* & , & & C[v] &= v^0 & , \\ B[v^*] &= v' & , & & C[v^*] &= n & , \\ B[v^*+1] &= v'' \\ & \vdots \\ B[v^*+n-1] &= v^{(n)} & . \end{aligned}$$

5. The elaboration of a formula

It seems attractive to treat the operators of the standard prelude, called "elementary operators" here, and the operators declared in the program in the same way as much as possible. Therefore, data concerning these elementary operators are stored in the first part of the value table beforehand. This is done by putting a unique negative number for each elementary operator in column *C*, as mentioned in the previous chapter. Likewise, part of the link table is filled beforehand, with a line (*l,v*) for each elementary operator.

For the elaboration of a formula a modified version of a well-known method is used, namely the method of transforming the formula into Reverse Polish. Instead of this form, the value of the formula is the result, however. This result is delivered in the value table, and a global variable *z* is made to point to its initial line. A piece of ALGOL 60 text will show this principle and some general constructions as well. In this text, the neater form

while ... do ...

is used instead of

for dummy:=dummy while ... do ...

Some general self-explanatory procedures are:

```
Boolean procedure LOOK (A); Boolean A;  
begin integer pold;  
    pold:=p; LOOK:=A; p:=pold  
end;
```

```
Boolean procedure LOOK2 (A,B); Boolean A,B;  
begin integer pold; pold:=p;  
    LOOK2:= if A then B else false;  
    p:=pold  
end;
```

Some procedures local to the procedure "formula" are only informally described here. The verb "to read" is used in this description in the sense

of the scanning process, with its possible side-effects, described in chapter 2. These procedures are:

Boolean procedure OPERATOR (*lwbprio, upbprio, optr, prio*);
value *lwbprio, upbprio*; integer *lwbprio, upbprio, optr, prio*;

[This procedure reads an operator, if possible, with a priority not lower than *lwbprio* and not higher than *upbprio*; if this attempt succeeds, this priority is assigned to *prio* and the found value table index of the operator is assigned to *optr*.]

procedure OPERAND;

[This procedure reads an "operand". Here "operand" means either a monadic formula or a secondary. The resulting value is delivered in the value table and pointed to by the global variable *z*, if parameter *q* of *FORMULA* equals *g*].

procedure DYADIC OPERATION(*v, w*); integer *v, w*;

[This procedure performs a dyadic operation. The value table index of the operator \underline{w} is found at the top *OPST*[*w*] of the operator stack. The value table indices of both operands are at the top of a value stack *VLST*, such that the effect of this procedure can be given by

$VLST[v-1] := z := VLST[v-1] \underline{w} VLST[v] ;$
 $v := v-1 ;$
 $w := w-1]$

If *OPST*[*w*] points to operator \underline{w} , then *PRIOST*[*w*] is the priority of \underline{w} ; *PRIOST* can be regarded as a second column of the operator stack. Now the main part of the procedure *FORMULA* is given by


```
Boolean procedure FORMULA (q); integer q;  
begin [declarations omitted]  
  
  if if LOOK (OPERATOR(10,10,dummy,dummy)) then true else  
  LOOK2 (SECONDARY(f), OPERATOR(0,9,dummy,dummy)) then  
  begin FORMULA:= true;  
    OPERAND;  
    if q=g then  
    begin VLST [1] := z ; v:=1 ;  
      PRIOST[0]:=0 ; w:=0 ;  
      while OPERATOR (0,9,optr,prio) do  
      begin while prio ≤ PRIOST [w] do  
        DYADIC OPERATION (v,w);  
        w:=w+1 ;  
        OPST [w]:=optr ; PRIOST [w]:=prio ;  
        OPERAND ;  
        v:=v+1 ; VLST[v]:= z  
      end;  
      while v > 1 do DYADIC OPERATION (v,w)  
    end else  
    while OPERATOR (0,9,dummy,dummy) do OPERAND  
  end   else FORMULA:= false  
end
```

Note that after a successful call of *FORMULA* with $q=g$, the global variable *z* has been made to point to the resulting value. This was done either by *DYADIC OPERATION* if the formula was dyadic, or by *OPERAND* if it was monadic. As to the task of the procedure *DYADIC OPERATION*, there are two possibilities. If the operator involved is elementary, the operation itself takes place within this procedure; if it is not, the operation is performed by using the operation declaration, contained in the ALGOL 68 program. This latter case has a strong analogy with the elaboration of a call, described in chapter 7.

6. The identification of identifiers and indications

Before the treatment of the elaboration of a call, something should be said about the identification method used. As mentioned in chapter 4, by the declaration

int i

the integer l^i , which corresponds uniquely to the tag i , is placed into the link table, together with a value table index v^i . This link table is filled consecutively starting at the top. At any moment it is known how far the link table has been filled, say up to line j , so for a new declaration, first j will be increased by one and then $L[j]$ and $V[j]$ will be used. Now if the identifier i has to be looked up, say for the elaboration of

$i := 1$,

a search for l^i is made in the link table, in the direction opposite to the way the link table was filled, i.e. starting at line j , then $j-1$, etc, until l^i is found. This search direction guarantees that the most recently elaborated declaration of the tag i will be taken, so if i is declared at the beginning of both a closed clause A and a closed clause B which is contained in A, an applied occurrence of i in the outermost reach of B will result in finding the declaration at the beginning of B.

Now on leaving B we have to get rid of this second declaration of i . This is accomplished by storing the current value of the link table pointer j' into a variable *jold* local to the Boolean procedure CLOSED CLAUSE on the entrance of this procedure, and by resetting j to the value *jold* on its exit. Moreover, all values in the value table referred to by the link table lines that are discarded at this step, are released as far as this is possible. More specifically, such a value is relaxed, which means that its reference counter, being the value in column D , is decreased by one. If, as a result of such a relaxation, the reference counter of a value obtains the value zero, then the value involved is actually released, which means that its initial line is added to a free list and that its auxiliary lines are also added to the free list if these lines are not at the same time initial lines of values; if they are, however, they are relaxed, possibly also resulting in releasing, depending on the value of their own reference counter. The elements of a multiple value are an

example of the latter kind of lines.

It is recognized that a straightforward use of this reference counter method had the disadvantage that values of a circular structure will not be released. This disadvantage is accepted in the present version, however. To the above description of the use of the link table, an extension is necessary for the correct identification of an identifier, with a given tag, whose applied occurrence appears within a routinetext, since the dynamically last encountered declaration of an identifier with this tag need not be the defining occurrence we have to look for. This problem is solved by making part of the linktable "invisible", in a way similar to a method used for "protection" as will be described in the next chapter.

In this chapter the identification of identifiers was sketched; nearly the same method, including the use of the link table and the value table, applies to the identification of indications. A minor complication occurs in the search for an operator indication, because here the mode of the operands must also be taken into account during the identification.

7. The call

Consider the example

```
proc plus = (int a,b) int: a + b ;  
  
plus (1,2) .
```

In the ALGOL 68 Report, the elaboration of a call is described by means of replacing the skip symbols of a routine by the corresponding actual parameters, thus constructing an identity declaration. The interpreter actually constructs such an identity declaration on the elaboration of a call, for which purpose a free part of the array X , immediately beyond the final program symbol $X[n]$, is used.

So for the above example, the identity declaration

```
int a=1 , b=2
```

will be constructed.

Then this identity declaration is elaborated like one that is contained in the program text. To do so, the program pointer p makes an "excursion", i.e. its value is saved, p is set to $n+1$, and afterwards p will be reset to the saved value. Next, the remaining part, i.e.

```
int:a+b ,
```

which is considered to be a cast, is elaborated by means of a similar excursion of p .

Of course, this description is too simplistic, as is shown by the following example:

```
int a:=1, b:=2;  
proc plus = (int a,b) int : a+b;  
plus (b,a) ,
```

where the identity declaration

int $a=b$, $b=a$

would arise, which is obviously wrong. Therefore the Report describes a "protection" process.

In the interpreter the protection problem is solved by making the formal parameters "invisible" when they enter the link table, which means that instead of the small integral values l^a and l^b in our example, the values $l^a+10000$ and $l^b+10000$, respectively, are taken. The term "invisible" is chosen because the incremented value of l^a is not "seen" on the search in the link table activated by the second part

$b=a$

of the identity declaration above, thus enabling the searching process to find the unaffected value l^a . After the completion of the identity declaration the formal parameters are made visible, by decreasing their l -values by 10000. So the formal parameters need not be changed in the remaining part of the routine.

8. The jump

It has often been stressed that frequent use of explicitly written jumps does generally not improve the quality of programs. Therefore there is a tendency to reduce the use of goto statements in ALGOL-like languages. Typically, the interpreter itself contains only a single goto statement, used as an exit in the case of an abnormal termination of the program. Although it was proved by Van Wijngaarden [5] that even in ALGOL 60, a jump can always be avoided, it is an ALGOL 68 concept and it has therefore been implemented in the interpreter. Basically, a jump results in giving the program pointer p a new value. We have to distinguish, however, between a local jump of the form

..... ; goto l_1 ; ... ; l_1 : ...

and a non-local jump like

...; (...; goto l_2 ;...);...);...; l_2 :...

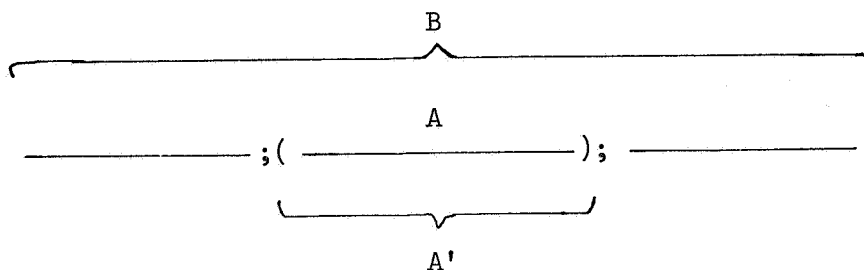
Especially for these forward jumps, it is appropriate to deal with defining occurrences of labels in the scan with $q=h$, i.e. before the more active scan with $q=g$.

If such an occurrence is found, the value l representing the label internally is put into the link table, together with the current value of the program pointer p . The latter value is supplied with a minus sign before putting it into column V , to be able to recognize it as a label. Moreover, in this label-searching scan the program pointer value p_2 pointing to the end of the serial clause is found, which we will need later on. During this scan, a new range, e.g. a closed clause, is skipped, since labels within it are not interesting at this stage. If now in the next scan (with $q=g$) of the serial clause a jump is encountered, the label involved is looked up in the link table, yielding the stored value of p . Now the jump is local if

$$p_1 < p < p_2 \tag{1}$$

where p_1 and p_2 are the starting and ending point of the serial clause, respectively. In this case, the found value of p replaces the current value of the program pointer and the scan is continued.

If (1) does not hold, however, the jump is non-local. Simply changing the program pointer is not allowed now, since this would have the following effect. The program text after the defining occurrence of the label involved would be treated as if it was contained in the smallest serial clause containing the jump. Therefore something more complicated has to be done now. To sketch the treatment of both the local and the non-local jump in more detail, the reader has to remember that the interpreter contains a set of Boolean procedures calling each other in a hierarchical way. The appropriate place for making the test (1) and changing the program pointer in the case of a local jump, is the Boolean procedure *SERIAL CLAUSE*, because p_1 and p_2 are local variables of it. The occurrence of a jump, however, is detected in the Boolean procedure *JUMP* and the normal situation is that the elaboration of a notion takes place in the Boolean procedure corresponding to that notion. The former of these procedures does not directly call the latter, like it is not customary in an army that a general gives a direct command to a soldier. Similarly, information from a soldier to a general usually has to pass a large number of intermediate levels. In our case the parameter q is used as a channel for this latter type of information transport, i.e. from *JUMP* to *SERIAL CLAUSE*. It passes in this upward direction through the procedures *HIP*, *PRIMARY*, *SECONDARY*, *TERTIARY*, and *UNIT*. Consequently the parameter q is a call-by-name parameter in ALGOL 60 terms. Now suppose that, in the Boolean procedure *SERIAL CLAUSE* the occurrence of a jump is perceived after a call of *UNIT* by comparing the values of q before and after this call, and that the found value of p does not satisfy (1), i.e. the jump is non-local. Furthermore, suppose that the serial clause under consideration, which we will call A , is preceded by an open symbol and followed by a close symbol, thus forming a closed clause A' , which is in turn contained in a serial clause B . So the situation is:



Now by virtue of the occurrence of a non-local jump in A, the remaining part of A is skipped, but the parameter q is left in its modified state. This means that a former but still existing activation of the Boolean procedures *SERIAL CLAUSE*, on analysing B, perceives the occurrence of a jump. Since $p1$ and $p2$ are local variables of *SERIAL CLAUSE*, test (1) will now be made for the wider range B etc.

9. Results and conclusions

Both ALGOL 68 and the interpreter discussed in the preceding chapters, are still subject to modifications. Therefore the present version of this interpreter, listed in appendix A, will not be the final one, but may be regarded as a first approximation. It consists of two separate programs. The first performs the lexical scan and produces output for the second. The interpreter has run successfully on a number of small but non-trivial ALGOL 68 programs, a few of which are listed in appendix B. The total time used by the Electrologica X8 computer for the listed programs was about 12 minutes.

As to the aspects mentioned in the introduction, the results are not discouraging if one does not assign too much weight to the requirements of efficiency, error handling and completeness. The most serious drawback for large-scale applications is the extensive use of core memory. The error messages are rather cryptical and situations may arise where the interpreter will not report an error but will deliver unpredictable results instead, when it is dealing with an incorrect program.

The language ALGOL 68 was not implemented completely. The implementation was not only based on the original Report [1], but also on a number of proposed modifications which will probably be accepted, such as the abandonment of the proceduring coercion.

Apart from these language revisions, a number of restrictions were introduced, including the following:

- no bits or bytes structures are available;
- no synchronization operations are available;
- no distinction is made between local and global generators;
- applied occurrences of mode indications must not dynamically precede their defining occurrences;
- as transput procedures, only print, newpage, newline and space are available in the present version; the layout of a printed number may differ from the description in the Report [1].

A more exact and definitive list of restrictions will be made when both a revised ALGOL 68 report has been published and some shortcomings of the interpreter have been eliminated.

The choice of ALGOL 60 and the absence of a machine dependent object code make the interpreter accessible, in principle, for both human beings and machines that can read ALGOL 60. As a final remark, it should be mentioned that some concepts of this language, such as blocks, recursion and conditional expressions have proved to be extremely facilitating tools in writing this interpreter.

References

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```

BEGIN COMMENT ALGOL 68 INTERPRETER, PART 1, L. AMMERRAL, MACHINE DEPENDENT PROCEDURES: FIXT(M,N,X) -- PRINTS X WITH M
DIGITS BEFORE AND N DIGITS AFTER THE POINT, ABSFIXT(M,N,X) -- LIKE FIXT BUT WITHOUT SIGN, NLCR -- NEW LINE
CARRIAGE RETURN, NEWPAGE -- GIVES A NEW PAGE, PRINT(X) -- PRINTS X IN A STANDARD FORM, PRINTTEXT($*) --
PRINTS THE STRING $, EXIT -- TERMINATES THE PROGRAM EXECUTION, RESYM -- READS A CHARACTER AND DELIVERS ITS
INTEGRAL REPRESENTATION, (RESYM IS AN INTEGER PROCEDURE WITHOUT PARAMETERS) PRSYM(N) -- PRINTS THE CHARACTER
WITH INTEGRAL REPRESENTATION N, CSYM(N) -- THE PUNCHING COUNTERPART OF PRSYM, CPUNCH(X) -- THE PUNCHING
COUNTERPART OF PRINT, FIXC(M,N,X) -- THE PUNCHING COUNTERPART OF FIXT)
INIEGER I;
EOR I:= I WILLE IREU QQ
BEGIN
    BEGIN SOURCEMENT PAPERTAPE TO MAGN,TAPE,L.AMMERRAL;
        BOOLEAN SEENI, SUC, NORMAL;
        INIEGER NRSNGI;
        INIEGER ABSV, AGN, AND, API, ARG, POLDI, AT, BEC, BOOL, BUS, BY, BYTES, CARRET, CL, CLI, CHAR, COL,
        COLI, COM, COMI, COMPL, CONJ, DIV, DO, DUM, EIT, ELEM, ENT, EQ, FALSE, FLEX, FOR, FROM, GE, GOON,
        GOONI, GOTO, GT, H, HEAP, I, IM, IND, INT, INTR, INTERROGATION, ISS, ISNT, J, L, LE, LESS, LMIS,
        LOC, LRNR, LWB, MOD, MODE, MIN, N, NA, NE, NIL, NLINE, NMIS, NOT, NO, NNR, NX, ODD, OF, OP, OPEN,
        OPENI, OR, OVER, P, PLIN, PLIT, PLUS, PNT, PNTI, PRIO, PROC, PX, QUOTE, QUOTEI, RE, REAL, REF,
        REPR, RND, S, SGN, SKP, STR, STRING, SPACE, SUB, TAB, TAG, TEN, TMS, TO, TPNR, TRUE, UN, UP, UPB,
        VERT, VERTI, VOID, WHL, TMAB, PSAB, PLTO, MAB, DVAB, MDAB, PLUSI, EGI, MINI, TIMSI, DIVI,
        PCT, NOTI, LESSI, GTI, ORI, ANDI, ZA, ZB, ZC, ZD, ZE, ZF, ZG, ZH, ZI, ZJ, ZK, ZL, ZM, ZN, ZO, ZP,
        ZQ, ZR, ZS, ZT, ZU, ZV, ZW, ZX, ZY, ZZ)
        REAL REALNR;
    END
    PROCEDURE PROUT(A); INIEGER ABBAY A;
    BEGIN NLCR; CSYM(119); CSYM(119);
    EOR I:= 0 STEP 1 UNTIL A(0) DO
        BEGIN I:= I - 1 I 10 * 10 = 1 THEN
            BEGIN NLCR; CSYM(119) END;
            FIXT(8, 0, A(I)); FIXC(8, 0, A(I))
        END
    END;
END;

```

```

155I 67; ISNT:= 88; BECI:= 84; AGN:= 132; OPENI:= 110; CL:= 119; PLIT:= 63; COLI:= 128; GOONI:= 163;
COMI:= 122; QUOTEI:= 174; VERTI:= 131; PNTI:= 30; FORI:= 168; FROMI:= 170; TOI:= 171; WHLI:= 172;
DOI:= 173; INTI:= 101; REALI:= 102; BOOLI:= 103; CHAR:= 104; VOIDI:= 92; STRI:= 95; COMPLI:= 106; TRUEI:= 38;
FALSEI:= 39; REFI:= 96; FLEXI:= 97; EITI:= 98; PROCI:= 99; UNI:= 100; STRINGI:= 109; MODEI:= 112; PRIOI:= 113;
LOCI:= 114; HEAPI:= 115; OPI:= 116; ATI:= 129; OFI:= 161; GOTOI:= 165; SKPI:= 166; NILI:= 167; NRSNGI:= 125;
ENTI:= 400031; ABSVI:= 400027; ARG:= 400035; CONJI:= 400034; LMI:= 400033; LWB:= 400024; ODDI:= 400028;
REI:= 400032; REPRI:= 400036; RNDI:= 400030; SNGI:= 400029; UPBI:= 400025; UP:= 400023; TMABI:= 400004;
TIMSI:= 400018; PSABI:= 400002; PLTO:= 400001; PLUSI:= 400016; MABI:= 400003; MINI:= 400017;
DVABI:= 400007; DIVI:= 400019; MDABI:= 400006; MODI:= 400021; OVABI:= 400005; OVERI:= 400020; NEI:= 400011;
NOTI:= 400026; LEI:= 400014; LESSI:= 400012; GEI:= 400015; GTI:= 400013; ORI:= 400008; ANDI:= 400009;
EOR:= 400010; PLITI:= 400022; ZAI:= 10; ZBI:= 11; ZCI:= 12; ZDI:= 13; ZEI:= 14; ZFI:= 15; ZGI:= 16; ZHI:= 17;
ZII:= 18; ZJI:= 19; ZKI:= 20; ZLI:= 21; ZMI:= 22; ZNI:= 23; ZO:= 24; ZPI:= 25; ZSI:= 26; ZRI:= 27; ZSI:= 28;
ZTI:= 29; ZU:= 30; ZVI:= 31; ZWI:= 32; ZXI:= 33; ZYI:= 34; ZZ:= 35; TAGI:= 3; INDI:= 4; API:= 120;
PLUSI:= 64; MMINI:= 65; LESSI:= 12; COMI:= 11; PNTI:= 88; TABI:= 89; QUOTEI:= 121; NOTI:= 76; GTI:= 74;
INTERROGATION:= 122; COLI:= 90; VERTI:= 127; SPACEI:= 93; TABI:= 118; PCTI:= 132; ORI:= 79; ANDI:= 80;
CARRETI:= 119; NXI:= 1000; NMISI:= 1000; NNRNI:= 500; NAIS:= 5000; NQI:= 200; OPENI:= 98; CLI:= 99;
GOONI:= 91; NLINEI:= 500; EGI:= 70; TIMSI:= 66; DIVI:= 67; SUBI:= 100; BUSI:= 101;
BEGIN INIEGER ABBAY X(0:MX), MISC(0:NMISI), A(1:NA), Q(1:NQ);
REAL ABBAY RNR(0:NRNR);
PROCEDURE OUT(I); VALUE I; INIEGER I;
BEGIN PX:= PX + 1; LE PX > NX THEN ER(ARRAY X EXC I) X(PX) := I END;
PROCEDURE ER(S); S:BEGIN S;
BEGIN NLCR; NLCR;

```

```

FOR I:= 1 STEP 1 UNTIL PX DO PRINT(X(I)); NLCR; NLCR; NLCR;
FOR I:= 1 STEP 1 UNTIL LMS DO PRINT(MISC(I)); NLCR; NLCR; NLCR; PRINTTEXT({ERROR I});
PRINTTEXT(S); EXIT
END;

```

```

BOOLEAN PROCEDURE LETGITS; LE BETWEEN(10, 9, 35) THEN
BEGIN INTEGER H;
LETGITS:= TRUE; N:= 1; Q(1):= S;
FOR I:= 1 WHILE BETWEEN(0, S, 35) DO
BEGIN N:= N + 1; Q(N):= S END;
LE NEWITEM(H) THEN
BEGIN LMS:= LMS + 1; LE LMS + N > NMIS THEN ER({MISC.ARRAY EXCT}); MISC(LMS):= N;
FOR J:= 1 STEP 1 UNTIL N DO MISC(LMS + J):= Q(J); OUT(10000 * TAG + LMS);
LMS:= LMS + N
END
ELSE OUT(10000 * TAG + H)

```

```

END
ELSE LETGITS:= FALSE;
PROCEDURE INTOUT(I); VALUE I; INTEGER I;
BEGIN INTEGER H;
N:= 1; Q(1):= I; LE NEWITEM(H) THEN
BEGIN LMS:= LMS + 2; LE LMS > NMIS THEN ER({MISC.ARRAY EXCT}); MISC(LMS - 1):= I;
MISC(LMS):= I
END;
OUT(10000 * INT + H)
END;

```

```

BOOLEAN PROCEDURE NEWITEM(H); INTEGER H;
BEGIN COMMENT IF ARRAY Q CONTAINS AN OLD ITEM THEN H WILL BE SET SUCH THAT IT POINTS TO THAT
OLD ITEM IN MISC, OTHERWISE H WILL POINT TO THE FIRST FREE CELL IN MISC;

```

```

BOOLEAN PROCEDURE DIFFERENT(N); VALUE N; INTEGER N;
BEGIN INTEGER J;
FOR I:= 1 WHILE (LE J > N THEN FALSE ELSE MISC(H + J) = Q(I)) DO J:= J + 1;
DIFFERENT:= J < N
END;

```

```

H:= 1;
FOR I:= 1 WHILE (LE MISC(H) = - 1 THEN FALSE ELSE (LE MISC(H) = N THEN TRUE ELSE
DIFFERENT(N) DO HI:= H + MISC(H) + 1; NEWITEM:= MISC(H) = - 1
END;

```

```

PROCEDURE REALOUT(R); VALUE R; REAL R;
BEGIN INTEGER H;
H:= 1;
FOR I:= 1 WHILE (LE H > LRNR THEN FALSE ELSE RNR(H) * R DO HI:= H + 1; LE H > LRNR THEN
BEGIN LRNR:= H; RNR(H):= R END;
PX:= PX + 1; X(PX):= 10000 * REAL + H
END;

```

```

PROCEDURE ON(M); VALUE M; INTEGER M;
BEGIN INTEGER H, J;
N:= M;
FOR J:= 1 STEP 1 UNTIL M DO Q(J):= A(P = M + J - 1); LE NEWITEM(H) THEN

```

```

BEGIN LMIS:= LMIS + 1; IF LMIS > N > NMIS THEN ER(MISC, EXC.); MISC(LMIS):= M;
FOR J:= 1 STEP 1 UNTIL M DO MISC(LMIS + J):= 0; J); OUT(100000 + IND + LMIS);
LMIS:= LMIS + M

```

```

END
ELSE OUT(100000 + IND + H)

```

END;

```

BOOLEAN PROCEDURE REAL DENOTATION(R); REAL R;
BEGIN INTEGER POLD;

```

```

REAL RR;
POLD:= P; IF FLOATING(RR) THEN
BEGIN REAL DENOTATION:= ISSUE; R:= RR; END.
ELSE IF VARPNT(RR) THEN
BEGIN REAL DENOTATION:= ISSUE; R:= RR; END.
ELSE
BEGIN REAL DENOTATION:= EALSE; P:= POLD; END

```

END;

```

BOOLEAN PROCEDURE VARPNT(RR); REAL RR;

```

```

BEGIN INTEGER NDU, INUM, N, POLD;
POLD:= P; IF INT DENOTATION(INUM, NDU) THEN RR:= INUM ELSE RR:= 0;
IF (LE ISPNT) THEN INT DENOTATION(INUM, N) ELSE FALSE; THEN
BEGIN RR:= RR + INUM / 10 + N; VARPNT:= ISSUE; END
ELSE
BEGIN VARPNT:= EALSE; P:= POLD; END

```

END;

```

BOOLEAN PROCEDURE EXPART(IEXP); INTEGER IEXP;

```

```

BEGIN INTEGER POLD, ISIGN, INUM, NDU;
BOOLEAN SUC;
POLD:= P; SUC:= EALSE; IF IS(TEN) THEN
BEGIN ISIGN:= LE IS(PLUS) THEN 1 ELSE IF IS(MIN) THEN -1 ELSE 1;
SUC:= INTDENOTATION(INUM, NDU)

```

```

END;
IF SUC THEN IEXP:= ISIGN * INUM ELSE P:= POLD; EXPART:= SUC

```

END;

```

BOOLEAN PROCEDURE FLOATING(RR); INTEGER RR;

```

```

BEGIN BOOLEAN SUC;
INTEGER POLD, INUM, NDU, IEXP;
REAL RRR;
SUC:= EALSE; POLD:= P; IF VARPNT(RRR) THEN
BEGIN IF EXPART(IEXP) THEN
BEGIN RR:= RRR; SUC:= ISSUE; END

```

```

END
ELSE IF INTDENOTATION(INUM, NDU) THEN
BEGIN IF EXPART(IEXP) THEN
BEGIN RR:= INUM; SUC:= ISSUE; END

```

```

END;
IF SUC THEN RR:= RR * 10 + IEXP ELSE P:= POLD; FLOATING:= SUC

```

END;

```

BOOLEAN PROCEDURE INTDENOTATION(I, N); INTEGER I, N;

```

```

BEGIN INTEGER S, POLD;
POLD:= P; IF BETWEEN(0, S, 9) THEN

```

```

BEGIN I:= S; N:= 1;
FOR DUM:= DUM WHILE BETWEEN(0, S, 9) DO.
  BEGIN I:= 10 * I + S; N:= N + 1 END;
INTDENOTATION:= ISUE
END
ELSE
BEGIN INTDENOTATION:= EALSE; P:= POLD END
END;

BOOLEAN PROCEDURE BETWEEN(MIN, S, MAX); VALUE MIN, MAX; INTEGER MIN, MAX, S;
BEGIN FOR DUM:= DUM WHILE A(P) = CARRET ∨ A(P) = SPACE OR P:= P + 1;
  LE A(P) ≥ MIN ∧ A(P) ≤ MAX THEN
  BEGIN S:= A(P); P:= P + 1; BETWEEN:= ISUE END
  ELSE BETWEEN:= EALSE
END;

BOOLEAN PROCEDURE IS(C); VALUE C; INTEGER C; LE A(P) = C THEN
  BEGIN P:= P + 1; IS:= ISUE END
  ELSE IS:= EALSE;

BOOLEAN PROCEDURE IS2(C1, C2); VALUE C1, C2; INTEGER C1, C2;
BEGIN INTEGER POLD;
  BOOLEAN SUC;
  POLD:= P; SUC:= LE IS(C1) THEN IS(C2) ELSE EALSE; LE -SUC THEN P:= POLD; IS2:= SUC
END;

BOOLEAN PROCEDURE IS3(C1, C2, C3); VALUE C1, C2, C3; INTEGER C1, C2, C3;
BEGIN INTEGER POLD;
  BOOLEAN SUC;
  POLD:= P; SUC:= LE IS(C1) THEN IS2(C2, C3) ELSE EALSE; LE -SUC THEN P:= POLD; IS3:= SUC
END;

BOOLEAN PROCEDURE IS4(C1, C2, C3, C4); VALUE C1, C2, C3, C4; INTEGER C1, C2, C3, C4;
BEGIN INTEGER POLD;
  BOOLEAN SUC;
  POLD:= P; SUC:= LE IS(C1) THEN IS3(C2, C3, C4) ELSE EALSE; LE -SUC THEN P:= POLD;
  IS4:= SUC
END;

BOOLEAN PROCEDURE IS8(C1, C2, C3, C4, C5, C6, C7, C8); VALUE C1, C2, C3, C4, C5, C6, C7, C8;
INTEGER C1, C2, C3, C4, C5, C6, C7, C8;
BEGIN INTEGER N, POLD;
  BOOLEAN SUC;
  POLD:= P; N:= LE C6 = 0 THEN 5 ELSE LE C7 = 0 THEN 6 ELSE LE C8 = 0 THEN 7 ELSE 8;
  SUC:= LE IS4(C1, C2, C3, C4) THEN (LE N = 5 THEN IS(C5) ELSE LE N = 6 THEN IS2(C5, C6)
  ELSE LE N = 7 THEN IS3(C5, C6, C7) ELSE IS4(C5, C6, C7, C8)) EALSE EALSE;
  LE -SUC THEN P:= POLD; IS8:= SUC
END;

BOOLEAN PROCEDURE IS12(C1, C2, C3, C4, C5, C6, C7, C8, C9, C10, C11, C12);
VALUE C1, C2, C3, C4, C5, C6, C7, C8, C9, C10, C11, C12;
INTEGER C1, C2, C3, C4, C5, C6, C7, C8, C9, C10, C11, C12;
BEGIN INTEGER N, POLD;
  BOOLEAN SUC;

```

```

POLD:=P) N1:=LE C10 = 0 THEN 9 ELSE LE C11 = 0 THEN 10 ELSE LE C12 = 0 THEN 11 ELSE 12)
SUC:= LE 188(C1, C2, C3, C4, C5, C6, C7, C8) THEN (IF N = 9 THEN 19(C9) ELSE IF N = 10
THEN 19(C9, C10) ELSE IF N = 11 THEN 19(C9, C10, C11) ELSE 19(C9, C10, C11, C12))
ELSE FALSE; LE ~SUC THEN P:= POLD; IS 12:= SUC
END)

```

```

BOOLEAN PROCEDURE ISU(C); VALUE C; INTEGER C1 ISU1= IS2(C, AP1)

```

```

BOOLEAN PROCEDURE ISU2(C1, C2); VALUE C1, C2; INTEGER C1, C2; ISU2:= IS3(C1, C2, AP1)

```

```

BOOLEAN PROCEDURE ISU3(C1, C2, C3); VALUE C1, C2, C3; INTEGER C1, C2, C3;
ISU3:= IS4(C1, C2, C3, AP1)

```

```

BOOLEAN PROCEDURE ISU4(C1, C2, C3, C4); VALUE C1, C2, C3, C4; INTEGER C1, C2, C3, C4;
ISU4:= IS5(C1, C2, C3, C4, AP1, 0, 0, 0)

```

```

BOOLEAN PROCEDURE ISU5(C1, C2, C3, C4, C5); VALUE C1, C2, C3, C4, C5; INTEGER C1, C2, C3, C4, C5;
ISU5:= IS6(C1, C2, C3, C4, C5, AP1, 0, 0)

```

```

BOOLEAN PROCEDURE ISU6(C1, C2, C3, C4, C5, C6); VALUE C1, C2, C3, C4, C5, C6;
INTEGER C1, C2, C3, C4, C5, C6; ISU6:= IS7(C1, C2, C3, C4, C5, C6, AP1, 0)

```

```

BOOLEAN PROCEDURE ISU7(C1, C2, C3, C4, C5, C6, C7); VALUE C1, C2, C3, C4, C5, C6, C7;
INTEGER C1, C2, C3, C4, C5, C6, C7; ISU7:= IS8(C1, C2, C3, C4, C5, C6, C7, AP1)

```

```

BOOLEAN PROCEDURE ISU8(C1, C2, C3, C4, C5, C6, C7, C8); VALUE C1, C2, C3, C4, C5, C6, C7, C8;
INTEGER C1, C2, C3, C4, C5, C6, C7, C8; ISU8:= IS12(C1, C2, C3, C4, C5, C6, C7, C8, AP1, 0, 0, 0)

```

```

P:= 0; NEWPAGE; NORMAL:= TRUE;
FOR I:= 1, 2, 3, 4 DO CSYM(19); CSYM(1); S1:= RESYM;
FOR J:= 1 WHILE S # INTERROGATION DO
BEGIN PRSYM(S); CSYM(5); LE S = CARRET ^ NORMAL ^ S = SPACE THEN ELSE
BEGIN LE S = AP1 THEN NORMAL:= ~NORMAL; P:= P + 1; LE P > NA THEN ER(
$SOURCE PROGRAM TOO LONG$); A(P) := S
END;
S:= RESYM
END;

```

```

A(P + 1) := INTERROGATION; NA:= P; PX:= LNRR:= 0; PLINI:= 2; P:= 1;
FOR HI:= 1 SIEE 1 UNILL NMIS DO MISC(HI) := -1; MISC(1) := 9; LMIS:= 36;
FOR J:= 2, ZP, ZI, 4, ZS, ZQ, ZR, ZT, 3, ZE, ZX, ZY, 2, ZL, ZN, 3, ZC, ZO, ZB, 6, ZA, ZR, ZC,
ZC, ZO, ZS, 3, ZS, ZI, ZN, 6, ZA, ZR, ZC, ZS, ZI, ZN, 3, ZT, ZA, ZN, 6, ZA, ZR, ZC, ZT, ZA,
ZM, 5, ZP, ZR, ZI, ZN, ZT, 2, ZR, ZE, 2, ZI, ZM, 7, ZN, ZE, ZW, ZP, ZA, ZG, ZE, 7, ZN, ZE,
ZM, ZL, ZI, ZN, 5, ZS, ZP, ZA, ZC, ZE DO
BEGIN LMIS:= LMIS + 1; MISC(LMIS) := J; END;

```

```

FOR I:= 1 WHILE P $ NA DO
BEGIN SUC:= SEEN1:= TRUE; LE IS(SPACE) THEN J POLD:= P;
LE LETGITS THEN ELSE LE REALDENOTATION(REALNR) THEN REALOUT(REALNR) ELSE LE
INTDENOTATION(INTNR, DUM) THEN INTOUT(INTNR) ELSE LE IS(AP1) THEN
BEGIN LE IS(ZA) THEN

```



```

BEGIN LE ISU2(ZB, ZS) THEN OUT(ABSV) ELSE LE ISU2(ZR, ZG) THEN OUT(ARG) ELSE LE
ISU1(ZT) THEN OUT(AT) ELSE SUCI:= FALSE
END
ELSE LE IS(ZB) THEN
BEGIN LE ISU4(ZE, ZG, ZI, ZN) THEN OUT(OPEN) ELSE LE ISU3(ZO, ZO, ZL) THEN
OUT(BOOL) ELSE LE ISU(ZY) THEN OUT(BY) ELSE SUCI:= FALSE
END
ELSE LE IS(ZC) THEN
BEGIN LE ISU3(ZA, ZS, ZE) THEN OUT(OPEN) ELSE LE ISU3(ZH, ZA, ZR) THEN OUT(CHAR)
ELSE LE ISU(ZO) THEN ISU6(ZO, ZM, ZM, ZE, ZN, ZT) THEN
BEGIN FOR I:= 1 WHILE ~(LE ISU8(API, ZC, ZO, ZM, ZM, ZE, ZN, ZT) THEN ISU6 ELSE
ISU3(API, ZC, ZO)) DO
BEGIN P:= P + 1) LE P > NA THEN ER(COMMENT) END
END
ELSE LE ISU4(ZO, ZM, ZP, ZL) THEN OUT(COMPL) ELSE LE ISU3(ZO, ZN, ZJ) THEN
OUT(CONJ) ELSE SUCI:= FALSE
END
ELSE LE ISU2(ZD, ZO) THEN OUT(DO) ELSE LE IS(ZE) THEN
BEGIN LE ISU3(ZL, ZS, ZE) THEN OUT(VERT) ELSE LE ISU3(ZL, ZS, ZP) THEN OUT(AGN)
ELSE LE ISU2(ZN, ZO) THEN OUT(CL) ELSE LE ISU3(ZI, ZI, ZH, ZE, ZR) THEN
OUT(FT) ELSE LE ISU3(ZN, ZT, ZI, ZE, ZR) THEN OUT(ENT) ELSE LE ISU3(ZS, ZA,
ZC) THEN OUT(CL) ELSE LE ISU3(ZX, ZI, ZT) THEN OUT(PNT) ELSE SUCI:= FALSE
END
ELSE LE IS(ZF) THEN
BEGIN LE ISU4(ZA, ZL, ZS, ZE) THEN OUT(FALSE) ELSE LE ISU(ZI) THEN OUT(CL) ELSE LE
ISU3(ZL, ZE, ZM) THEN OUT(FLEX) ELSE LE ISU2(ZO, ZR) THEN OUT(FOR) ELSE LE
ISU3(ZR, ZO, ZM) THEN OUT(PROM) ELSE SUCI:= FALSE
END
ELSE LE ISU4(ZG, ZO, ZT, ZO) THEN OUT(GOTO) ELSE LE ISU4(ZH, ZE, ZA, ZP) THEN
OUT(HEAP) ELSE LE IS(ZI) THEN
BEGIN LE ISU3(ZS, ZN, ZT) THEN OUT(ISNT) ELSE LE ISU2(ZN, ZT) THEN OUT(INT) ELSE
LE ISU(ZN) THEN OUT(VERT) ELSE LE ISU(ZM) THEN OUT(IM) ELSE LE ISU(ZP) THEN
OUT(OPEN) ELSE LE IS(API) THEN OUT(PLIT) ELSE SUCI:= FALSE
END
ELSE LE IS(ZL) THEN
BEGIN LE ISU2(ZD, ZC) THEN OUT(LOC) ELSE LE ISU2(ZW, ZB) THEN OUT(LWB) ELSE SUCI:=
FALSE
END
ELSE LE ISU4(ZM, ZO, ZD, ZE) THEN OUT(MODE) ELSE LE ISU3(ZN, ZI, ZL) THEN OUT(NIL)
BEGIN LE ISU2(ZD, ZD) THEN OUT(ODD) ELSE LE ISU(ZP) THEN OUT(OP) ELSE LE ISU(ZP)
THEN OUT(OP) ELSE LE ISU3(ZU, ZS, ZE) THEN OUT(AGN) ELSE LE ISU2(ZU, ZT)
THEN OUT(VERT) ELSE SUCI:= FALSE
END
ELSE LE IS(ZP) THEN
BEGIN LE ISU7(ZR, ZI, ZO, ZR, ZI, ZT, ZY) THEN OUT(PRIO) ELSE LE ISU3(ZR, ZO, ZC)
THEN OUT(PROC) ELSE SUCI:= FALSE
END
ELSE LE IS(ZR) THEN
BEGIN LE ISU3(ZE, ZA, ZL) THEN OUT(REAL) ELSE LE ISU2(ZE, ZF) THEN OUT(REF) ELSE
LE ISU(ZE) THEN OUT(RE) ELSE LE ISU3(ZE, ZP, ZR) THEN OUT(REPR) ELSE LE
ISU4(ZO, ZU, ZN, ZD) THEN OUT(RND) ELSE SUCI:= FALSE
END
ELSE LE IS(ZS) THEN
BEGIN LE ISU3(ZI, ZG, ZN) THEN OUT(SGN) ELSE LE ISU3(ZK, ZI, ZP) THEN OUT(SKIP)
ELSE LE ISU5(ZT, ZR, ZI, ZN, ZG) THEN OUT(STRING) ELSE LE ISU5(ZT, ZR, ZU,
ZC, ZT) THEN OUT(STR) ELSE SUCI:= FALSE
END
ELSE LE IS(ZT) THEN
BEGIN LE ISU3(ZH, ZE, ZN) THEN OUT(VERT) ELSE LE ISU(ZO) THEN OUT(TO) ELSE LE

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```

ISU3(ZR, ZU, ZE) THEN OUT(TOUE) ELSE SUCI# EASE#
END
ELSE LE IS(ZU) THEN
BEGIN LE ISU4(ZN, ZI, ZO, ZN) THEN OUT(UN) ELSE LE ISU2(ZP, ZB) THEN OUT(UPB) ELSE
SUCI# EASE#
END
ELSE LE ISU4(ZV, ZO, ZI, ZO) THEN OUT(VOID) ELSE LE ISU5(ZW, ZM, ZI, ZL, ZE) THEN
OUT(WHL) ELSE SUCI# EASE# LE -SUC IREN P# POLDI
END
ELSE LE IS(OPEN1) THEN IRUE ELSE IS(SUB) THEN OUT(OPEN) ELSE LE LE IS(CLI) THEN IRUE
ELSE IS(BUS) THEN OUT(CL) ELSE LE IS(GOON1) THEN OUT(GOON) ELSE LE IS(COM1) THEN
OUT(COM) ELSE LE IS(COL1) THEN OUT(LE IS(EG1, COL1) THEN ISS ELSE LE IS(EG1) THEN BEC
ELSE COL) ELSE LE IS(VERT1) THEN OUT(LE IS(COL1) THEN AGN ELSE VERT) ELSE LE
IS(TIM1) THEN OUT(LE IS(TIM1) THEN UP ELSE LE IS(COL1, EOI) THEN TMAB ELSE TIM)
ELSE LE IS(PLUS1) THEN OUT(LE IS(COL1, EOI) THEN PRAB ELSE LE IS(EG1, COL1) THEN
PLD ELSE PLUS) ELSE LE IS(MINI) THEN OUT(LE IS(COL1, EOI) THEN MAB ELSE MIN) ELSE LE
IS(DIVI) THEN OUT(LE IS(COL1, EOI) THEN DVAB ELSE DIV) ELSE LE IS(PCT) THEN OUT(LE
IS(TIM1) THEN (LE IS(COL1, EOI) THEN MDAB ELSE MOD) ELSE LE IS(COL1, EOI) THEN OVAB
ELSE OVER) ELSE LE IS(NOT1) THEN OUT(LE IS(EG1) THEN NE ELSE NOT) ELSE LE IS(LESS1)
THEN OUT(LE IS(EG1) THEN LE ELSE LESS) ELSE LE IS(OT1) THEN OUT(LE IS(EG1) THEN GE
ELSE GT) ELSE LE IS(EG1) THEN OUT(EG) ELSE LE IS(PNT1) THEN OUT(PNT) ELSE LE IS(AND1)
THEN OUT(AND) ELSE LE IS(PNT1) THEN OUT(PNT) ELSE LE IS(NRSGN) THEN
BEGIN FOR I:= 1 WHILE -IS(NRSGN) DO
BEGIN P# P + 1; LE P > NA THEN ER({COMMENT?}) END
END
ELSE LE IS(QUOTE1) THEN
BEGIN LE A(P + 1) = QUOTE1 THEN
BEGIN INTEGER M;
N:= 1; Q(I):= A(P); P:= P + 2; LE NEWITEM(M) THEN
BEGIN LMIS:= LMIS + 2; LE LMIS > NMIS THEN ER({MISC ARRAY EXC.});
MISC(LMIS - 1):= 1; MISC(LMIS):= Q(I)
END;
OUT(10000 * CHAR * M)
END
ELSE
BEGIN INTEGER M;
N:= 0;
FOR I:= 1 WHILE -IS(QUOTE1) DO
BEGIN N:= N + 1; LE N > NO THEN ER({O ARRAY EXC.}); Q(N):= A(P); P:= P + 1 END;
LE NEWITEM(M) THEN
BEGIN LE LMIS + N + 1 > NMIS THEN ER({MISC ARRAY EXC.}); LMIS:= LMIS + 1;
MISC(LMIS):= N;
FOR I:= 1 SIER 1 UNILL N DO MISC(LMIS + I):= Q(I); LMIS:= LMIS + N
END;
OUT(10000 * STRING * M)
END
END
ELSE SEEN1:= EASE# LE -SUC IREN
BEGIN J:= 0; P:= P + 1;
FOR I:= 1 WHILE A(P) # API DO
BEGIN J:= J + 1; P:= P + 1 END;
OW(J); P:= P + 1
END;
LE P > NA ^ SEEN1 ^ SUC IREN
BEGIN NLCR; PRINTTEXT({WRONGI}); ABSFIXT(4, 0, A(P)); ABSFIXT(4, 0, P); P:= P + 1 END;
END;
CSYM(INTERROGATION); CSYM(119); X(0):= PX; MISC(0):= LMIS; RNR(0):= LNR; PROUT(X);
PROUT(MISC); CSYM(119); NLCR; NLCR;
FOR L:= 0 SIER 1 UNILL LNR UU
BEGIN PRINT(RNR(L)); CPUNCH(RNR(L)) END

```

END

END

END

END

```

BEGIN COMMENT ALGOL 68 INTERPRETER, PART 2. L. AMMERAAL, MACHINE DEPENDENT PROCEDURES; FIXT(M,N,X) -- PRINTS X WITH M
DIGITS BEFORE AND N DIGITS AFTER THE POINT, ABSFIXT(M,N,X) -- LIKE FIXT BUT WITHOUT SIGN, NLCR -- NEW LINE
CARRIAGE RETURN, NEWPAGE -- GIVES A NEW PAGE, SPACE(N) -- GIVES N BLANK SPACES ON THE CURRENT LINE, PRINT(X) --
PRINTS X IN A STANDARD FORM, PRINTTEXT(45) -- PRINTS THE STRING S, EXIT -- TERMINATES THE PROGRAM EXECUTION,
READ -- READS A NUMBER, NUMBERS MAY BE SEPARATED BY TWO OR MORE BLANKS (READ IS A REAL PROCEDURE WITHOUT
PARAMETERS) RESYM -- READS A CHARACTER AND DECLIVERS ITS INTEGRAL REPRESENTATION, (RESYM IS AN INTEGER PROCEDURE
WITHOUT PARAMETERS) PRSYM(N) -- PRINTS THE CHARACTER WITH INTEGRAL REPRESENTATION N;
INTEGER PLTC, PSAB, MAB, TMAB, DVAB, OVAB, MDAB, OR, AND, EQ, NE, LESS, GT, LE, GE, PLUS, MIN, TMS, DIV, OVER,
XOD, UP, PLIT, LWB, UPB, NOT, ABSV, ODD, SGN, RAD, ENT, RE, IM, CONJ, REPR, NEXP, NLN, NCOS, NACS,
NSIN, NASN, NTAN, LRE, LIM, STBL, STIN, STRL, STCR, STCR, STNG, RINT, RRL, RSTS, NATN, STRW, STVD, JGLB, PART,
FEI, FEZ, LEN1, LEN2, MASP, NASP, I1, PP, MM, V1, V2, MINL, M1, M2, BV1, BV2;
BOOLEAN UNCO, MORF;
REAL PI, EV1, EV2, RE1, RE2, IM1, IM2;
INTEGER ARRAY ARR1(50);
INTEGER AGN, AT, ATL, BEC, BITS, BOOL, BY, CHAR, CL, COL, COLL, COM, DO, DUM, EIT, CML, F, FALSE, FIRM, FLEX,
FM, FOR, FROM, FS, G, GOON, GOTO, H, HEAP, I, J, IND, INT, IS, ISAT, L, LOC, MEK, MODE, NEW, NFST, NIL, OF,
OP, OPEN, OPROUT, P, PLAS, PNT, PRIO, PRNT, PROC, PX, OOLD, Q, REAL, REF, ROW, SKP, STR, STRING, STRONG, TAG,
TO, TRUE, UN, V, VOID, VRL, WHL, Z, Z1, NA, NX, NL, NMIS, NM1, NNR, NR1;
EQB I := 1 WHILE TRUE DO
START: AGN:= 132; AT:= 129; ATL:= 1; BEC:= 84; BITS:= 107; BOOL:= 103; BY:= 170; CHAR:= 104; CL:= 119;
COL:= 128; COLL:= 2; COM:= 122; DIV:= 56; DO:= 173; DUM:= 1; EIT:= 98; ENT:= 62; EQ:= 82; F:= 3;
FALSE:= 39; FIRM:= 4; FLEX:= 97; FML:= 5; FOR:= 188; FROM:= 169; G:= 6; GOON:= 163; GOTO:= 165;
H:= 7; HEAP:= 115; J:= 0; INT:= 101; IS:= 87; ISNT:= 88; LESS:= 52; MIN:= 81; LOC:= 114;
LWB:= 60; UPB:= 61; MEK:= 8; MODE:= 112; NEW:= 9; NIL:= 167; OF:= 161; OP:= 116; OPEN:= 118;
OPROUT:= 14; OVER:= 57; P:= 1; PLUS:= 78; PNT:= 36; PRIO:= 113; PROC:= 99; REAL:= 102; REF:= 96;
RND:= 74; ROW:= 11; SKP:= 166; STR:= 95; STRING:= 109; STRONG:= 12; TAG:= 3; TMS:= 83; TO:= 174;
TRUE:= 38; UN:= 100; VOID:= 92; VRL:= 131; VRL:= 13; WHL:= 172; MASP:= 20; NAI:= 2000; NX:= 1000;
NL:= 500; NMIS:= 1000; NNR:= 500; NFST:= 10; PLTO:= 42; PSAB:= 43; MAB:= 44; TMAB:= 45; DVAB:= 46;
OVAB:= 47; MOAB:= 48; OR:= 49; AND:= 50; EQ:= 51; NE:= 52; LESS:= 53; GT:= 54; LE:= 55; GE:= 56; PLUS:= 57;
MIN:= 58; TMS:= 59; DIV:= 60; OVER:= 61; MOD:= 62; UPB:= 63; PLIT:= 64; LWB:= 65; UPB:= 66; NOT:= 67;
ABSV:= 68; ODD:= 69; SGN:= 70; RND:= 71; ENT:= 72; REI:= 73; INI:= 74; CONJ:= 75; ARG:= 76; REPR:= 77;
PI:= 4 * ARCTAN(1); NSQR:= 909; NEXP:= 908; NLN:= 907; NCOS:= 906; NACS:= 905; NSINT:= 904;
NASN:= 903; NTAN:= 902; NATN:= 901; CML:= 106;
BEGIN INTEGER ARRAY X(11NX), MISC(11NMIS), A, B, C, D(11NA), LIND, VIND(11NL);
REAL ARRAY E(11NA), RRR(11NR);
INTEGER ARRAY VASP, WASP(11MASP);
PROCEDURE ADMC(JOLD); VALUE JOLD; INTEGER JOLD; COMMENT ADMINISTRATION TO CLOSE A RANGE;
BEGIN FOR J:= J SIZE - 1 WHILE JOLD + 2 DO RX(VIND(J)); J:= JOLD EQQ;
PROCEDURE ADMC(JOLD); INTEGER JOLD; COMMENT ADMINISTRATION TO OPEN A RANGE;
BEGIN JOLD:= J; J:= J + 1; LIND(J):= PART; VIND(J):= 1 END;
BOOLEAN PROCEDURE ANDP(A, B); BOOLEAN A, B; COMMENT ADVANCE P ONLY IF A AND B;
BEGIN INTEGER POLD;
POLD:= P; IF IE A THEN B ELSE FALSE THEN ANDP:= TRUE ELSE
BEGIN ANDP:= FALSE; P:= POLD END;
END;
REAL PROCEDURE ARGC(R, I); VALUE R, I; REAL R, I; COMMENT ARGUMENT OF THE COMPLEX NUMBER (R, I);
ARGC:= IE ABS(R) > ABS(I) THEN ARCTAN(I / R) + PI / 2 * (IE I < 0 THEN SIGN(R) - 1 ELSE 1 * SIGN(R))
ELSE - ARCTAN(R / I) + PI / 2 * SIGN(I);
INTEGER PROCEDURE ASK;

```

```

BEGIN INTEGER VI;
  LE FS = 0 THEN ER(4NO FREE SPACE)} ASKIF.VI:= FS; FS:= D(V); DIV:= 1
END;

```

```

INTEGER PROCEDURE ASKN(N); VALUE N; INTEGER NI;
BEGIN INTEGER V, VO, VI, VN, I;
  BOOLEAN EXH;
  VI:= VI; FS:= FS; VO:= -1; EXH:= V = 0; I:= 1;
  FOR I:= 1 WHILE ~EXH & I < N DO
    BEGIN VN:= D(V); LE VN = 0 THEN EXH:= TRUE ELSE
      BEGIN LE VN = V + 1 THEN I:= I + 1 ELSE
        BEGIN I:= 1; VI:= VN; VO:= V END;
      V:= VN
    END
  END;

```

```

END;
LE EXH THEN ER(4NO FREE SPACE)} ELSE
  BEGIN ASKN:= VI; LE VO = -1 THEN FS:= D(V) ELSE D(VO)= D(V) END;
  FOR V:= VI + N - 1 STEP - 1 UNTIL VI DO D(V)= 1
END;

```

```

PROCEDURE ASTL(O); INTEGER O; COMMENT ASSIGNATION TAIL;
BEGIN INTEGER VD, VC, VGT, VCT, I, VG, VM, VMT, FEC, N, NLT, FEG;
  BEGIN M(AZI)= DEF, 4DEST NO NAME}; VO:= Z; VC:= B(VD); M(UNIT(O), 4NO SOURCE});
  M(COER(STRONG, VC, O), 4ASS MODE}); LE O = G THEN
    BEGIN VG:= Z; MORF:= FALSE; LE A(VC) = ROW THEN
      BEGIN VGT:= B(VG); VCT:= B(VC); NI:= C(VC); LE A(VCT) = 1 THEN
        BEGIN FOR I:= 1 STEP 1 UNTIL N DO M(A(VCT + I)) = A(VGT + I) ^ B(VCT + I) = B(VGT
          + I), 4MULT ASS BOUND});
        END
      ELSE
        BEGIN VM:= NINST(VG); VMT:= B(VN); A(VMT)= A(VCT); B(VD)= VM; RX(VC); VC:= VM;
          VCT:= B(VC)
        END
      END;
    FEG:= B(VGT) + C(VCT) - 1; FEG:= B(VGT) + C(VGT) - 1; NLT:= E(VGT);
    FOR I:= 1 STEP 1 UNTIL NLT DO SPS(FEG - 1 + 1, FEC - 1 + 1)
  END
  ELSE SPS(VG, VC); RX(VG); Z:= VD
END

```

```

END
ELSE M(UNIT(F), 4SOURCE}
END;

```

```

BOOLEAN PROCEDURE BOP(S, Y); VALUE S; INTEGER S1, S2; BOOLEAN B1, B2; LE PP = S THEN
  BEGIN BOP:= TRUE; PUTBOOL(Y) END;
  ELSE BOP:= FALSE;

```

```

PROCEDURE BOP2(S1, B1, S2, B2); INTEGER S1, S2; BOOLEAN B1, B2; LE PP = S1 THEN PUTBOOL(B1) ELSE
  BEGIN M(PP = S2, 4BOOL OP.); PUTBOOL(B2) END;

```

```

BOOLEAN PROCEDURE BOP4(S1, B1, S2, B2, S3, B3, S4, B4); INTEGER S1, S2, S3, S4; BOOLEAN B1, B2, B3, B4;
  BEGIN BOP4:= TRUE;
  LE PP = S1 THEN PUTBOOL(B1) ELSE LE PP = S2 THEN PUTBOOL(B2) ELSE LE PP = S3 THEN PUTBOOL(B3)
  ELSE LE PP = S4 THEN PUTBOOL(B4) ELSE BOP4:= FALSE
  END;

```

Vp 8: ?
Vp 9?

```

BOOLEAN PROCEDURE CAST(Q) INTEGER Q)
BEGIN
  INTEGER T;
  LE LK(DCL(VRL, DUM, F), LE INP(OPEN) THEN IBUE ELSE INP(COL)) THEN
  BEGIN
    CAST:= IBUE; MDCL(VRL, T, Q), $ CAST $}} LE INP(COL) THEN M(UNIT(Q),
    $ NO UNIT IN CAST $} ELSE M(ENCL(Q), $ NO ENCLOSED CL. IN CAST $});
    LE Q = G THEN MORF:= FALSE; M(COER(STRONG, T, Q), $ CAST COERCION $}
  END
  ELSE CAST:= FALSE
END;

PROCEDURE CDIN(Q, GOUT); INTEGER Q, GOUT; COMMENT CONDITION IN CLAUSE;
BEGIN
  BOOLEAN BL;
  INTEGER Q0, Q1;
  BL:= BIZ) = 1; EMP; Q0:= Q1:= LE BL THEN G ELSE F; M(SCL(Q1), $COND. IN $); LE Q1 # Q0 THEN Q1= F;
  GOUT:= LE Q0 = G THEN F ELSE G
END;

PROCEDURE CFIN(Q, GOUT); INTEGER Q, GOUT; COMMENT CONFORMITY IN CLAUSE;
BEGIN
  INTEGER U, V, W, Q0, Q1, Q2, JOLD, L;
  BOOLEAN JMP, SUC, DONE;
  U:= Z; V:= CIU; JMP:= DONE:= FALSE;
  EBB I:= 1, I WHILE INP(COM) DO
  BEGIN
    Q0:= Q1:= LE JMP v DONE THEN F ELSE G; M(INP(OPEN), $OPEN IN CONF. IN CL. $}; ADMO(JOLD);
    MDCL(FML, W, Q1), $DCL IN CONF. IN CL. $}; LE INTAG. L) THEN POS(L, W, Q1);
    SUC:= LE Q1 = G THEN NEG(V, W) ELSE FALSE; Q2:= LE SUC THEN G ELSE F; M(INP(COL),
    $ IN CONF. IN CLAUSE MISSING $}; M(UNIT(Q2), $UNIT IN CONF. IN CLAUSE $});
    LE SUC THEN DONE:= IBUE; LE Q0 = G ^ Q2 = F THEN JMP:= IBUE; ADMO(JOLD)
  END;
  LE JMP THEN Q1= F; RX(U); GOUT:= LE JMP v DONE THEN F ELSE G
END;

PROCEDURE CHCL(Q); INTEGER Q; COMMENT CHOOSER CLAUSE;
BEGIN
  INTEGER AV, GOUT;
  LE G = G THEN
  BEGIN
    EGR I:= 1 WHILE DREF v DPRO(Q) DO ; AV:= AIZ);
    LE AV = BOOL THEN CDIN(Q, GOUT) ELSE LE AV = INT THEN CSIN(Q, GOUT) ELSE
    BEGIN
      M(AV = UN, $CHOISE. MODE $}); CFIN(Q, GOUT); ENQ;
    LE INP(VERT) THEN M(SCL(GOUT),
    $OUT CLAUSE $} ELSE LE INP(AGN) THEN CHCL(GOUT) ELSE LE GOUT = G THEN
    BEGIN
      Z:= ASK; AIZ)= VOID END
    END
  ELSE P:= MATCH - 1
END;

PROCEDURE CLLC(Q); INTEGER Q; COMMENT COLLATERAL CLAUSE ;
BEGIN
  INTEGER N, V, W, I;
  M(INP(COM), $AFTER SER. CLAUSE $}; V:= ASK; AIV)= COLL; C(V); N:= LSTC(UNIT(F)) + 1;
  BIV)= W:= ASKN(N); MOVE(Z, W); EMP;
  EBB I:= 2 SIZE I UNTIL N DO
  BEGIN
    LE I > 2 THEN M(INP(COM), $COLL. CL. COMMA $}; M(UNIT(Q), $COLL. CLAUSE $}); LE Q = G THEN
    BEGIN
      BEGIN MOVE(Z, W + 1 - 1); EMP END
    END;
    Z:= V; M(INP(CL), $ ) MISSING IN COLL. CL. $}
  END;

```

```

PROCEDURE CLT(Q); INTEGERS Q; COMMENT CALL TAIL;
BEGIN INTEGERS V, W, N, PR, JOLD, PX0, I, POLD, J1, J2;
REAL Y;
MORF := TRUE; V := Z; W := B(V); N := B(W); PR := C(W); LE PP := ~ 1000 THEN
  BEGIN M(UNIT(C), {PRINT}); GOUT(Z); Z := V; M(INP(CL), {PRINT CL}); END
ELSE LE PP < ~ 996 THEN
  BEGIN PR := MATCH; LE PP = ~ 999 THEN NEWPAGE ELSE LE PP = ~ 998 THEN NLCR ELSE SPACE(1) END
ELSE LE PP $ NAIN THEN
  BEGIN M(UNIT(Q), {ST,FUN}); M(INP(CL), {ST,F,CL,}); M(COER{STRONG, STRL, Q}, {ARG,ST,FUN});
  LE Q = G THEN
    BEGIN PR := PR; Y := E(Z);
    LE ROP4(NSOR, SORT(Y), NEXP, EXP(Y), NLN, LN(Y), NCOS, COS(Y)) THEN ELSE LE
    ROP4(NACS, ARGC(Y), SORT(1 - Y + Y), NSIN, SIN(Y), NASN, ARGC(SORT(1 - Y + Y),
    Y), NTAJ, SIN(Y) / COS(Y)) THEN ELSE M(ROP(NATJ, ARCTAN(Y)), {ILL,ST,FUN});
  END
END
END
END
ADMO(JOLD); J1 := JGLB - 1; J2 := J;
FOR J1 := J1 + 1 WHILE LIND(J1) # PART QO; INVS(J1, J2); PX0 := PX; PI := P - 1;
FOR I := 1 STEP 1 UNTIL N DO
  BEGIN CTX(PR); WTX(QO); WTX(LE I = N THEN GOON ELSE COM); END;
  POLD := P; P := PX0 + 1; M(LYDQ(Q), {PARAM IN CALL}); PI := POLD; M(INP(CL),
  {MISSING IN CALL}); POLD := P; PI := PR + 1; M(CAST(Q), {CALL}); PI := POLD; PX := PX0;
  VISE(J1, J2); ADM(JOLD)
END
END
END;

```

```

BOOLEAN PROCEDURE COER(SORT, APST, Q); VALUE SORT, APST, Q; COMMENT COERCION;
BEGIN COER := TRUE;
  LE Q # G THEN ELSE LE A(APST) = VOID ^ SORT = STRONG THEN ELSE LE (AIZ) = VOID v AIZ) = REF
  ^ BIZ) = 0) ^ SORT = STRONG THEN
  BEGIN EMP; Z := NINST(APST) END
  ELSE LE AIZ) = COLL THEN DCOL(APST) ELSE
  BEGIN INTEGERS APRI, I, N, V, VO, V1;
  BOOLEAN TRY, DONE;
  REAL R;
  APRI := Z; TRY := TRUE;
  FOR I := 1 WHILE LE TRY THEN ~MEQ(APRI, APST) ELSE FALSE DO
    BEGIN LE SORT = MEK THEN
      BEGIN TRY := LE DREF THEN ISUE ELSE DPRO(Q); APRI := Z END
      ELSE
      BEGIN COMMENT FIRM OR STRONG;
      LE A(APST) = UN THEN
      BEGIN V := B(APST); I := 0; NI = C(V);
      FOR I := 1 WHILE LE I < N THEN ~MEQ(B(V + I), APRI) ELSE FALSE DO (I + 1)
      LE I < N THEN
      BEGIN Z := NINST(APST); C(Z) := NINST(APRI); RX(APRI); APRI := Z END
      ELSE TRY := FALSE
      END
      ELSE LE SORT = FIRM THEN
      BEGIN LE A(APRI) = REF THEN
      BEGIN LE MEQ(BI(APRI), APST) THEN
      BEGIN MDREF, {COERC, DREF,}; APRI := Z END
      ELSE TRY := FALSE
      END
      END
      ELSE LE A(APRI) = PROC THEN
      BEGIN LE MEQ(BI(APRI) + 1, APST) THEN M(DPRO(Q),
      {COERC, DPROCT}) ELSE TRY := FALSE
      END
      END
      END

```

Vp xi ?

```

END
ELSE TRY:= FALSE
END
ELSE LE SORT = STRONG THEN
BEGIN LE DREF THEN APRI:= Z ELSE LE DPRO(9) THEN APRI:= Z ELSE
BEGIN DONE:= FALSE; LE CPL1(APST, REI, IMI) THEN
BEGIN LE APRI:= INT THEN
BEGIN RI:= BIAPRI; PUTREAL(R); RX(APRI); APRI:= Z; END;
LE APRI:= REAL THEN
BEGIN PUTCPL(EAPRI), 0; RX(APRI); DONE:= ISSUE; APRI:= Z; END
END;
LE DONE THEN ELSE
BEGIN LE APST:= REAL ^ A(APRI) = INT THEN
BEGIN RI:= BIAPRI; PUTREAL(R); RX(APRI); APRI:= Z; DONE:= ISSUE; END
ELSE LE A(APST) = ROW ^ A(APRI) = ROW THEN
BEGIN V:= NINST(APRI); N:= C(APST); V0:= ASKN(N + 2); A[V0]:= ROW;
B[V0]:= V; V1:= V0 + 1; C[V0]:= N; A[V1]:= 1; B[V1]:= V; C[V1]:= 1;
FOR I:= 1 STEP 1 UNTIL N DO A[V1 + I]:= B[V1 + I]; C[V1 + I]:= 1;
E[V1]:= 1; RX(APRI); APRI:= Z; V0:= V; DONE:= ISSUE
END
ELSE TRY:= FALSE
END;
LE -DONE THEN TRY:= FALSE
END
END
END
END;
END;
COER:= TRY
END
END
END;
END;
END;
PROCEDURE COPGEN(I, A, B, X, Y); VALUE A, B; INTEGER I, A, B, X, Y
END I:= A STEP 1 UNTIL B DO Y:= X;
END;
END;
PROCEDURE COPY(V, W); VALUE V, W; INTEGER V, W
BEGIN A[W]:= A[V]; B[W]:= B[V]; C[W]:= C[V]; E[W]:= E[V]; END;
END;
PROCEDURE COPYN(V, W, N); VALUE V, W, N; INTEGER V, W, N;
BEGIN INTEGER I, M;
M:= N - 1;
FOR I:= 0 STEP 1 UNTIL M DO COPY(V + 1, W + 1)
END;
END;
BOOLEAN PROCEDURE CPL1(V, RE, IM); VALUE V; INTEGER V; REAL RE, IM;
BEGIN INTEGER BV, V1, V2;
CPL1:= FALSE; LE A[V] = STR THEN
BEGIN LE C[V] = 2 THEN
BEGIN BV:= B[V]; LE C[BV] = LRE ^ C(BV + 1) = LIM THEN
BEGIN V1:= B[BV]; V2:= B[BV + 1]; LE A[V1] = REAL ^ A[V2] = REAL THEN
BEGIN CPL1:= ISSUE; RE:= E[V1]; IM:= E[V2]; END
END
END
END
END;
END;
END;
END;

```



```

BOOLEAN PROCEDURE CPL2; CPL2:= LE CPL1(V1, RE1, IM1) THEN CPL1(V2, RE2, IM2) ELSE FALSE;

```

```

BOOLEAN PROCEDURE CPOP(S, RE, IM); INTEGER S; REAL RE, IM; LE PP = S THEN
BEGIN CPOP:= TRUE; PUTCPL(RE, IM) END
ELSE CPOP:= FALSE;

```

```

PROCEDURE CSIN(Q, COUT); INTEGER Q, COUT; COMMENT CASE IN CLAUSE;

```

```

BEGIN INTEGER N, I, QO, QI;
BOOLEAN JMP;
N:= B(Z); EMP; I:= 0; JMP:= FALSE;
FOR I:= 1, I WHILE INP(COM) DO
BEGIN I:= I + 1; QO:= QI; LE I = N THEN G ELSE F; MUNIT(QI), {UNIT IN CASE CLAUSE};
LE QO = G ^ QI = F THEN JMP:= TRUE
END;
LE JMP THEN Q:= F; COUT:= LE JMP ^ N > 1 ^ N & I THEN F ELSE G
END;

```

```

PROCEDURE CTX(P1); INTEGER P1; COMMENT COPY TO X;

```

```

BEGIN INTEGER K, XI;
K:= 0; P1:= P1 + 1; XI:= X(P1);
FOR I:= 1 WHILE K > 0 ^ (XI ≠ COM ^ XI ≠ GOON ^ XI ≠ CL) DO
BEGIN LE XI = OPEN THEN K:= K + 1 ELSE LE XI = CL THEN K:= K - 1; PX:= PX + 1; X(PX):= XI;
P1:= P1 + 1; XI:= X(P1)
END
END;

```

```

END;

```

```

BOOLEAN PROCEDURE DCL(VIC, T, Q); VALUE VIC, Q; INTEGER VIC, T, Q; SOMMENI DECLAREN;

```

```

BEGIN INTEGER L, J, V, POLD, VSH, V1, T1, NEXT;
BOOLEAN POS;

```

```

BOOLEAN PROCEDURE O(S); VALUE S; INTEGER S; LE INP(S) THEN

```

```

BEGIN O:= TRUE; LE Q = G THEN
BEGIN T:= ASK; A(T):= S; ENQ
END
ELSE O:= FALSE; POLD:= P; LE INP(OPEN) THEN P:= MATCH; NEXT:= X(P);
POS:= LE NEXT > 91 ^ NEXT < 119 THEN TRUE ELSE IN(IND, L); P:= POLD; LE POS THEN
BEGIN DCL:= TRUE;
LE O(INT) THEN ELSE LE O(REAL) THEN ELSE LE O(BOOL) THEN ELSE LE O(CHAR) THEN ELSE
LE O(VOID) THEN ELSE LE O(STRING) THEN
BEGIN LE O = G THEN
BEGIN A(T):= ROW; B(T):= V; ASK(N); C(T):= J; A(V):= 0; B(V):= V + 2; C(V):= 1;
E(V):= 1; V1:= V + 1; A(V1):= 1; B(V1):= 0; C(V1):= 1; A(V + 2):= CHAR
END
END

```

```

END;

```

```

ELSE LE INP(CMPL) THEN

```

```

BEGIN LE O = G THEN
BEGIN PUTCPL(0, 0); T:= Z END
END

```

```

ELSE LE STOR(VIC, T, Q) THEN ELSE LE RFOR(VIC, T, Q) THEN ELSE LE RWRD(VIC, T, Q) THEN
ELSE LE PROR(VIC, T, Q) THEN ELSE LE UNDR(VIC, T, Q) THEN ELSE LE IN(IND, L)
THEN (LE SRCH(L, J) THEN A(VIND(J)) = NEW ELSE FALSE) ELSE FALSE THEN
BEGIN LE O = G THEN
BEGIN V:= VIND(J); VSH:= B(V); LE VSH < 0 THEN
BEGIN POLD:= P; P:= C(V); B(V):= V1; ASK; M(DCL(ATL, T1, Q), {NEWMODE DCL});
P:= POLD; COPY(T1, V1); FREE(T1); T:= V1; B(V):= - 1
END
ELSE T:= VSH
END

```

```

END;
ELSE DCL:= FALSE
END
ELSE DCL:= FALSE
END;

PROCEDURE DCOL(APST); VALUE APST; INTEGER APST; COMMENT DECOLLATERATION)
BEGIN
  INTEGER VCOL, N, I, FECL, ZI, VI, V2, BAPS, STRV, MSTR, MZ;
  VCOL:= Z; N:= C(VCOL); FECL:= B(VCOL); LE A(APST); LE A(APST); STR I;
  BEGIN M(CIAPST) = N, {COLL.CL, TO STRUCT}; ZI:= ASKN(N + 1); BAPS:= BIAPST; COPY(APST, Z);
  ZI:= Z + 1; COPY(BAPS, ZI, N); BIZI:= ZI; FREE(VCOL);
  EOB I:= 1 SIER 1 UNILL N QQ
  BEGIN ZI:= FECL * 1 + 1; STRV:= B(BAPS * 1 + 1); MSTR:= A(STRV); MZ:= A(ZI);
  LE MZ = INT ^ MSTR = REAL ^ MZ = CHAR ^ MSTR * ROW I; THEN M(COER(STRV)); BIZI - 1 + 1 := Z
  {WID, IN COLL} ELSE LE MZ = COLL I; THEN DCOL(STRV); BIZI - 1 + 1 := Z
END;
Z:= ZI - 1
END
ELSE
BEGIN M(AIAPST) = ROW, {COLL.CL NO ROW OR STRUCT}; Z:= ASKN(N + 3); AIZI:= ROW;
  BIZI:= VI:= Z + 1; C(ZI):= 1; M(LE A(BIAPST)) = 0 I; THEN I; ELSE A(BIAPST) + 1 := 1;
  {COLL TO ROW, BOUND}; A(V1):= 1; B(V1):= V1 + 2; C(V1):= 1; E(V1):= N; V2:= V1 + 1;
  A(V2):= 1; B(V2):= N; C(V2):= 1; COPYN(FECL, V2 + 1, N); FREE(VCOL); FREE(FECL, N)
END;
RMCL(APST)
END;

BOOLEAN PROCEDURE DEN(Q); VALUE Q; INTEGER Q; COMMENT DENOTATION)
BEGIN
  INTEGER L;
  BOOLEAN SUC;
  SUC:= I; THEN LE IN(INT, L) I; THEN
  BEGIN LE Q = G I; THEN PUTINT(MISC(L + 1)) END
  ELSE LE IN(REAL, L) I; THEN
  BEGIN LE Q = G I; THEN PUTREAL(RNR(L)) END
  ELSE LE INP(TRUE) I; THEN
  BEGIN LE Q = G I; THEN PUTBOOL(TRUE) END
  ELSE LE INP(FALSE) I; THEN
  BEGIN LE Q = G I; THEN PUTBOOL(FALSE) END
  ELSE LE IN(CHAR, L) I; THEN
  BEGIN LE Q = G I; THEN PUTCHAR(MISC(L + 1)) END
  ELSE LE IN(STRING, L) I; THEN
  BEGIN LE Q = G I; THEN PUTSTRING(L) END
  ELSE LE IN(BITS, L) I; THEN
  BEGIN LE Q = G I; THEN PUTBITS(MISC(L)) END
  ELSE SUC:= FALSE; DEN:= SUC; LE Q = G ^ SUC I; THEN MORP:= FALSE
END;

INTEGER PROCEDURE DEV(V); VALUE V; INTEGER V; COMMENT DEVELOP)
BEGIN
  INTEGER N, I, VI, VD, N1, VN, NN, IN;
  BOOLEAN U;
  NN:= N; C(V); U:= FALSE;
  EOB I:= 1 SIER 1 UNILL N QQ
  BEGIN VI:= V + 1 - 1; LE A(VI) = UN I; THEN
  BEGIN VD:= B(VI); U:= I; THEN; N1:= C(VD); NN:= NN + N1 - 1 END
END;
LE U I; THEN
BEGIN DEV:= VN:= ASKN(NN); IN:= 0)

```

```

COR I:= 1 SIZE 1 UNTIL N DO
  BEGIN V:= V + 1 - 1; LE A(V) = UN THEN
    BEGIN VD:= B(V); N1:= C(VD);
      FOR J:= 1 SIZE 1 UNTIL N1 DO
        BEGIN IN:= IN + 1; B(VN + IN - 1); B(VD + J - 1); END
      END
    END
  ELSE
    BEGIN IN:= IN + 1; B(VN + IN - 1); B(V) END
  END;
PREEN(V, N)
END;
ELSE DEV:= V
END;

```

```

PROCEDURE DM(V); VALUE V; INTEGER V; COMMENT DISMISS;
BEGIN INTEGER AV, BV, CV, N, I, NV, BBV, ABV;
  AV:= A(V); BV:= B(V); CV:= C(V); LE AV = REF THEN RX(BV) ELSE LE AV = PROC THEN
    BEGIN ADV:= A(BV); BBV:= B(BV); NV:= LE BBV > 0 THEN BBV ELSE LE BBV = - 10 THEN 1 ELSE 2;
      FOR I:= 0 SIZE 1 UNTIL NV DO RS(BV + 1 - I); PREEN(BV, NV + 2)
    END
  ELSE LE AV = STR THEN
    BEGIN FOR I:= 1 SIZE 1 UNTIL CV DO RX(BV + 1 - I); PREEN(BV, CV) END
  ELSE LE AV = ROW THEN
    BEGIN FREEN(BV, CV + 1); RXN(B(BV), E(BV)) END;
  ELSE LE AV = 0 THEN
    BEGIN IF CV > 0 THEN RX(CV); N:= C(BV);
      BEGIN FOR I:= 1 SIZE 1 UNTIL N DO RS(B(BV + 1 - I)); FREEN(BV, N)
    END
  END;
END;

```

```

BOOLEAN PROCEDURE DPRO(Q); INTEGER Q; COMMENT DEPROCEED;
LE LE A(Z) = PROC THEN B(Z) = 0 ^ C(B(Z)) > 0 ELSE FALSE THEN
  BEGIN INTEGER POLD, V;
    POLD:= P; V:= Z; P:= C(B(Z)); M(CAST(Q), {DEPROCT}); RX(V); P:= POLD; DPRO:= ISSUE
  END
ELSE DPRO:= FALSE;

```

```

BOOLEAN PROCEDURE DREF; COMMENT DEREFERENCING;
LE A(Z) = REF THEN
  BEGIN INTEGER V;
    V:= Z; Z:= NINST(B(V)); RX(V); DREF:= ISSUE.
  END
ELSE DREF:= FALSE;

```

```

INTEGER PROCEDURE DTC; COMMENT DECLARER, TAG COUNT;
BEGIN INTEGER N, POLD;
  POLD:= P; LE LE DCL(FML, DUM, F) THEN IN(TAG, DUM) ELSE FALSE THEN
    BEGIN N:= 1;
      FOR I:= 1 WHILE LE LE INP(COM) THEN DCL(FML, DUM, F) ^ ISSUE ELSE FALSE THEN IN(TAG, DUM)
        ELSE FALSE DO N:= N + 1
    END
  ELSE N:= 0; DTC:= N; P:= POLD
END;

```

```

PROCEDURE EMP; RX(Z);

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1005

```

BOOLEAN PROCEDURE ENCL(Q); INTEGER Q; COMMENT ENCLOSED CLAUSE;
BEGIN
  INTEGER JOLD, POLD;
  ENCL:= FALSE; IF Q # G THEN
    BEGIN IF INP(OPEN) THEN
      BEGIN ENCL:= TRUE; P:= MATCH END
    END
  ELSE IF INP(OPEN) THEN
    BEGIN POLD:= P; ADMC(JOLD); IF SCL(Q) THEN
      BEGIN ENCL:= TRUE; IF INP(CL) THEN ELSE IF INP(VERT) THEN
        BEGIN IF Q = G THEN
          BEGIN CHCL(Q); M(INP(CL), {CHOOSER CL}); END
        ELSE P:= MATCH
      END
    ELSE IF Q = G THEN CLC(Q) ELSE P:= MATCH; ADMC(JOLD)
    END
  ELSE P:= POLD - 1
  END
END;

```

```

PROCEDURE ER(S); STRING S;
BEGIN NLCR; PRINTTEXT({ERROR I}); PRINTTEXT(S); SPACE(3); PRINTTEXT({ERROR AT POSITION});
  ABSFIXT(6, 0, P); NLCR; SODQ START
END;

```

```

PROCEDURE FILL;
BEGIN I:= I + 1; ARR(I):= MM END;

```

```

BOOLEAN PROCEDURE FORM(Q); INTEGER Q; COMMENT FORMULA;
BEGIN
  INTEGER W, VDEX, ODEX, PRIQ;
  INTEGER ABBY OPST, PRST, VALST(INFST);

```

```

  BOOLEAN PROCEDURE OPND; COMMENT OPERAND;
  BEGIN
    INTEGER OPMD, PMD, W, U, J;
    REAL R;

```

```

    IF OP(10, 10, J, DUM) THEN
      BEGIN
        N(OPND, $NO OPERAND AFTER MON, OPERATOR); OPND:= TRUE; IF Q = G THEN
          BEGIN ERR I:= I, I WHILE IF COER(FIRM, PMD, Q) THEN FALSE ELSE FURT(LINDI), := 10,
            J) DO

```

```

          BEGIN W:= B(VINDI); PMD:= B(W + 2) END;
          M(U > 0, $FILL.MON, OPTOR); P:= C(W); U:= Z; OPMD:= A(U);
          IF PP > 0 THEN OPT(1, PP, U, DUM) ELSE
            BEGIN P:= -P; IF OPMD = INT THEN

```

```

              BEGIN I:= B(U);
                IF IOP4(ABSV, ABS(I), MIN, -, 1, PLUS, 1, SGN, SIGN(I)) THEN ELSE IF
                  BOP(ODD, I # 1, 2 * 2) THEN ELSE
                    BEGIN M(PP = REPR, $FILL.MON, INT OP); PUTCHAR(I) .END

```

```

              END
            ELSE IF OPMD = REAL THEN

```

```

              BEGIN R:= E(U); I:= R;
                IF ROP4(ABSV, ABS(R), MIN, -, R, PLUS, R, 0, 0) THEN ELSE M(IOP4(SGN,
                  SIGN(R), ENT, ENTIER(R), RND, 1, 0, 0), $FILL.MON, REAL OP);

```

```

              END
            ELSE IF OPMD = ROW THEN
              BEGIN I:= B(U) + 1; IF IOP(UPB, B(I)) THEN ELSE M(IOP(LWB, A(I)),
                $FILL.MON, ROW OP);
            END

```

Sp. ? Smaller?

```

ELSE LE OPMD = CHAR THEN M(IOP(ABS, B(U)),
  $ILL.MON.CHAROP) ELSE LE OPMD = BOOL THEN M(IOP(ABS, B(U)),
  $ILL.MON.BOOL.OP) ELSE
  BEGIN M(CPLI(U, REI, IMI), $ILL.MON.OP);
  LE ROP4(RE, REI, IM, IMI, ABSV, SORT(REI * REI + IMI * IMI), 0, 0) THEN
  ELSE LE CPOP(PLUS, REI, IMI) THEN ELSE LE CPOP(MIN, -REI, -IMI)
  THEN ELSE LE CPOP(CONJ, REI, -IMI) THEN ELSE
  BEGIN M(PP = ARG * (REI * 0 + IMI * 0), $ILL.MON.CPL.OP);
  PUTREAL(ARGC(REI, IMI))
END
END;
  RK(U)
END
END

```

```

END
ELSE OPND := SEC(Q)
END OPND;
END
END

```

```

BOOLEAN PROCEDURE OP(PRL, PRH, OPTR, PRIO); VALUE PRL, PRH; INTEGER PRL, PRH, OPTR, PRIO;
COMMENT OPERATOR;
IF Q = G THEN
  BEGIN INTEGER JJ, POLD, LJ

```

```

  BOOLEAN PROCEDURE DIPP;
  BEGIN PRIO := B(BVIND(JJ)); DIFF := PRIO < PRL + PRIO > PRH END;
  JJ := J; POLD := P; OP := FALSE; LE IN(IND, L) THEN
  BEGIN
    FOR I := 1 WHILE LE JJ > 0 THEN (LE LE L IND(JJ) THEN A(VIND(JJ)) = PROC ELSE
    FALSE THEN DIFF ELSE TRUE) ELSE FALSE DO JJ := JJ - 1; LE JJ > 0 THEN
    BEGIN OP := TRUE; OPTR := JJ END
    ELSE P := POLD
  END
END

```

```

END
ELSE OP := IN(IND, DUM);
PROCEDURE OPRT(N, PRTN, V, W); VALUE N, PRTN, V, W; INTEGER N, PRTN, V, W;
COMMENT OPERATOR ROUTINE;
INTEGER PX0, POLD, I, JOLD;
POLD := P; ADMO(JOLD); PX0 := PX;
FOR I := 1 STEP 1 UNTIL N DO
  BEGIN CTX(PRTN); WTX(EQ); WTX(OPROUT); WTX(LE I = 1 THEN V ELSE W);
  WTX(LE I = N THEN GOON ELSE COM)
END;
P := PX0 + 1; PX := PX0; M(IYD(N), $OPERAND); P := PRTN + 1; M(CAST(Q), $OP DECL);
P := POLD; ADMC(JOLD)
END;

```

```

PROCEDURE DYOP(V, O); INTEGER V, O; COMMENT DYADIC OPERATION;
BEGIN INTEGER PRY, W, MP1, MP2, L, J, ABY1, RR, X, K;
  BEGIN Y, DD, REP, IMP, REP1;
  BOOLEAN TRY;
  V2 := VLST(V); V := V - 1; V1 := VLST(V); J := OPST(O); O := 0 - 1; TRY := TRUE; L := LIND(J);
  FOR I := 1 WHILE TRY DO
    BEGIN W := BVIND(J); MP1 := B(W + 2); MP2 := B(W + 3); Z := V1; TRY := B(W);
    LE COER(FIRM, MP1, O) THEN
      BEGIN V1 := Z; Z := V2; LE COER(FIRM, MP2, O) THEN
        BEGIN V2 := Z; TRY := FALSE END
      END
    END
  END

```

```

ELSE TRY:= FURT(L, PRY, J)
END
ELSE TRY:= FURT(L, PRY, J)
END;
M(J) > 0, $ILL.DYADIC OPERATOR$); PPI:= C(W); LE PP > 0 THEN OPRT(2, PP, V1, V2) ELSE
BEGIN PPI:= -PP; M1:= A(V1); M2:= A(V2); BV1:= B(V1); BV2:= B(V2); EV1:= E(V1);
EV2:= E(V2); LE PP < PLUS THEN
BEGIN LE PP < EQ THEN
BEGIN LE PP < OR THEN
BEGIN LE PP = PLUS THEN
BEGIN M(M2 = REF, $+*?); M1STCH(V1, BV2), $ILL.OP.$); SUMS; SPS(Z, BV2);
EMP; Z:= NRT(BV2)
END
ELSE
BEGIN M(M1 = REF, $LEFT OP. NO NAME$); ABV1:= A(BV1); LE ABV1 = INT THEN
BEGIN M(M2 = INT, $RIGHT OP. NO INT$); X:= B(BV1);
LE IOP4(PSAB, X + BV2, MAB, X * BV2, TMAB, X * BV2, OVAB, X / BV2)
THEN ELSE
BEGIN RR:= X - X / BV2 * BV2;
M(IOP(MDAB, LE RR < 0 THEN RR + ABS(BV2) ELSE RR), $IL ABOPT$)
END;
SPS(Z, BV1); EMP; Z:= NRT(BV1)
END
ELSE LE ABV1 = REAL THEN
BEGIN Y:= E(BV1); LE M2 = INT THEN EV2:= BV2;
M(ROP4(PSAB, Y + EV2, MAB, Y * EV2, TMAB, Y * EV2, DVAB, Y / EV2);
$ILL.AB OPT$); SPS(Z, BV1); EMP; Z:= NRT(BV1)
END
ELSE
BEGIN M(STCH(BV1, V2) ^ PP = PSAB, $ILL.+*?); SUMS; SPS(Z, BV1); EMP;
Z:= NRT(BV1)
END
END
END
END
ELSE BOP2(OR, BV1 = 1 ^ BV2 = 1, AND, BV1 = 1 ^ BV2 = 1)
END
ELSE LE M1 = INT ^ M2 = INT THEN
BEGIN LE BOP4(EQ, BV1 = BV2, NE, BV1 # BV2, LESS, BV1 < BV2, GT, BV1 > BV2)
THEN ELSE BOP2(LE, BV1 < BV2, GE, BV1 > BV2)
END
ELSE LE RL2 THEN
BEGIN LE BOP4(EQ, EV1 = EV2, NE, EV1 # EV2, LESS, EV1 < EV2, GT, EV1 > EV2)
THEN ELSE BOP2(LE, EV1 < EV2, GE, EV1 > EV2)
END
ELSE LE $CRL THEN
BEGIN LE BOP4(EQ, LE UNEQ THEN LEN1 + LEN2 = 0 ELSE LEN1 = LEN2, NE, LE UNEQ
THEN LEN1 + LEN2 > 0 ELSE LEN1 # LEN2, LESS, LSTR(FE1, LEN1, FE2,
LEN2), GT, LSTR(FE2, LEN2, FE1, LEN1)) THEN ELSE BOP2(LE, $LSTR(FE2, LEN2,
FE1, LEN1), GE, $LSTR(FE1, LEN1, FE2, LEN2))
END
ELSE LE M1 = BOOL ^ M2 = BOOL THEN BOP2(EQ, BV1 = BV2, NE, BV1 # BV2) ELSE LE
CPL2 THEN BOP2(EQ, RE1 = RE2 ^ IM1 = IM2, NE, RE1 # RE2 ^ IM1 # IM2)
END
PP<PLUS
ELSE LE PP < PLIT THEN
BEGIN LE M1 = INT ^ M2 = INT THEN
BEGIN LE IOP4(PLUS, BV1 + BV2, MIN, BV1 - BV2, TIMES, BV1 * BV2, OVER, BV1 / BV2)
THEN ELSE LE IOP4(DIV, BV1 / BV2) THEN ELSE LE IOP4(UP, BV1 + BV2) THEN
ELSE
BEGIN RR:= BV1 = BV1 / BV2 * BV2;
PUTINT(LE RR < 0 THEN RR + ABS(BV2) ELSE RR)

```

```

END
ELSE JE STCH(V1, V2) THEN
  BEGIN M(PP = PLUS, $FILL CHAR OR STRING OP); SUMS END
ELSE JE PP < UP THEN
  BEGIN LE RL2 THEN M(ROP4PLUS, EV1 + EV2, MIN, EV1 - EV2, TMS, EV1 * EV2, DIV,
    EV1 / EV2), $REAL OVER OR MOD; ELSE JE CPL2 THEN
    BEGIN JE CPOP(PLUS, RE1 + RE2, IM1 + IM2) THEN ELSE JE CPOP(MIN, RE1 - RE2,
      IM1 - IM2) THEN ELSE JE CPOP(TMS, RE1 * RE2 - IM1 * IM2, RE1 * IM2 +
        IM1 * RE2) THEN ELSE
        BEGIN DD:= RE2 * RE2 + IM2 * IM2;
          M(CPOP(DIV, (RE1 * RE2 + IM1 * IM2) / DD, (- RE1 * IM2 + IM1 * RE2)
            / DD), $FILL.COMPLEX OP.);
        END
      END
    END
  END
ELSE
  BEGIN M(M2 = INT, $EXP NO INT); JE M1 = REAL THEN PUTREAL(EV1 + BV2) ELSE
    BEGIN M(CPL1(V1, RE1, IM1), $FILL.UP); REP:= 1; IMP:= 0; KI= ABS(BV2);
      FOR I:= 1 SIZE 1 UNTIL K DO
        BEGIN REP:= REP * RE1 - IMP * IM1; IMP:= IMP * RE1 + REP * IM1;
          REP:= REP;
        END;
      JE BV2 < 0 THEN
        BEGIN DD:= REP * REP + IMP * IMP; PUTCPL(REP / DD, - IMP / DD) END
      ELSE PUTCPL(REP, IMP)
    END
  END
  END UP
END
ELSE JE PP = PLIT THEN EV1:= BV1 ELSE M(M1 = REAL, $PLIT);
  BEGIN JE M1 = INT THEN EV1:= BV1 ELSE M(M2 = REAL, $PLIT); PUTCPL(EV1, EV2);
    JE M2 = INT THEN EV2:= BV2 ELSE M(M2 = REAL, $PLIT);
  END PLIT
ELSE
  BEGIN M(PP = LWB * PP = UPB) * BV1 2.1 ^ BV1 5 C(V2) ^ M1 = INT ^ M2 = ROW,
    $FILL.DYADIC OP.); PUTINT(JE PP = LWB THEN A(BV2 + BV1) ELSE B(BV2 + BV1))
  END LWB OR UPB;
  RX(V1) RX(V2)
END;
VST(V):= Z
END DYOP;

PROCEDURE STVL;
  BEGIN VDEX:= VDEX + 1; M(VDEX $ NFST, $VALUE STACK EXC.); VLAST(VDEX):= Z END;

PROCEDURE STOR;
  BEGIN ODEX:= ODEX + 1; M(ODEX $ NFST, $OP.STACK EXC.); OPST(ODEX):= W; PRST(ODEX):= PRIO END;

COMMENT FORMULA STARTS HERE;
JE JE LOOK(OP(10, 10, DUM, DUM)) THEN ISSUE ELSE LK2(SEC(F), OP(0, 9, DUM, DUM)) THEN
  BEGIN FORM:= ISSUE; M(OPND, $NO OPERAND); JE O = G THEN
    BEGIN VDEX:= ODEX + 1; PRST(0):= 0; STVL; MORF:= ISSUE;
      FOR I:= 1 WHILE OP(0, 9, W, PRIO) DO
        BEGIN FOR I:= 1 WHILE PRIO $ PRST(ODEX) DO DYOP(VDEX, ODEX); STOR; M(OPND,
          $NO OPERAND AFTER OPERATOR); STVL
        END;
      FOR I:= 1 WHILE VDEX > 1 DO DYOP(VDEX, ODEX)
    END
  END
ELSE

```

```

FOR I:= 1 WHILE ANDR(OP(O, 9, DUM, DUM), OPND) DO:
  END
  ELSE FORM:= FALSE
  END FORM;

PROCEDURE FREE(V); VALUE V; INTEGER V;
BEGIN D(V):= FS; FS:= V END;

PROCEDURE FREE(V, N); VALUE V, N; INTEGER V, N;
  BEGIN INTEGER W;
    W:= V + N - 1;
    FOR V:= V STEP 1 UNTIL W DO FREE(V)
  END;

BOOLEAN PROCEDURE FURT(L, PR, J); VALUE L, PR; INTEGER L, PR, J; COMMENT FURTHER SEARCH FOR OPERATOR;
  BEGIN J:= J - 1;
    FOR I:= 1 WHILE LE J > 0 THEN (LE L = LIND(J) ^ A(VIND(J)) = PROC THEN B(S(VIND(J))) # PR ELSE
      TRUE) ELSE FALSE DO J:= J - 1; FURT:= J > 0
  END;

BOOLEAN PROCEDURE GEN(O); VALUE O; INTEGER O; COMMENT GENERATOR;
  LE LE IN(LOC) THEN TRUE ELSE INP(HEAP) THEN
  BEGIN INTEGER T;
    GEN:= TRUE; M(DCL(ATL, T, O), #NO DCL IN GEN); LE O = G THEN
    BEGIN Z:= NRT(T); MORF:= FALSE END
  END
  ELSE GEN:= FALSE;

PROCEDURE GOUT(V); VALUE V; INTEGER V; COMMENT GENERAL OUTPUT;
  BEGIN INTEGER AZ, BZ, CZ, O, N, FE, I;
    O:= G; Z:= V;
    FOR I:= 1 WHILE DREF V DPRO(O) DO, M(O = G, #IMPL, RESTR, OUTPUT JUMP); AZ:= A(Z); BZ:= B(Z);
    CZ:= C(Z); LE AZ = STR THEN
    BEGIN FOR I:= 1 STEP 1 UNTIL CZ DO GOUT(B(BZ - 1 + I)); END
    ELSE LE AZ = ROW THEN
    BEGIN N:= E(BZ); FE:= B(BZ);
      FOR I:= 1 STEP 1 UNTIL N DO GOUT(FE - 1 + I)
    END
    ELSE LE AZ = UN THEN GOUT(CZ) ELSE LE AZ = COLL THEN
    BEGIN FOR I:= 1 STEP 1 UNTIL CZ DO GOUT(BZ - 1 + I) END
    ELSE LE AZ = INT ^ AZ = BOOL THEN PRINT(BZ) ELSE LE AZ = REAL THEN PRINT(E(Z)) ELSE LE AZ =
    CHAR THEN PRSYN(BZ) ELSE LE AZ = PROC THEN
    BEGIN PP:= C(BZ);
      LE PP = - 999 THEN NEWPAGE ELSE LE PP = - 998 THEN NLCR ELSE LE PP = - 997 THEN SPACE(1)
      ELSE ER(#KILL, OUTP.#)
    END
    ELSE ER(#KILL, OUTPUT#)
  END;

BOOLEAN PROCEDURE HIP(O); INTEGER O;
HIP:= LE SKIP(O) THEN TRUE ELSE LE JUMP(O) THEN TRUE ELSE LE NIHL(O) THEN TRUE ELSE VAC(O);

BOOLEAN PROCEDURE IDP(O); VALUE O; INTEGER O; COMMENT MODE IDENTIFIER;
  BEGIN INTEGER L, J, POLD;

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POLD:= P; LE IN(TAG, L) THEN
BEGIN IDFI:= TRUE; LE Q = G THEN
BEGIN LE SRCH(L, J) THEN
BEGIN Z:= VIND(J); LE Z < 0 THEN
BEGIN IDFI:= FALSE; P:= POLD END.
ELSE
BEGIN DIZ:= DIZ + 1; MORP:= TRUE END.
END
ELSE
BEGIN NLCR; NLCR;
FOR J:= 1 STEP 1 UNTIL MISC(L) DO ABSFIX(4, 0, MISC(L + J)); ER(
UNKNOWN IDENTIFIER)
END
END
END;
ELSE IDFI:= FALSE
END;

BOOLEAN PROCEDURE IN(K, L); VALUE KI; INTEGER K, L; COMMENT READS NEXT ELEMENT OF X, IF THIS HAS MARK K;
BEGIN
INTEGER XX;
XX:= X(P); LE XX = 10000 = K THEN
BEGIN IN:= TRUE; P:= P + 1; L:= XX - K * 10000 END
ELSE IN:= FALSE
END;

BOOLEAN PROCEDURE INP(S); VALUE S; INTEGER S; COMMENT READS NEXT ELEMENT OF X, IF THIS IS S;
BEGIN
INTEGER L;
LE X(P) = S THEN
BEGIN INP:= TRUE; P:= P + 1 END
ELSE INP:= LE S = EQ THEN ANDP(IN(IND, L), L = 10) ELSE FALSE
END;

PROCEDURE INVS(J1, J2); VALUE J1, J2; INTEGER J1, J2; COMMENT MAKES PART OF THE LINK TABLE INVISIBLE;
BEGIN
INTEGER JJ;
FOR JJ:= J1 + 1 STEP 1 UNTIL J2 DO LE LIND(JJ) < 1000 THEN LIND(JJ):= LIND(JJ) + 1000
END;

BOOLEAN PROCEDURE IOP(S, Y); VALUE S; INTEGER S, Y; LE PP = S THEN
BEGIN IOP:= TRUE; PUTINT(Y) END
ELSE IOP:= FALSE;

BOOLEAN PROCEDURE IOP4(S1, I1, S2, I2, S3, I3, S4, I4); INTEGER S1, I1, S2, I2, S3, I3, S4, I4;
BEGIN IOP4:= TRUE;
LE PP = S1 THEN PUTINT(I1) ELSE LE PP = S2 THEN PUTINT(I2) ELSE LE PP = S3 THEN PUTINT(I3)
ELSE LE PP = S4 THEN PUTINT(I4) ELSE IOP4:= FALSE.
END;

BOOLEAN PROCEDURE IRN(Q); VALUE Q; INTEGER Q; IRDN:= LE IRON(Q) THEN TRUE ELSE VDN(Q);
COMMENT IDENTIFIER DECLARATION;

PROCEDURE IRL(Q, IT); VALUE IT; INTEGER Q, IT; COMMENT IDENTITY RELATION TAIL;
BEGIN
BOOLEAN ISI, IDNT;
INTEGER I, J, K, U1, U2;

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INTEGER ARRAY N, U, V(1:2);
IF Q # G THEN M(TERT(F), IDENT.REL.) ELSE
BEGIN
  MORF:= EALSE;
  FOR I:= 1 WHILE DPRO(Q) DO; U1:= V(1); Z1:= I1:= IS; M(TERT(Q), {ID,IR,M,S.});
  FOR I:= 1 WHILE OPRO(Q) DO; U2:= V(2); Z2:= M(AU1) = REF ^ A(Z) = REF,
  {NO NAMES IN ID, REL.}; JE BIU1 = BIZ THEN IDNT:= ISBE ELSE
  BEGIN
    FOR I:= 1, 2 DO
      BEGIN V(1):= U(1); N(1):= 1;
        FOR I:= 1 WHILE A(V(1)) = REF DO
          BEGIN N(1):= N(1) + 1; V(1):= B(V(1)) END
        END;
        JE N(1) > N(2) THEN 1 ELSE 2; K:= 3 - J1
        FOR I:= N(1) - N(K) STEP - 1 UNTIL 1 DO U(J1):= B(U(J1)); IDNT:= B(U(J1)) = B(U(K));
        END;
        RX(U1); RX(U2); Z1:= ASK; A(Z1):= BOOL; B(Z1):= JE IDNT = IS1 THEN 1 ELSE 0
      END;
    END;
  END;
  BOOLEAN PROCEDURE IYDF(G, V); VALUE G, V; INTEGER O, V; COMMENT IDENTITY DEFINITION;
  BEGIN
    INTEGER L;
    JE ANDP(IN(TAG, L), INP(EQ)) THEN
      BEGIN IYDF:= ISBE; JE G = G THEN
        BEGIN JE INP(OPROUT) THEN
          BEGIN Z:= X(P); P:= P + 1 END
          ELSE M(UNIT(Q), {NO UNIT IN IDENTITY DEF}); M(COER(STRONG, V, G),
            {MODE IN IDENTITY DEF}); POS(L + 10000, Z, G)
          END
          ELSE JE INP(OPROUT) THEN P:= P + 1 ELSE M(UNIT(F), {IDENTITY DEF});
        END
      ELSE IYDF:= EALSE
    END;
  END;
  BOOLEAN PROCEDURE IYDN(Q); VALUE Q; INTEGER O; COMMENT IDENTITY DECLARATION;
  BEGIN
    BOOLEAN SUC;
    INTEGER JOLD;
    IYDN:= SUC:= JE LE Q = F THEN ISBE ELSE LOOK(IYDN(F)) THEN (JE LIST(PIDN(Q)) THEN ISBE ELSE
      LIST(MIDN(Q)) ELSE EALSE); JE SUC THEN VISB(JOLD, J)
    END;
  END;
  BOOLEAN PROCEDURE JUMP(O); INTEGER O;
  BEGIN
    INTEGER L, J;
    JE INP(COTO) THEN; JE IN(TAG, L) THEN
      BEGIN JUMP:= ISBE; JE Q = G THEN
        BEGIN Q:= F; M(SRCH(L, J), {LABEL NOT FOUND}); PLAB:= VIND(J); M(PLAB > 0, {LABEL});
        END
      END
    END
    ELSE JUMP:= EALSE
  END;
  BOOLEAN PROCEDURE LAB(O); VALUE O; INTEGER O; COMMENT DEFINING OCC. OF LABELS;
  BEGIN
    INTEGER L;
    JE ANDP(IN(TAG, L), INP(COL)) THEN
      BEGIN LAB:= ISBE; JE Q = H THEN
        BEGIN J:= J + 1; LIND(J):= L; VIND(J):= P END.
      END
    END
  END;

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ELSE LAB1= FALSE
END;

BOOLEAN PROCEDURE LBUN(Q); INTEGER Q;
BEGIN IF SEQ(LAB(Q)) THEN ; LBUN:= UNIT(Q) END;

COMMENT LABELLED UNIT;

BOOLEAN PROCEDURE LIST(B); BOOLEAN B; IF B THEN
BEGIN FOR I:= 1 WHILE ANDP(INP(COM), B) DO ; LIST:= TRUE END
ELSE LIST:= FALSE;

BOOLEAN PROCEDURE LK2(A, B); BOOLEAN A, B;
BEGIN INTEGER POLD;
POLD:= P; LK2:= IF A THEN B ELSE FALSE; P:= POLD
END;

BOOLEAN PROCEDURE LOOK(A); BOOLEAN A;
BEGIN INTEGER POLD;
POLD:= P; LOOK:= A; P:= POLD
END;

BOOLEAN PROCEDURE LOOP(Q); INTEGER Q;
BEGIN INTEGER L, V, ST, ENDV, P1, Q1, Q0, JOLD, P2, CV, Z1;
BOOLEAN EX1, EXH, HALT, BND, FRST, JMP;

BOOLEAN PROCEDURE CONT;
BEGIN IF Q1 = G THEN CV:= B(V); Q0:= Q1; IF -FRST THEN P1:= P1;
EXH:= IF Q1 ≠ G THEN TRUE ELSE IF -BND → ST = 0 THEN FALSE ELSE SIGN(ST) * (CV - ENDV)
> 0; IF INP(WHL) THEN
BEGIN M(SCL(Q1), {AFTER WHILE}); JMP:= Q1 ≠ Q0; M(COER(MEEK, STBL, Q1),
→MODE AFTER WHILE); EXH:= EXH ∨ B(Z) = 0
END;
M(INP(DO), {DO MISSING}); IF FRST THEN
BEGIN CONT:= TRUE; IF Q1 = G ^ EXH THEN Q1:= P; FRST:= FALSE END
ELSE
BEGIN HALT:= Q1 ≠ G ∨ EXH; IF HALT THEN P:= P2; CONT:= -HALT END
END;

IF (IF SPY(FOR) THEN TRUE ELSE IF SPY(FROM) THEN TRUE ELSE IF SPY(BY) THEN TRUE ELSE IF SPY(TO)
THEN TRUE ELSE IF SPY(WHL) THEN TRUE ELSE SPY(DO)) THEN
BEGIN LOOP:= TRUE; EX1:= Q = G; IF EX1 THEN
BEGIN ADMO(JOLD); V:= ASK; A1(V):= INT END;
IF INP(FOR) THEN
BEGIN M(INTAG, L), {NO ID AFTER FOR}; POS(L, V, Q) END;
IF INP(FROM) THEN
BEGIN M(UNIT(Q), {UNIT AFTER FROM}); IF EX1 THEN B(V):= MEEKINT(Q) END
ELSE IF EX1 THEN B(V):= 1; IF INP(BY) THEN
BEGIN M(UNIT(Q), {UNIT AFTER BY}); IF EX1 THEN ST:= MEEKINT(Q) END
ELSE IF EX1 THEN ST:= 1; IF INP(TO) THEN
BEGIN M(UNIT(Q), {UNIT AFTER TO}); IF EX1 THEN
BEGIN ENDV:= MEEKINT(Q); BND:= TRUE END
END
ELSE BND:= FALSE; JMP:= FALSE; P1:= P; FRST:= TRUE; Q1:= Q;
FOR I:= 1 WHILE CONT DO
BEGIN Q0:= Q1; M(UNIT(Q1), {UNIT AFTER DO}); P2:= P; Z1:= Z; IF Q1 = G THEN
BEGIN B(V):= B(V) + ST END

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END;
LE JMP V Q1 * Q0 THEN Q1 = P; LE EX1 THEN ADMC(JOLD)
ELSE LOOP := FALSE
END;

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INTEGER PROCEDURE LSTC(B); BOOLEAN B; COMMENT COUNTS THE MEMBERS OF A LIST;

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BEGIN INTEGER K, POLD;
POLD := P; LE B THEN
BEGIN K := 1;
FOR I := 1 WHILE ANDP(INP(COM), B) DO K := K + 1; P := POLD; LSTC := K
END
ELSE LSTC := 0
END;

```

```

BOOLEAN PROCEDURE LSTR(F1, L1, F2, L2); VALUE F1, L1, F2, L2; INTEGER P1, L1, P2, L2;
COMMENT FOR COMPARISON OF STRINGS;
LSTR := LE UNEQ THEN B(F1 - 1 + 1) < B(F2 - 1 + 1) ELSE L1 < L2;

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PROCEDURE M(B, S); BOOLEAN B; STRING S; COMMENT MUST
LE B THEN ER(S);

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BOOLEAN PROCEDURE MADE(V, W); VALUE V, W; INTEGER V, W; COMMENT LOOK IF ASSUMPTION OF MODE EQ, WAS MADE;
BEGIN INTEGER I;
I := 0;
FOR I := 1 + 1 WHILE LE I ≤ NASP THEN ~(V = VASP[I] ^ W = WASP[I]) ^ V = VASP[I] ^ W = WASP[I];
ELSE FALSE DO; MADE := I ≤ NASP
END;

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```

PROCEDURE MAKE(V, W); VALUE V, W; INTEGER V, W; COMMENT MAKE ASSUMPTION; MODES EQ;
BEGIN NASP := NASP + 1; M(NASP ≤ NASP, $SPACE FOR MODE EQ, EXH.$); VASP(NASP) := V; WASP(NASP) := W; END;

```

```

INTEGER PROCEDURE MATCH; COMMENT SKIP TO MATCHING CLOSE SYMBOL;
BEGIN INTEGER I;
I := 1;
FOR I := 1 WHILE I > 0 DO LE INP(OL) THEN I := 1 - 1 ELSE LE INP(OPEN) THEN I := 1 + 1 ELSE P := P
+ 1; MATCH := P
END;

```

```

BOOLEAN PROCEDURE MOD(Q); VALUE Q; INTEGER Q; COMMENT MODE DEFINITION;
BEGIN INTEGER L, P1, V;
LE IN(IND, L) THEN
BEGIN MOD := ISBE; M(INP(EQ), $ NO IN MODE DEF, $); P1 := P; M(DCL(ATL, DUM, P),
$ NO ACT DCL IN MODE DEF $); LE Q = H THEN
BEGIN V := ASK; A(V) := NEW; B(V) := - 1; C(V) := P1; POS(L, V, G); END
END
ELSE MOD := FALSE
END;

```

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BOOLEAN PROCEDURE MDN(G); VALUE G; INTEGER G; COMMENT MODE DECLARATION;
LE INP(MODE) THEN
BEGIN MDN := ISBE; M(LIST(MDN(G)), $ NO MODE DEF $); END

```

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ELSE MONI:= EALSE;
INTEGER PROCEDURE MEKINT(Q); INTEGER Q; COMMENT MEK COERCION TO INTEGRAL;
BEGIN M(COER(MEK, STIN, Q), $ NO MEK INT UNIT WHERE REQUIRED $); MEKINT:= B(Z); EMP END;

BOOLEAN PROCEDURE MEQ(V, W); VALUE V, W; INTEGER V, W; COMMENT MEQ EQ. USED BY MEQ;
BEGIN INTEGER AV, AW, BV, BW, CV, CW, I, NV, NW;
  AV:= A(V); AW:= A(W); I:= V = W ^ AV = W ^ NV = I, NV, NW;
  BEGIN BV:= B(V); BW:= B(W); CV:= C(V); CW:= C(W);
    I:= AV = REF THEN MEQ:= MEQ(BV, BW) ELSE I:= AV = STR THEN
      BEGIN I:= MADE(V, W) THEN MEQ:= TRUE ELSE I:= CV = CW THEN
        FOR I:= 1 + 1 WHILE I = CV THEN EALSE ELSE I:= C(BV + I) ^ C(BW + I) THEN
          EALSE ELSE MEQ:= MEQ(BV + I), B(BW + I)) DO; MEQ:= I = CV
        END
      ELSE MEQ:= EALSE
    END
  ELSE I:= AV = ROW THEN MEQ:= I:= CV = CW THEN MEQ(BV, B(BW)) ELSE EALSE ELSE I:= AV =
    UN THEN
      BEGIN NV:= C(BV); NW:= C(BW); I:= 1;
        FOR I:= 1 + 1 WHILE I = NV THEN EALSE ELSE UNMB(BV + I, BW) DO; I:= NV THEN
          BEGIN I:= 1;
            FOR I:= 1 + 1 WHILE I = NW THEN EALSE ELSE UNMB(BW + I, BV) DO;
              MEQ:= I = NW
            END
          ELSE MEQ:= EALSE
        END
      ELSE I:= AV = PROC THEN
        BEGIN NV:= B(BV); I:= NV = 10 THEN NV:= 1 ELSE I:= NV < 0 THEN NV:= 2; NW:= B(BW);
          I:= NW = -10 THEN NV:= 1 ELSE I:= NW < 0 THEN NV:= 2; I:= NV = NW THEN
            BEGIN I:= 0;
              FOR I:= 1 + 1 WHILE I = NV + 1 WHILE I = NV + 2 THEN EALSE ELSE MEQ(BV + I), B(BW + I)) DO;
                MEQ:= I = NV + 2
              END
            ELSE MEQ:= EALSE
          END
        ELSE MEQ:= TRUE
        END
      ELSE MEQ:= EALSE
    END;

COMMENT MODE IDENTITY DECLARATION;

BOOLEAN PROCEDURE MIDN(Q); VALUE Q; INTEGER Q;
BEGIN INTEGER V, POLD;
  BOOLEAN SUC;
  POLD:= P; SUC:= EALSE; I:= DCL(PML, V, Q) THEN
    BEGIN SUC:= LIST(YDF(Q, V)); I:= Q = G THEN RS(V) END;
    MIDN:= SUC; I:= -SUC THEN P:= POLD
  END;

PROCEDURE MOVE(V, W); VALUE V, W; INTEGER V, W;

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BEGIN INTEGER AV, BV, CV, U, W1, N, FE2, I;
AV := A[V]; BV := B[V]; CV := C[V]; COPY(V, W); LE AV = REF THEN
BEGIN U := BV; D[U] := D[U] + 1 END
ELSE LE AV = STR THEN
BEGIN W1 := B[W]; ASKN(CV); COPY(BV, W1, CV);
FOR I := 1 STEP 1 UNTIL CV DO B[W1 + I - 1] := NINST(B[BV + I - 1]);
END
ELSE LE AV = ROW THEN
BEGIN B[W] := W1; ASKN(CV + 1); COPY(BV, W1, CV + 1); N := E[BV]; LE N = 0 THEN N := 1;
B[W1] := FE2; ASKN(N); MOVEN(B[BV], FE2, N)
END
ELSE LE AV = UN THEN
BEGIN N := C[BV]; B[W] := U; ASKN(N); C[U] := N; LE CV > 0 THEN C[W] := NINST(CV);
FOR I := 1 STEP 1 UNTIL N DO B[U + I + 1] := NINST(B[BV + I + 1]);
END
ELSE LE AV = PROC THEN
BEGIN N := B[BV]; B[W] := W1; ASKN(N + 2); COPY(BV, W1);
FOR I := 0 STEP 1 UNTIL N DO B[W1 + I + 1] := NINST(B[BV + I + 1]);
END
END;
END;

```

```

PROCEDURE MOVEN(V, W, N); VALUE V, W, N; INTEGER V, W, N;
BEGIN INTEGER I, M;
M := N - 1;
FOR I := 0 STEP 1 UNTIL M DO MOVE(V + I, W + I)
END;

```

```

BOOLEAN PROCEDURE MVN(O); VALUE O; INTEGER O; COMMENT MODE VARIABLE DECLARATION;
BEGIN INTEGER V, POLD;
BOOLEAN SUC;
POLD := P; SUC := FALSE; LE INP(HEAP) THEN I LE DCL(ATL, V, O) THEN
BEGIN SUC := LIST(VDF(O, V)); LE O = G THEN RS(V) END;
MYDN(SUC) LE -SUC THEN P := POLD
END;

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BOOLEAN PROCEDURE NIHL(O); VALUE O; INTEGER O; LE INP(NIL) THEN
BEGIN NIHL := TRUE; LE O = G THEN
BEGIN Z := ASK; A[Z] := REF; B[Z] := 0 END
ELSE NIHL := FALSE;

```

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INTEGER PROCEDURE NINST(V); VALUE V; INTEGER V; COMMENT NEW INSTANCE OF;
BEGIN INTEGER W;
NINST := W; ASK; MOVE(V, W)
END;

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INTEGER PROCEDURE NRT(V); VALUE V; INTEGER V; COMMENT NEW REFERENCE TO;
BEGIN INTEGER W;
NRT := W; ASK; A[W] := REF; B[W] := V; D[V] := D[V] + 1
END;

```

```

PROCEDURE OD(OP, M1, M2, MD, PRIO); INTEGER OP, M1, M2, MD, PRIO; COMMENT DYADIC OPERATOR;
BEGIN Z := ASKN(S); A[Z] := PROC; B[Z] := Z + 1; A[Z1] := 0; B[Z1] := PRIO; C[Z1] := - OP;
B[Z1 + 1] := MD; B[Z1 + 2] := M1; B[Z1 + 3] := M2; LE OP = OLD THEN
BEGIN L := L + 1; OLD := OP END;

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END;
POS(L, Z, G)

PROCEDURE OM(OP, M1, M2); INTEGER OP, M1, M2; COMMENT MONADIC OPERATOR;
BEGIN Z := ASK(N(4)); A(Z) := PROC B(Z) := Z + 1; A(Z) := 0; B(Z) := -10; C(Z) := - OP;
  B(Z + 1) := M1; B(Z + 2) := M2; LE OP + OOLD THEN
  BEGIN L := L + 1; OOLD := OP END;
  POS(L, Z, G)
END;

BOOLEAN PROCEDURE OPON(Q); VALUE Q; INTEGER Q; COMMENT OPERATION DECL. ;
LE INP(OP) THEN
  BEGIN INTEGER L, U;
  OPDN := TRUE; M(LE IN(IND, L)) THEN INP(EO) ELSE EALSE, † AFTER OPSYMBOL †; M(RTXT(Q),
  † ROUTINTEXT MISSING IN OP, DCL †); LE Q = G THEN
  BEGIN POS(L, Z, G); U := B(Z); B(U) := LE B(U) = 1 THEN -10 ELSE - 1 END;
  END
  ELSE OPDN := EALSE;
BOOLEAN PROCEDURE PACK(X); BOOLEAN X; PACK := ANDP(INP(OPEN), ANDP(X, INP(CL)));
BOOLEAN PROCEDURE PIDN(Q); VALUE Q; INTEGER Q; COMMENT PROC IDENTITY DECL. ;
BEGIN INTEGER L;
  LE ANDP(INP(PROC), ANDP(IN(TAG, L), INP(EO))) THEN
  BEGIN PIDN := TRUE; M(RTXT(Q), † NO ROUTINTEXT AFTER †); POS(L + 10000, Z, Q) END
  ELSE PIDN := EALSE
END;

BOOLEAN PROCEDURE POLD, N, I, T1, L, W;
  BEGIN INTEGER VIC, T, Q; VALUE VIC, Q; INTEGER VIC, T, Q;
  POLD := † LE Q = G THEN
  BEGIN LE PLAN(VIC, T, F) THEN
  BEGIN PLAN := TRUE; P := POLD; T := ASK;
  N := LE INP(OPEN) THEN (LE VIC = FML THEN DTC ELSE LSTC(DCL(VRL, T, F))) ELSE 0;
  A(T) := PROC B(T) := W := ASK(N + 2); A(W) := 0; B(W) := 0;
  EQB I := 1 SIER 1 UNTIL N DO
  BEGIN MDCL(VIC, T, Q) † VIC = FML † I > 1, † DCL MISSING IN PLAN †;
  LE VIC = FML THEN M(IN(TAG, L), † MISSING IDENTIFIER IN PLAN †);
  B(W + 1 + I) := T1; LE I < N THEN M(INP(COM),
  † MISSING, IN PLAN †) ELSE M(INP(CL), † MISSING) IN PLAN †;
  END;
  M(DCL(VRL, T1, Q), † DCL MISSING AFTER PLAN †); B(W + 1) := T1
  END
  ELSE PLAN := EALSE
  END
  ELSE
  BEGIN LE INP(OPEN) THEN P := MATCH; PLAN := DCL(VRL, T, F) END
  END;

PROCEDURE POS(L, V, Q); VALUE L, V, Q; INTEGER L, V, Q; COMMENT POSSESS;
  LE Q = G THEN
  BEGIN JGLB := J + 1; LIND(J) := L; VIND(J) := V END;
  BOOLEAN PROCEDURE PRDR(VIC, T, Q); VALUE VIC, Q; INTEGER VIC, Q; COMMENT PROCEDURE DECLARATOR;

```

```

PROR:= ANDP(INP(PROC), PLAN(VRL, Y, O));

BOOLEAN PROCEDURE PRIM(Q); INTEGER O; COMMENT PRIMARY;
BEGIN BOOLEAN PR;
  INTEGER S;
  PR:= PR + LE DEN(Q) THEN ISSUE ELSE LE IDP(Q) THEN ISSUE ELSE LE MIP(Q) THEN ISSUE ELSE LE
  CAST(Q) THEN ISSUE ELSE ENCL(Q);
  FOR I:= 1 WHILE LE PR THEN INP(OPEN) ELSE FALSE DO
  BEGIN IF O = 6 THEN
    BEGIN FOR I:= 1 WHILE A(Z) = REF ^ A(B(Z)) + ROW DO DREF; S:= A(Z);
      LE S = PROC THEN CLTL(Q) ELSE SLTL(Q)
    END
  ELSE P:= MATCH
  END;
END;

PROCEDURE PTOP(N, OP, PR); INTEGER N, OP, PR; COMMENT MOVES N OPERATORS FROM ARR TO VALUE TABLE;
BEGIN INTEGER JJ;
  FOR JJ:= 1 STEP 1 UNTIL N DO
  BEGIN I:= (JJ - 1) * 3; OD(OP, ARR(I + 1), ARR(I + 2), ARR(I + 3), PR) END;
  I:= 0
END;

PROCEDURE PUTBITS(X); VALUE X; INTEGER X;
BEGIN Z:= ASK; A(Z):= BITS; B(Z):= X END;

PROCEDURE PUTBOOL(X); VALUE X; BOOLEAN X;
BEGIN Z:= ASK; A(Z):= BOOL; B(Z):= LE X THEN 1 ELSE 0 END;

PROCEDURE PUTCHAR(X); VALUE X; INTEGER X;
BEGIN Z:= ASK; A(Z):= CHAR; B(Z):= X END;

PROCEDURE PUTCPL(RE, IM); VALUE RE, IM; REAL RE, IM;
BEGIN INTEGER V, VRE, VIM, VI;
  Z:= ASK; A(Z):= STR; B(Z):= V + ASKN(2); C(Z):= 2; V1:= V + 1; B(V):= VRE:= ASK; C(V):= LRE;
  B(V1):= VIM:= ASK; C(V1):= LIM; A(VRE):= REAL; C(VRE):= RE; E(VIM):= IM
END;

PROCEDURE PUTINT(I); VALUE I; INTEGER I;
BEGIN Z:= ASK; A(Z):= INT; B(Z):= I END;

PROCEDURE PUTREAL(R); VALUE R; REAL R;
BEGIN Z:= ASK; A(Z):= REAL; E(Z):= R END;

PROCEDURE PUTSTRING(L); VALUE L; INTEGER L;
BEGIN INTEGER W, I, N, K;
  N:= MISC(L); K:= LE N > 1 THEN N ELSE 1; Z:= ASKN(3 + K); I:= Z + 1; W:= Z + 2; A(Z):= ROW;
  B(Z):= V; C(Z):= 1; A(V):= 0; B(V):= W + 1; C(V):= 1; E(V):= K; A(W):= 1; B(W):= N; C(W):= 1;
  COPGEN(1, 1, K, CHAR, A(W + 1)); COPGEN(1, 1, N, MISC(L + 1), B(W + 1))
END;

```



```

BOOLEAN PROCEDURE PYDN(Q); VALUE Q; INTEGER Q; COMMENT PROC VARIABLE DECL. ;
BEGIN
  INTEGER L, V;
  IF ANDP(LE INP(HEAP) ~ ISSUE THEN INP(PROC) ELSE FALSE, ANDP(INITAG, L), INP(BEC)) THEN
    BEGIN
      PYDN:= ISSUE; M(RTXI(Q), † NO ROUTINETEXT †); LE Q * G THEN
        BEGIN
          V:= ASK; POS(L, V, G); ALVJ:= REF; B(V):= Z END
        END
      ELSE PYDN:= FALSE
    END;
END;

BOOLEAN PROCEDURE PYDF(Q); VALUE Q; INTEGER Q; COMMENT PRIORITY DEF. ;
BEGIN
  INTEGER L, LI, J, V;
  LE IN(IND, L) THEN
    BEGIN
      PYDF:= ISSUE; M(INP(Q), † MISSING IN PRIODEF †); M(IN(INT, LI),
        † PRIONUMBER MISSING †); LE Q * G THEN
        BEGIN
          M(SRCH(L, J), † PRIODEF FOR UNKNOWN OP. †);
          FOR V:= B(VIND(J)), ALV) WHILE V > 0 DO LE B(V) † -.10 THEN B(V):= MISC(LI + †);
        END
      END
    END
  ELSE PYDF:= FALSE
END;

BOOLEAN PROCEDURE PYDN(Q); VALUE Q; INTEGER Q; COMMENT PRIORITY DECL. ;
LE INP(PRIO) THEN
  BEGIN
    PYDN:= ISSUE; M(LIST(PYDF(Q)), † PRIORITY DEF MISSING †) END
  ELSE PYDN:= FALSE;
END;

BOOLEAN PROCEDURE RFRD(VIC, T, Q); VALUE VIC, Q; INTEGER VIC, T, Q;
COMMENT REFERENCE TO MODE DECLARATOR;
LE INP(REF) THEN
  BEGIN
    INTEGER V;
    RFRD:= ISSUE; M(DCL(LE VIC * FML * X(P) † REF THEN FML ELSE VRL, V, Q), † AFTER REF †);
    LE Q * G THEN
      BEGIN
        T:= NRT(V); RX(V) END
      END
    ELSE RFRD:= FALSE;
  END;

BOOLEAN PROCEDURE RL2; COMMENT WIDENING TO REAL FOR INT/REAL OPERATIONS;
BEGIN
  BOOLEAN BRL2;
  BRL2:= ISSUE; LE M1 = INT THEN EV1:= B(V) ELSE BRL2:= M1 = REAL;
  LE M2 = INT THEN EV2:= B(V) ELSE BRL2:= BRL2 ^ M2 = REAL; RL2:= BRL2
END;

PROCEDURE RMCL(APST); VALUE APST; INTEGER APST; COMMENT REMOVE COLLAT. ;
LE ~NEG(Z, APST) THEN
  BEGIN
    INTEGER NSUB, BZ, FEZ, NELZ, NCL, VNEW, VNEW1, NELN, FEN, VNN, VCL1, V1, BPS1, FEP, V1,
    VD, VDI, ZO, N, I, BFE1;
    M(A(Z) = A(APST), † RMCL †); LE A(Z) = ROW THEN
      BEGIN
        NSUB:= A(Z); BZ:= B(Z); FEZ:= B(BZ); NELZ:= E(BZ); M(A(FEZ) = COLL, † NO COLL. IN ROW †);
        NCL:= C(FEZ); LE NSUB < C(APST) THEN
          BEGIN
            VNEW:= ASKN(NSUB + 3); A(VNEW):= ROW; B(VNEW):= VNEW + 1;
            C(VNEW):= NSUB + 1; A(VNEW1):= 1; NELN:= NELZ * NCL; E(VNEW1):= NELN;
            B(VNEW1):= FEN; ASKN(NELN); C(VNEW1):= 1;
            FOR I:= 1 SIZE 1 UNTIL NSUB DO
              BEGIN
                A(VNEW1 + I):= A(BZ + I); B(VNEW1 + I):= B(BZ + I);
                C(VNEW1 + I):= C(BZ + I) * NCL
              END
            END
          END
        ELSE
          BEGIN
            VNEW:= ASKN(NSUB + 3); A(VNEW):= ROW; B(VNEW):= VNEW + 1;
            C(VNEW):= NSUB + 1; A(VNEW1):= 1; NELN:= NELZ * NCL; E(VNEW1):= NELN;
            B(VNEW1):= FEN; ASKN(NELN); C(VNEW1):= 1;
            FOR I:= 1 SIZE 1 UNTIL NSUB DO
              BEGIN
                A(VNEW1 + I):= A(BZ + I); B(VNEW1 + I):= B(BZ + I);
                C(VNEW1 + I):= C(BZ + I) * NCL
              END
            END
          END
        END
      END
    END
  END;

```

```

END;
VNN:= VNEW1 * 1 + NSUB; A[VNN]:= 1; B[VNN]:= NCL; C[VNN]:= 1;
FOR I:= 1 SIZE 1 UNTIL NELZ DO
  BEGIN VCL:= FEZ - 1 + I; M[VCL] := COLL ^ C[VCL] * NCL, #N ELTS IN RMCL;
  V1:= B[VCL]; COPYN(V1, FEZ + (I - 1) * NCL, NCL); FREE(V1, NCL)
END;
Z:= VNEW; RMCL(APST)
END
ELSE
  BEGIN M(NSUB = C(APST), #RMCL NSUB); BPST:= B(APST); FEZ:= B(BPST); M(A[FEZ] * STR,
    #NO STRUCT IN RMCL); M(C[FEZ] * NCL, #N STRUCT ELTS IN RMCL); BFEZ:= B[FEZ] - 1;
    FOR I:= 1 SIZE 1 UNTIL NELZ DO
      BEGIN V1:= FEZ - 1 + I; M[V1] := COLL, #ELT OF ROW NOT COLL; A[V1]:= STR;
        V1:= B[V1] - 1; B[V1]:= VD:= ASK(NCL); VD:= VD - 1; Z0:= Z;
        FOR L:= 1 SIZE 1 UNTIL NCL DO
          BEGIN Z:= V1 + L; LE A[Z] := COLL THEN DCOL(B[FEZ + L]); B[VD + L]:= Z;
            C[VD + L]:= C[FEZ + L]
          END;
        Z:= Z0
      END
    END
  END
END

```

```

END
ELSE
  BEGIN M(A[Z] = STR, #NO ROW OR STR IN RMCL); B[Z] := B(Z); BPST:= B(APST); Z0:= Z; N:= C(Z);
    M(N = C(APST), #N ELTS OF STRUCT IN RMCL);
    FOR I:= 1 SIZE 1 UNTIL N DO
      BEGIN Z:= B[0Z - 1 + I]; DCOL(B[BPST - 1 + I]); B[0Z - 1 + I]:= Z END;
        Z:= Z0
      END
    END
  END
END;

```

```

PROCEDURE RNC(1, S); VALUE I; INTEGER I; STRING S;
BEGIN NLCR; NLCR; ABSFIX(6, 0, 1); PRINTTEXT(#TH INDEXES TOO); PRINTTEXT(0); ER(1) END;

```

```

BOOLEAN PROCEDURE ROP(S, Y); VALUE S; INTEGER S; REAL Y; LE PP = S THEN
  BEGIN ROP:= ISSUE; PUTREAL(Y) END
  ELSE ROP:= FALSE;

```

```

BOOLEAN PROCEDURE ROP4(S1, R1, S2, R2, S3, R3, S4, R4); INTEGER S1, S2, S3, S4; REAL R1, R2, R3, R4;
  LE PP = S1 THEN PUTREAL(R1) ELSE LE PP = S2 THEN PUTREAL(R2) ELSE LE PP = S3 THEN PUTREAL(R3)
  ELSE LE PP = S4 THEN PUTREAL(R4) ELSE ROP4:= FALSE
END;

```

```

PROCEDURE RS(V); VALUE V; INTEGER V; LE V > 0 THEN
  BEGIN DM(V); FREE(V) END;

```

```

BOOLEAN PROCEDURE RTXT(Q); VALUE Q; INTEGER Q; COMMENT ROUTINE TEXT;
  LE LK2(PLAN(FML, DUM, F), INP(COL)) THEN
  BEGIN INTEGER P1, V;
    RTXT:= ISSUE; P1:= P1 M(LE LE PLAN(FML, V, Q) THEN INP(COL) ELSE FALSE THEN UNIT(P) ELSE FALSE,
    #ROUTINETEXT); LE Q = 0 THEN
    BEGIN Z:= V; CIB(V):= P1; MORP:= ISSUE END
  END
  ELSE RTXT:= FALSE;

```

Vp: 7 Swell?

Vp: 7

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```

BOOLEAN PROCEDURE RWDR(VIC, T, Q) VALUE VIC, Q; INTEGER VIC, T, Q; SOURCE: ROW DECLARATOR;
BEGIN INTEGER POLD, V, ST, N, VD, I, NLT, VM, V1, J;

BOOLEAN PROCEDURE RWRW(VIC, Q) VALUE VIC, Q; INTEGER VIC, Q;
BEGIN BOOLEAN WORK;
INTEGER J, LB, UB;
WORK := Q = G; JE WORK THEN J := VD + 1; RWRW := EALSI; JE VIC = ATL THEN
BEGIN JE UNIT(Q) THEN
BEGIN RWRW := ISSUE; M(COER(MEEK, STIN, Q), $ROWROWER MODE?); JE INP(COL) THEN
BEGIN LB := A(J) + 1; BIZ := BIZ; EMP END;
M(UNIT(Q), $MISSING-U,B,?); M(COER(MEEK, STIN, Q), $U,B. MODE IN DCL?);
JE WORK THEN
BEGIN UB := BIZ; BIZ := BIZ; EMP END;
ELSE JE WORK THEN
BEGIN LB := A(J) + 1; UB := BIZ; BIZ := BIZ; EMP END;
JE WORK THEN NLT := NLT + (JE UB < LB THEN 0 ELSE UB - LB + 1)
END
END
ELSE
BEGIN JE WORK THEN A(J) := BIZ; JE RWRW := ISSUE; JE UNIT(Q) ~ WORK THEN A(J) := MEEKINT(Q);
JE INP(COL) THEN J := JE UNIT(Q) ~ WORK THEN B(J) := MEEKINT(Q);
JE WORK THEN NLT := NLT + (JE B(J) < A(J) THEN 0 ELSE B(J) - A(J) + 1)
END
END;

LE Q = G THEN
BEGIN POLD := P; JE RWDR(VIC, T, F) THEN
BEGIN RWDR := ISSUE; P := POLD; T := V; ASK; A(V) := ROW; ST := JE INP(FLX) THEN 0 ELSE 1;
JE INP(FLX) THEN J := M(INP(OPEN), $MISSING IN DCL?); CIV := NIS; LSTC(RWRW(VIC, F));
B(V) := VD := ASK(N + 1); A(VD) := ST; C(VD) := J; NLT := 1;
FOR I := 1 STEP 1 UNTIL N DO
BEGIN RWRW(VIC, G) JE I = N THEN M(INP(COL), $MISSING IN DCL?) ELSE M(INP(COM),
$MISSING, IN ROWDCL?)
END;
M(DCL(JE XIP) = REF THEN VRL ELSE VIC, VM, Q), $ROW DCL?);
JE VIC = VRL THEN NLT := 1; E(VD) := NLT; JE NLT < 1 THEN B(VD) := VM ELSE
BEGIN B(VD) := V1 := ASK(NELT);
FOR I := 1 STEP 1 UNTIL NLT DO MOVE(VM, V1 + 1 - 1); RS(VM);
END;
JE N > 0 THEN
BEGIN C(VD + N) := 1;
FOR I := N STEP - 1 UNTIL 2 DO
BEGIN J := VD + 1; C(J - 1) := (B(J) - A(J) + 1) * C(J); END
END;
END
ELSE RWDR := EALSI
END
ELSE RWDR := ANDP(INP(FLX) ~ ISSUE, ANDP(PACK(LIST(RWRW(VIC, F)) ~ ISSUE), DCL(VIC, T, F));
END;

PROCEDURE RX(V) VALUE V; INTEGER V; JE V > 0 THEN
BEGIN DIV := DIV - 1; JE DIV < 0 THEN
BEGIN DP(V); FREE(V) END
END;

```

```

PROCEDURE RXN(V, N); VALUE V, N; INTEGER V, N;
BEGIN INTEGER V;
  W := V * N - 1;
  FOR V := V SIER 1 UNTIL W DO RX(V)
END;

BOOLEAN PROCEDURE SCL(Q); INTEGER Q; COMMENT SERIAL CLAUSE;
BEGIN INTEGER P1, P2, Q1, Q2;
  BOOLEAN LAB1, QF;

  BOOLEAN PROCEDURE JMP; LE Q = G ^ Q1 = F THEN
  BEGIN JMP := TRUE; LE PLAB > P1 ^ PLAB < P2 THEN
  BEGIN P1 = PLAB; Q1 = G END
  ELSE
  BEGIN Q1 = P; P1 = P2; QF := TRUE END
  ELSE JMP := FALSE;

  BOOLEAN PROCEDURE SEP; LE IMP(GOON) THEN
  BEGIN SEP := TRUE; LE JMP THEN ELSE LE Q = G THEN
  BEGIN LE -MORF THEN ELSE
  FOR I := 1 WHILE DPRO(Q) > DREF DO I EMP
  END
  ELSE LE IMP(PNT) THEN
  ELSE JMP THEN ELSE LE Q = G THEN P1 = P2 END
  BEGIN SEP := JMP ^ FALSE; P1 = P; SCL := FALSE; QF := FALSE;
  LE (LE Q = G THEN SCL(M) ELSE TRUE) THEN
  BEGIN P2 = P; P1 = P1; LAB1 := FALSE;
  FOR I := 0, 1 WHILE LE SEP THEN (LE Q = G ^ QF THEN P < P2 ELSE TRUE) ELSE FALSE DO
  BEGIN Q1 := Q; LE LAB(Q1) THEN LAB1 := TRUE;
  LE (LE LAB1 THEN ELSE ELSE LIST(UNION(Q1))) THEN ELSE
  BEGIN LE I = 0 THEN SCL := LBUN(Q1) ELSE M(LBUN(Q1), $ AFTER) OR ($); SCL := TRUE
  END
  END
END;

BOOLEAN PROCEDURE SCL1; COMMENT STRING OR CHAR PREP. FOR LSTR, I
LE STCH(V1, V2) THEN
BEGIN SCL1 := TRUE; MINL := LE LEN1 < LEN2 THEN LEN1 ELSE LEN2; I := 0;
FOR I := 1 + 1 WHILE LE I < MINL THEN BIFEL := I + 1; B(FE2 - I + 1) ELSE FALSE DO;
  UNEQ := I < MINL
END
END;

BOOLEAN PROCEDURE SEC(Q); INTEGER Q; SEC := LE GEN(Q) THEN TRUE ELSE LE SEL(Q) THEN TRUE ELSE PRIM(Q);
COMMENT SECONDARY;

BOOLEAN PROCEDURE SEL(Q); INTEGER Q; COMMENT SELECTION;
BEGIN INTEGER L, I, W, N, U;
  BOOLEAN NAME;
  LE ANDP(IN(TAG, L), IMP(OF)) THEN
  BEGIN SEL := TRUE; M(SEC(Q), $ AFTER OF); LE WEAK(Q) THEN
  BEGIN NAME := ALZ) = REF; LE NAME THEN ZIB BIZ; MORF := M(AZ) = STR, $SEL);
  W := B(Z); N := C(Z); I := 1;
  FOR I := 1 WHILE LE I > N THEN ELSE ELSE C(W + I - 1) + L DO I := I + 1; M(I) & N;

```

↑SELECTOR↑) U:= B(W * I - 1) Z:= LE NAME THEN NRT(U) ELSE NINST(U)

END
ELSE SEL:= EALSE
END/

BOOLEAN PROCEDURE SEQ(A); BOOLEAN A; COMMENT SEQUENCE;

IF A THEN
BEGIN SEQ:= ISSUE;
FOR I:= 1 WHILE A DO
END
ELSE SEQ:= EALSE;

BOOLEAN PROCEDURE SKIP(Q); INTEGER Q; LE INP(SKIP) THEN

BEGIN SKIP:= ISSUE; LE Q = 0 THEN
BEGIN Z:= ASK; A[Z]:= VOID END
END
ELSE SKIP:= EALSE;

PROCEDURE SLT(L(Q)); INTEGER Q; COMMENT SLICE TAIL;

BEGIN INTEGER VO, VD, N, FE, K, OFST, I, N1, J, NLT, LTR, UTR, W, VNW, FENW, VREF, VDI, TS1, L1;
BOOLEAN NAME, TSIR;
MORF:= ISSUE; NAME:= A[Z] = REF; LE NAME THEN VREF:= Z; VO:= LE NAME THEN B[VREF] ELSE Z;
M(ATVO) = ROW, † SLICE †; VD:= B[VO]; N1:= N(C[VO]); FE:= B[VD]; OFST:= C[VD];
VD1:= ASKN(N + 1); COPYN(VD, VD1, N + 1); A[VD1]:= 1; NLT:= E[VD];
FOR I:= 1 STEP 1 UNTIL N DO
BEGIN TSIR:= UNIT(Q); J:= VD1 + 1; LE TSIR THEN TS1:= MEKINT(Q);
LE SPY(COL) ∨ SPY(AT) ∨ TSIR THEN
BEGIN COMMENT TRIMMER;
LTR:= LE TSIR THEN TS1 ELSE A[J]; UTR:= B[J]; LE INP(COL) THEN
BEGIN LE UNIT(Q) THEN UTR:= MEKINT(Q) END;
LE INP(AT) THEN
BEGIN M(UNIT(Q), † AFTER AT †); L1:= MEKINT(Q) END
ELSE L1:= 1; LE LTR < A[J] THEN RGE(L, † LOW †); LE UTR > B[J] THEN RGE(L,
† HIGH †); OFST:= OFST + (LTR - A[J]) + C[J]; A[J]:= L1; B[J]:= L1 - LTR + UTR

END
ELSE
BEGIN COMMENT SUBSCRIPT;
K:= TS1; OFST:= OFST + (K - A[J]) + C[J]; B[J]:= - 1; N1:= N1 - 1
END;
IF I < N THEN M(INP(COM), † TOO FEW TRIMSCRIPT †) ELSE M(INP(CL),
† TOO MANY TRIMSCRIPTS †)

END;
C[VD1]:= OFST; LE N1 > 0 THEN
BEGIN LE N1 < N THEN
BEGIN W:= ASKN(N1 + 1);
FOR I:= 1 STEP 1 UNTIL N DO LE D(I) ≠ - 1 THEN
BEGIN COPY(VD1 + 1, W + 1); B[W]:= W + 1 END;
FREEN(VD1, N + 1)

END
ELSE W:= VD1; VNW:= ASK; A[VNW]:= ROW; B[VNW]:= W; C[VNW]:= N1

END
ELSE
BEGIN VNW:= FE + OFST - 1; FREEN(VD1, N + 1) END;
LE NAME THEN Z:= NRT(VNW) ELSE LE N1 > 0 THEN
BEGIN B[W]:= FENW:= ASKN(NLT); MOVEN(FE, FENW, NLT); Z:= VNW END
ELSE
BEGIN Z:= ASK; MOVE(VNW, Z) END

```

END;

BOOLEAN PROCEDURE SOFTC(Q); INTEGER Q; COMMENT SOFT COERCION;
BEGIN FOR I:= 1 WHILE LE Q = G THEN DPRO(Q) ELSE FALSE DO ; SOFTC:= Q = G END;

PROCEDURE SPS(V, W); VALUE V, W; INTEGER V, W; COMMENT SUPERSEDE;
BEGIN OM(W); MOVE(V, W) END;

BOOLEAN PROCEDURE SPY(S); VALUE S; INTEGER S; LE INP(S) THEN
BEGIN SPY:= TRUE; P:= 1 END
ELSE SPY:= FALSE;

BOOLEAN PROCEDURE SRCH(L, K); VALUE L; INTEGER L, K; COMMENT SEARCH IN LINK TABLE;
BEGIN INTEGER JJ;
JJ:= J;
FOR I:= 1 WHILE LE JJ > 0 THEN L ← LIND(JJ) ELSE FALSE DO JJ:= JJ - 1; K:= JJ; SRCH:= JJ > 0
END;

BOOLEAN PROCEDURE STCH(U, V); VALUE U, V; INTEGER U, V;
COMMENT OBTAINS FIRST ELEMENTS AND LENGTHS OF TWO STRINGS IF POSSIBLE;
STCH:= LE STRC(U, FE1, LEN1) THEN STRC(V, FE2, LEN2) ELSE FALSE;

BOOLEAN PROCEDURE STRD(VIC, T, Q); VALUE VIC, Q; INTEGER VIC, T, Q; COMMENT STRUCT DECLARATOR;
LE INP(STR) THEN
BEGIN INTEGER V, VD, N, V1, I, VM, L;
BOOLEAN F1;
STOR:= TRUE; M(IN(OPEN), {NO(AFTER STRUCT)}); LE Q = G THEN
BEGIN V:= ASK; A(V):= STR; C(V):= N:= DTG; B(V):= VD:= ASK(N); F1:= TRUE;
FOR I:= 1 STEP 1 UNTIL N DO
BEGIN LE F1 THEN F1:= FALSE ELSE M(IN(PCOM), {MISSING, IN STRUCTDCL});
M(OC(LE X(P) = REP THEN VRL ELSE VIC, VM, Q), {MISSING DCL}); B(VD + 1 := 1); I:= VM;
M(IN(TAG, L), {MISSING TAG}); C(VD + 1 := 1); I:= L
END;
T:= V; M(IN(CL), {MISSING}) IN STRUCTDCL;
END
ELSE P:= MATCH
ELSE STOR:= FALSE;

PROCEDURE STPN(F, N); VALUE F, N; INTEGER F, N; COMMENT STANDARD FUNCTION;
BEGIN OM( - F, STRL, STRL) B(21):= 1; L:= L + N END;

BOOLEAN PROCEDURE STRC(V, FE, LEN); VALUE V; INTEGER V, FE, LEN;
BEGIN INTEGER VD;
COMMENT LOOK IF V IS STRING OR CHAR;
STRC:= FALSE; LE A(V) = ROW THEN
BEGIN LE C(V) = 1 THEN
BEGIN VD:= B(V); FE:= B(VD); LE A(FE) = CHAR THEN
BEGIN STRC:= TRUE; LEN:= B(VD + 1) END
END
ELSE LE A(V) = CHAR THEN
BEGIN STRC:= TRUE; FE:= V; LEN:= 1 END
END;

```

```

PROCEDURE SUMS; COMMENT FORMS SUM OF TWO STRINGS;
BEGIN
  INTEGER N, Z1, Z2, Z3;
  N := LEN1 + LEN2; Z1 := ASKN(N + 3); A(Z1) := ROW; B(Z1) := Z1 * Z + 1; C(Z1) := 1; A(Z1) := 0;
  B(Z1) := Z3 * Z + 3; C(Z1) := 1; E(Z1) := N; Z2 := Z + 2; A(Z2) := 1; B(Z2) := N; C(Z2) := 1;
  COPYN(FE1, Z3, LEN1); COPYN(FE2, Z3 + LEN1, LEN2)
END;

```

```

BOOLEAN PROCEDURE TERT(Q); INTEGER Q; TERT := LE FORM(Q) THEN TRUE ELSE SEC(Q);

```

```

COMMENT TERTIARY;

```

```

BOOLEAN PROCEDURE UNDN(Q); VALUE Q; INTEGER Q; COMMENT UNITARY DECLARATION;

```

```

BEGIN
  BOOLEAN SUC;
  BEGIN
    INTEGER P0;
    P0 := P;
    UNDN := SUC := LE MDN(Q) THEN TRUE ELSE LE PYDN(Q) THEN TRUE ELSE LE OPDN(Q) THEN TRUE ELSE
    IRDN(Q) LE SUC THEN
    BEGIN
      LE X(P) + GOON ^ X(P) + COM THEN
      BEGIN UNDN := SUC := FALSE; P := P0 END;
    END;
    LE SUC ^ Q = G THEN
    BEGIN D(I) := D(I) + 1; Z := 1 END;
  END;
END;

```

```

BOOLEAN PROCEDURE UNDR(VIC, T, Q); VALUE VIC, Q; INTEGER VIC, T, Q; COMMENT UNION DECLARATOR;
LE INP(UN) THEN
BEGIN
  INTEGER N, W, U, I;
  UNDR := TRUE; M(INP(OPEN), $MISSING(IN UNIONDCL $)) LE Q = G THEN
  BEGIN
    N := LSTC(DCL(VRL, T, F)); T := ASK; A(T) := UN; W := B(T) := ASKN(N); C(T) := Q; C(W) := N;
    FOR I := 1 STEP 1 UNTIL N DO
      BEGIN
        M(DCL(VRL, U, Q)); $DCL MISSING IN UNIONDCL $; B(W + 1) := 1; U := U;
        LE I < N THEN M(INP(COM), $, MISSING IN UNIONDCL $) ELSE M(INP(CL),
        $) MISSING IN UNIONDCL $;
      END;
    B(T) := DEV(W)
  END
  ELSE P := MATCH
  ELSE UNDR := FALSE;

```

```

BOOLEAN PROCEDURE UNIT(Q); INTEGER Q; LE LOOP(Q) THEN UNIT := TRUE ELSE
BEGIN
  INTEGER POLD, L, NEXT;
  BOOLEAN POS;
  POLD := P; LE INP(OPEN) THEN P := MATCH; NEXT := X(P);
  POS := LE NEXT > 91 ^ NEXT < 119 THEN TRUE ELSE IN(IND, L); P := POLD;
  LE LE POS THEN RTXT(Q) ELSE FALSE THEN UNIT := TRUE ELSE LE TERT(Q) THEN
  BEGIN
    UNIT := TRUE;
    LE INP(SEC) THEN ASTL(Q) ELSE LE INP(IS) THEN IRTL(Q, IS) ELSE LE INP(ISNT) THEN
    IRTL(Q, ISNT)
  END
  ELSE UNIT := FALSE
END;

```

```

BOOLEAN PROCEDURE UNMB(V, W); VALUE V, W; INTEGER V, W; COMMENT UNION MEMBER;

```

```

BEGIN INTEGER N, I;
  N:= C(W); I:= 1;
  FOR I:= 1 + 1 WHILE LE I = N THEN ELSE ELSE -MEO(BIV), B(W + 1)) DO ; UNMB:= I + N
END;

BOOLEAN PROCEDURE VAC(Q); VALUE Q; INTEGER Q; LE ANDP(INP(OPEN), INP(CL)) THEN
BEGIN VACT:= IBUE; LE Q = G THEN
  BEGIN Z:= ASK; A(Z):= VOID END.
END
ELSE VACT:= EALSE;

BOOLEAN PROCEDURE VOF(Q, U); VALUE Q, U; INTEGER Q, U; COMMENT VARIABLE DEFINITION;
BEGIN INTEGER L, V, W;
  LE IN(YAG, L) THEN
  BEGIN VDF:= IBUE; LE Q = G THEN
  BEGIN V:= NINST(U); W:= NRT(V); LE INP(BEC) THEN
  BEGIN MUNIT(Q), (AFTER:=); M(COER(STRONG, U, Q), (SOURCE MADE IN VAR DEF));
    SPS(Z, V) EMP
  END;
  RX(V); POS(L, W, Q)
  END
  ELSE LE LE INP(BEC) THEN UNIT(P) ELSE EALSE. THEN
  END
  ELSE VDF:= EALSE
END;

BOOLEAN PROCEDURE VON(Q); VALUE Q; INTEGER Q; VON:= LE PVON(Q) THEN IBUE ELSE MYDN(Q);
COMMENT VARIABLE DECLARATION;

PROCEDURE VISR(J1, J2); VALUE J1, J2; INTEGER J1, J2; COMMENT MAKE VISIBLE;
BEGIN INTEGER JJ;
  FOR JJ:= J1 + 1 STEP 1 UNTIL J2 DO LE LIND(JJ) ≥ 1000 THEN LIND(JJ):= LIND(JJ) - 1000
  END;

BOOLEAN PROCEDURE WEAK(Q); INTEGER Q; COMMENT WEAK COERCION;
BEGIN FOR I:= 1 WHILE LE Q = G THEN (LE LE LE AIZ) = REF THEN AIB(Z)) = REF ELSE EALSE THEN DREF
  ELSE EALSE THEN IBUE ELSE DPRO(Q) ELSE EALSE DO ; WEAK:= Q = G
END;

PROCEDURE WTX(S); VALUE S; INTEGER S;
BEGIN PX:= PX + 1; X(PX):= S. END;

COMMENT WRITE TO X;
PART:= 1000001;
FOR I:= 2 STEP 1 UNTIL NA DO (I - 1):= 1; DIMA:= 0; FBI:= 1;
COMMENT NOW FIRST FOLLOWS STANDARD PRELUDE INFORMATION;
LRE:= 91; LIM:= 94; MISC(3):= MISC(2); MISC(2):= 1; PUTBOOL(IBUE); STBL:= Z; PUTINT(1); STIN:= Z;
PUTREAL(4 + ARCTAN(1)); STRL:= Z; PUTCP(1, 0, 1, 0); STCP:= Z; PUTCHAR(1); STCR:= Z; PUTSTRING(1)
STNG:= Z; RINT:= NRT(STIN); RRL:= NRT(STRL); RSTG:= NRT(STNG); STVD:= ASK; A(STVD):= VOID;
STRV:= ASK; A(STRV):= ROW; L:= 1; OOLD:= 0; OD(PLTO, STCR, RSTG, 1));
OD(PLTO, STNG, RSTG, 1));
FOR MM:= RINT, STIN, RINT, RRL, STIN, RRL, STRL, RRL, RSTG, STCR, RSTG, STNG, STNG, RSTG DO FILL;
PTOP(5, PSAB, 1); PTOP(3, MAB, 1); PTOP(3, THAB, 1); OD(OVAB, RINT, STIN, RINT, 1));
OD(MDAB, RINT, STIN, RINT, 1); OD(OVAB, RRL, STIN, RRL, 1); OD(OVAB, RRL, STRL, RRL, 1);
OD(OR, STBL, STBL, 2); OD(AND, STBL, STBL, 3);
FOR MM:= STIN, STIN, STBL, STBL, STRL, STRL, STBL, STBL, STIN, STBL, STCR, STCR, STBL, STBL,

```



```

STCR, STNG, STBL, STNG, STCR, STBL, STNG, STNG, STBL, STBL, STBL, STBL, STBL, STBL, STBL, STBL, STBL, STBL, STBL, STBL,
STCP, STBL, STCP, STIN, STBL, STCP, STBL, STBL, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP,
PTOP(14, NE, 4)) PTOP(8, LESS, 5)) PTOP(8, GT, 5)) PTOP(8, LE, 5)) PTOP(8, GE, 5))
EOB NM1= STIN, STIN, STIN, STIN, STBL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STIN, STIN, STCP, STCP,
STRL, STCP, STCP, STCP, STCP, STIN, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP,
STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG, STNG,
OM(PLUS, STRL, STRL)) PTOP(9, MIN, 6)) OM(MIN, STIN, STIN)) OM(MIN, STRL, STRL)) PTOP(9, TIMS, 7))
PTOP(9, DIV, 7)) OD(OVER, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN, STIN,
OD(SPLIT, STIN, STIN, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP,
OD(SPLIT, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL, STRL,
OD(UP, STCP, STIN, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP, STCP,
OD(UPB, STIN, STRV, STRV, STIN, 8)) OM(UPB, STRV, STRV, STIN)) OM(NOT, STBL, STBL, STBL, STBL, STBL, STBL, STBL, STBL, STIN))
OM(ABSV, STIN, STIN)) OM(ABSV, STRL, STRL)) OM(SGN, STIN, STIN)) OM(SGN, STRL, STRL)) OM(CONJ, STCP, STCP, STCP, STCP,
OM(OOD, STIN, STBL)) OM(SGN, STIN, STIN)) OM(SGN, STRL, STRL)) OM(IM, STCP, STRL)) OM(CONJ, STCP, STRL))
OM(ARG, STCP, STRL)) OM(REPR, STIN, STCR)) L:= L + 1; POS(L, STRL, G)) L:= L + 2; STFN(NSOR, 4))
STFN(NEXP, 3)) STFN(NLN, 2)) STFN(NCOS, 3)) STFN(NCOS, 6)) STFN(NSIN, 3)) STFN(NSIN, 6))
STFN(NTAN, 3)) STFN(NATN, 6)) OM(LUDD, STRL, STVD)) L:= L + 1; OM(1999, STIN, STVD)) L:= L + 7;
OM(1998, STIN, STVD)) L:= L + 7; OM(1997, STIN, STVD)) L:= L + 5; I:= READ; JE I = - 1 THEN EXIT;
M(I) = 1, {INCORRECT INPUT?} NEWPAGE;
EOB I:= RESYM XMLE I + 122 DD PRSYM(1)) NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR;
PXI= NXI:= READ;
EOB I:= 1 STEP 1 UNTIL NXI DD X(1):= READ; NM1:= READ;
EOB L:= 1 STEP 1 UNTIL NM1 DD MISC(L):= READ; NR1:= READ;
EOB L:= 1 STEP 1 UNTIL NR1 DD RNRL(L):= READ; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR; NLCR;
{PROGRAM};
    
```

END
END
END

Appendix B

```
( 'COMMENT' RECURSIVE FACTORIAL OPERATOR 'COMMENT'  
'OP' 'FAC' = ('INT' N) 'INT' : ( N#0 | 1 | N# 'FAC'(N-1))  
PRINT( 'FAC' 5 )  
)
```

```
'BEGIN' 'CO' 'COMPLEX SQUARE ROOT' 'CO'  
'COMPL' Z'  
'PROCT' COMPSORT = ('COMPL' Z) 'COMPL'  
'BEGIN' 'REAL' X'RE' Z, Y='IM' Z'  
'REAL' RP=SQRT((ABS' X + SQRT(X**2+Y**2))/2)'  
'REAL' IP = (RP*0|Y/(2*RP))'  
(X>=0|RP 'I' IP |ABS' IP 'I' (Y>=0|RP|-RP))  
'END'  
Z=COMPSORT(-1)  
PRINT('RE' Z, 'I' 'IM' Z)  
Z=COMPSORT(1 'I' -1)  
PRINT(NEWLINE, 'RE' Z, 'I' 'IM' Z)  
'END'
```

+.109868413469* 1 I -.450898605617*- 0

```
'BEGIN' 'CO' 'TOWERS OF HANOI' 'CO'  
'PROC' P = ('INT' A,B,K) 'VOID':  
'IF' K > 0 'THEN'  
  P(A,6-A-B,K-1)  
  PRINT((NEWLINE,"FROM",A,"TO",B,"PIECE",K))  
  P(6-A-B,B,K-1)  
'FI'  
  P(1,2,4)  
'END'
```

FROM	+1	TO	+3	PIECE	+1
FROM	+3	TO	+2	PIECE	+2
FROM	+1	TO	+3	PIECE	+1
FROM	+2	TO	+1	PIECE	+3
FROM	+1	TO	+3	PIECE	+2
FROM	+3	TO	+2	PIECE	+1
FROM	+2	TO	+2	PIECE	+4
FROM	+1	TO	+1	PIECE	+1
FROM	+3	TO	+2	PIECE	+2
FROM	+2	TO	+1	PIECE	+1
FROM	+3	TO	+3	PIECE	+3
FROM	+1	TO	+2	PIECE	+1
FROM	+3	TO	+3	PIECE	+2
FROM	+1	TO	+2	PIECE	+1

```
'BEGIN' # TRANSLATION FROM DECIMAL TO ROMAN #
'MODE' 'M';STRUCT' ('INT' VALUE, 'STRING' R);
'PROC' ROMANS('INT' NUMBER) 'STRING';
'BEGIN' 'INT' N:=NUMBER, 'STRING' RESULT,
[] 'M' TABLE =
((1000,"M"),(900,"CM"),(500,"D"),(400,"CD"),
(100,"C"),(90,"XC"),(50,"L"),(40,"XL"),
(10,"X"),(9,"IX"),(5,"V"),(4,"IV"),(1,"I"));
'FOR' I 'TO' 'UPB' TABLE 'DO';
'BEGIN' 'INT' V = VALUE 'OF' TABLE[I],
'STRING' R = R 'OF' TABLE[I];
'WHILE' V <= N 'DO' (RESULT +=R) N -=V;
'END';
RESULT
'END';
PRINT(ROMAN(1968))
'END'
```

MCMLXVIII

```
'BEGIN' # TRANSLATION FROM ROMAN TO DECIMAL #
'PROC' CHAR IN STRING = ('CHAR', C, 'REF', 'INT', I, 'STRING', S) 'BOOL' I
('FOR', K 'TO', 'UPB', S 'DO' ((S[K]||=K'I)) 'FALSE', L: 'TRUE' I)
'PROC' VALUE OF ROMAN = ('STRING', TEXT) 'INT' I
  'IF', TEXT = "" 'THEN' 0 'ELSE'
    'TOP', 'ABS', F ('CHAR', S) 'INT' I
    'CASE', 'INT', P: CHAR IN STRING(S,P,"IVXLCDM") P 'IN'
      1,5,10,50,100,500,1000
    'ESAC' I
    'INT', V, MAXV:=0, MAXP:
    'FOR', P 'TO', 'UPB', TEXT 'DO'
      'IF', (V<=ABS' TEXT[P]) > MAXV 'THEN' MAXP:=P MAXV:=V 'F' I I
    MAXV = VALUE OF ROMAN (TEXT[I: MAXP-I])
      + VALUE OF ROMAN (TEXT[MAXP+1: I])
  'F' I I
'PRINT(VALUE OF ROMAN("MCMCLXVII"))
'END'
```

*1968

```
'BEGIN' # IN SITU PERMUTATION #  
  (15) 'REAL' X=(3.14, 2.72, .577, -33.0, 1.41);  
  'PROC' PERMVEC = ('REF', (1)) 'REAL' VEC, (1) 'INT' P 'VOID';  
  'FOR' J 'TOT' 'UPB' P 'DO'  
    'BEGIN' 'INT' K=PI(J) 'WHILE' K>J 'DO' K=P(K);  
      'IF' K<J 'THEN'  
        'REAL' M=VEC(J), 'INT' L=P(K);  
        'WHILE' L<= J 'DO'  
          'BEGIN' VEC(K)=VEC(L); K=L; L=P(K) 'END';  
          VEC(K)=M  
        'IF'  
      'END'  
    'PERMVEC(X, (3,5,4,1,2)); PRINT(X)  
  'END'
```

+ .5770000000002e- 0 + .14100000000000e+ 1 -33 + .3139999999999e+ 1 + .27200000000000e+ 1

```
'BEGIN' # NUMERICAL QUADRATURE #  
'PROC' INT1 ('REAL' A,B,EPS,  
'BEGIN' 'PROC' INT1B ('REAL' X0,X4,F0,F2,F4) 'REAL'  
'BEGIN' 'REAL' X2# (X0+X4)/2; 'REAL' F1=Y((X0+X2)/2),  
F3 = Y((X2+X4)/2);  
'REAL' T;  
'IF' T# 6*F2-4*(F1+F3)+F0+F4) 'ABS' T < DELTAMAX 'THEN'  
(X4-X0)*(4*(F1+F3)+2*F2+F0+F4-T/15) 'ELSE'  
INT1(X0,X2,F0,F1,F2) + INT1(X2,X4,F2,F3,F4)  
'FI'  
'END' #INT1#,  
'REAL' DELTAMAX# 180.*EPS/'ABS' (B-A);  
INT1(A,B,Y(A),Y((A+B)/2),Y(B))/12  
'END';  
PRINT(INT(0,2,.01),('REAL'X) 'REAL'X**2-3*X) #-3.3333...#  
'END'
```

-.3333333333332* 1