IEEE-1588 Standard Version 2

-A Tutorial-

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October 2, 2006
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Outline

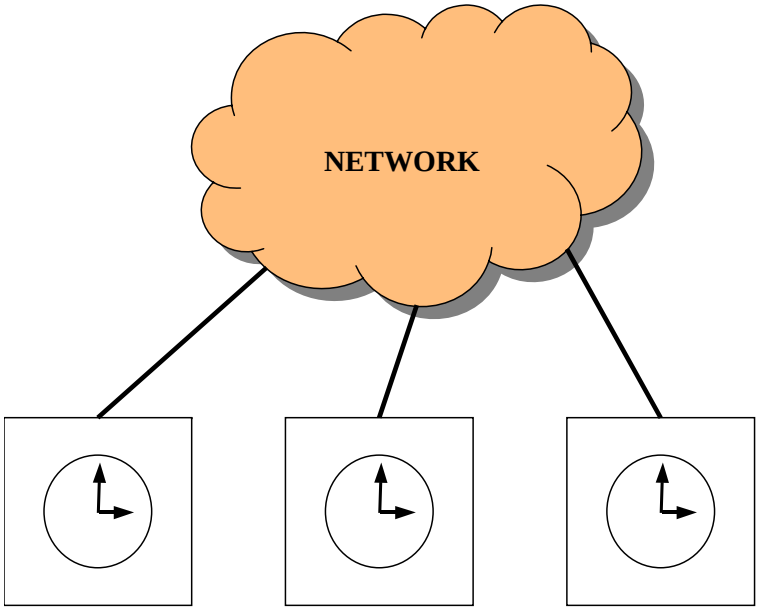
1. What is IEEE 1588
2. Motivation for version 2
3. Version 1 Clocks
4. New Types of Clocks in Version 2
5. Comparison of clock types
6. Message changes in version 2
7. Other changes
8. Status

What is IEEE 1588

- **IEEE 1588 synchronizes real-time clocks in the nodes of a distributed networked system.**
- **Enables a new methodology for measurement and control**

  BASED ON TIME

  NOT ON TIME-OF-RECEIPT-BASED EVENT NOTIFICATION.



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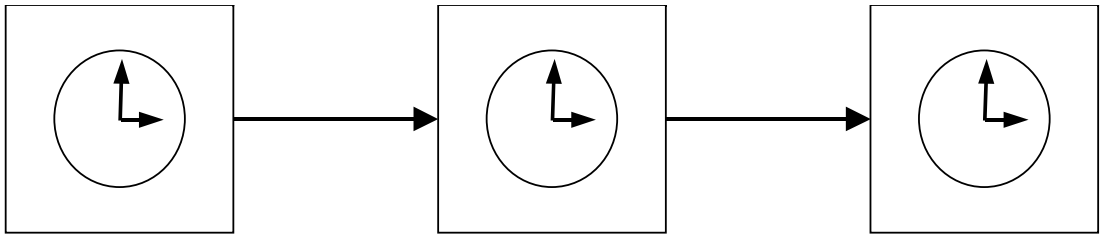
Major motivation for version 2

- **New application areas with new requirements: Telecom, IEEE 802.1as, power industry**
- **New requirements in original applications (IA, T&M, military)**
  - **Higher accuracy**
  - **Varied update rates**
  - **Linear topology (in addition to hierarchical)**
  - **Rapid reconfiguration after network changes**
  - **Fault tolerance**

V1 Synchronization Basics

Step 1: Organize the clocks into a master-slave hierarchy (based on observing the clock property information contained in multicast Sync messages)

Step 2: Each slave synchronizes to its master (based on Sync, Delay_Req, Follow_Up, and Delay_Resp messages exchanged between master and its slave)

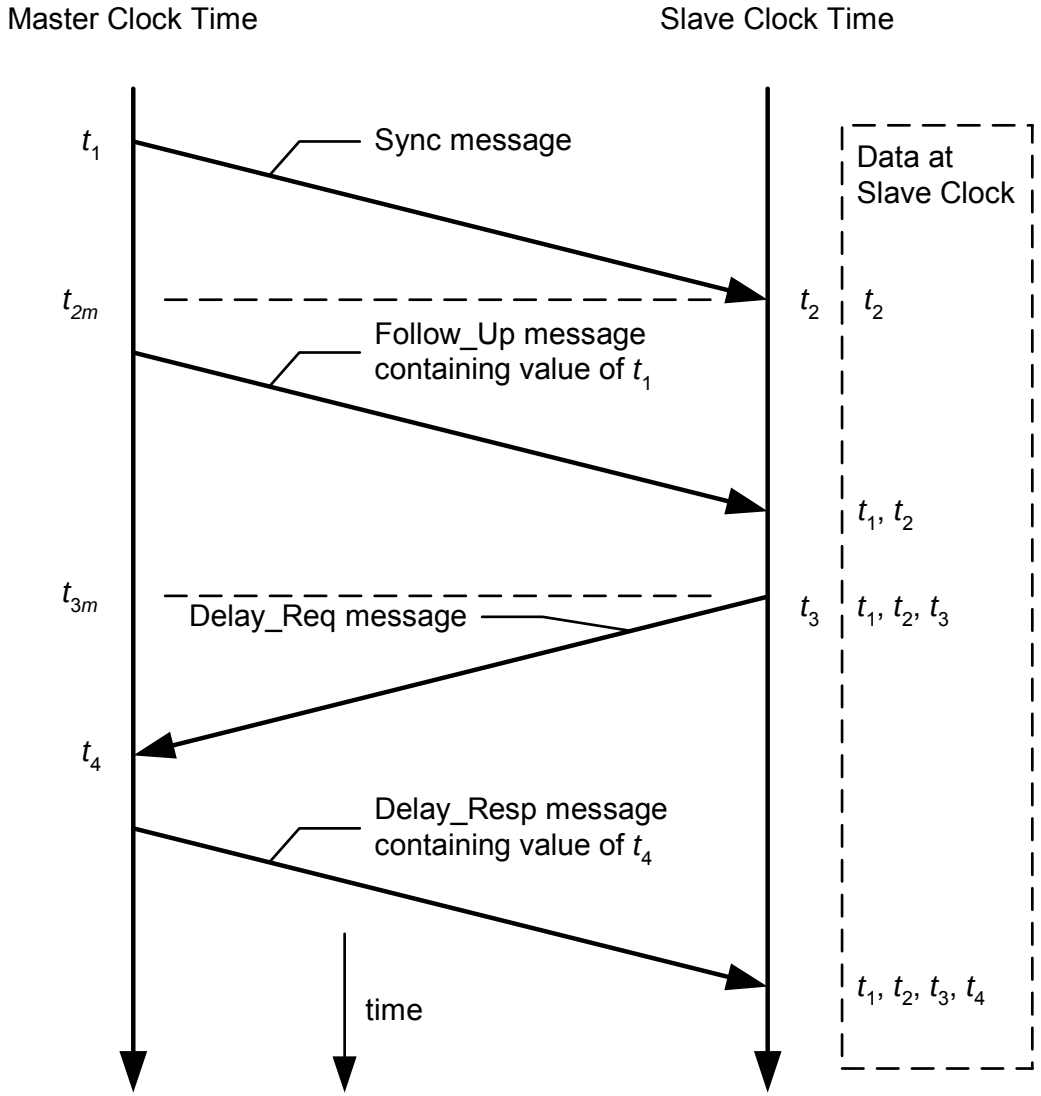


Grandmaster Clock
This clock determines the time base for the system

Slave to the Grandmaster Clock and Master to its Slave

Slave to its Master

Timing diagram



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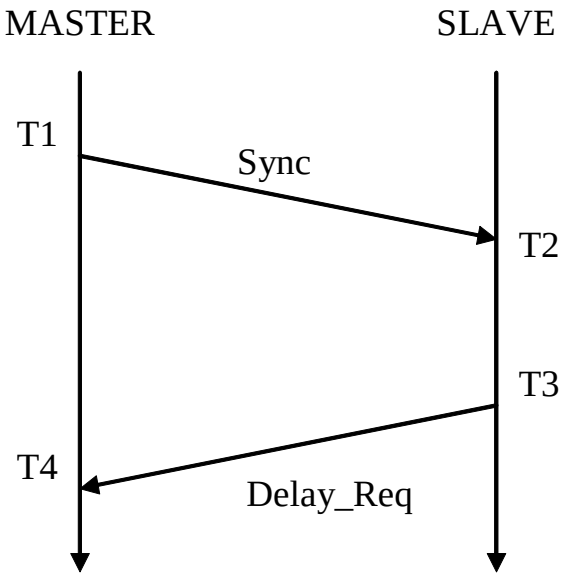
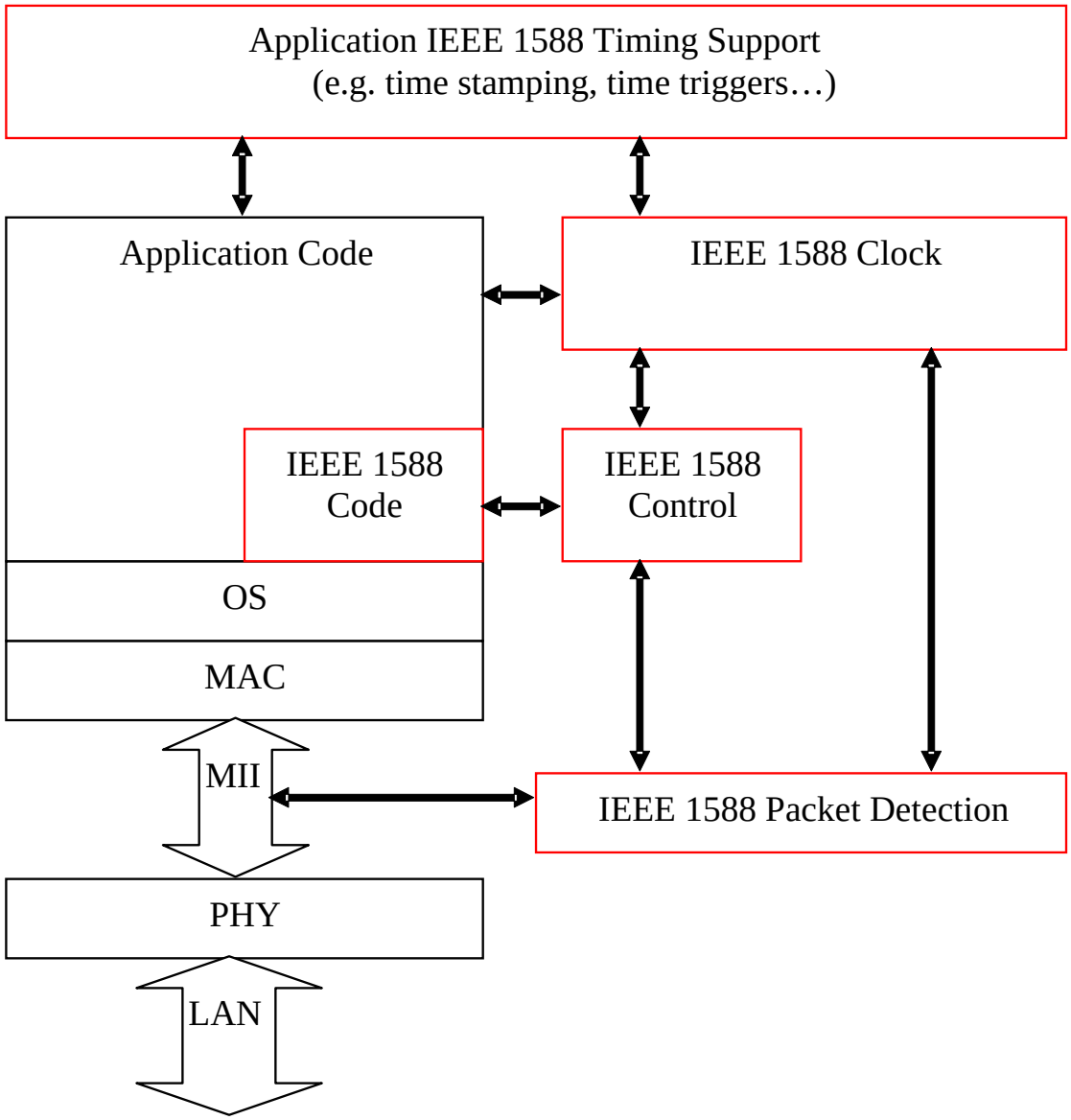
IEEE 1588





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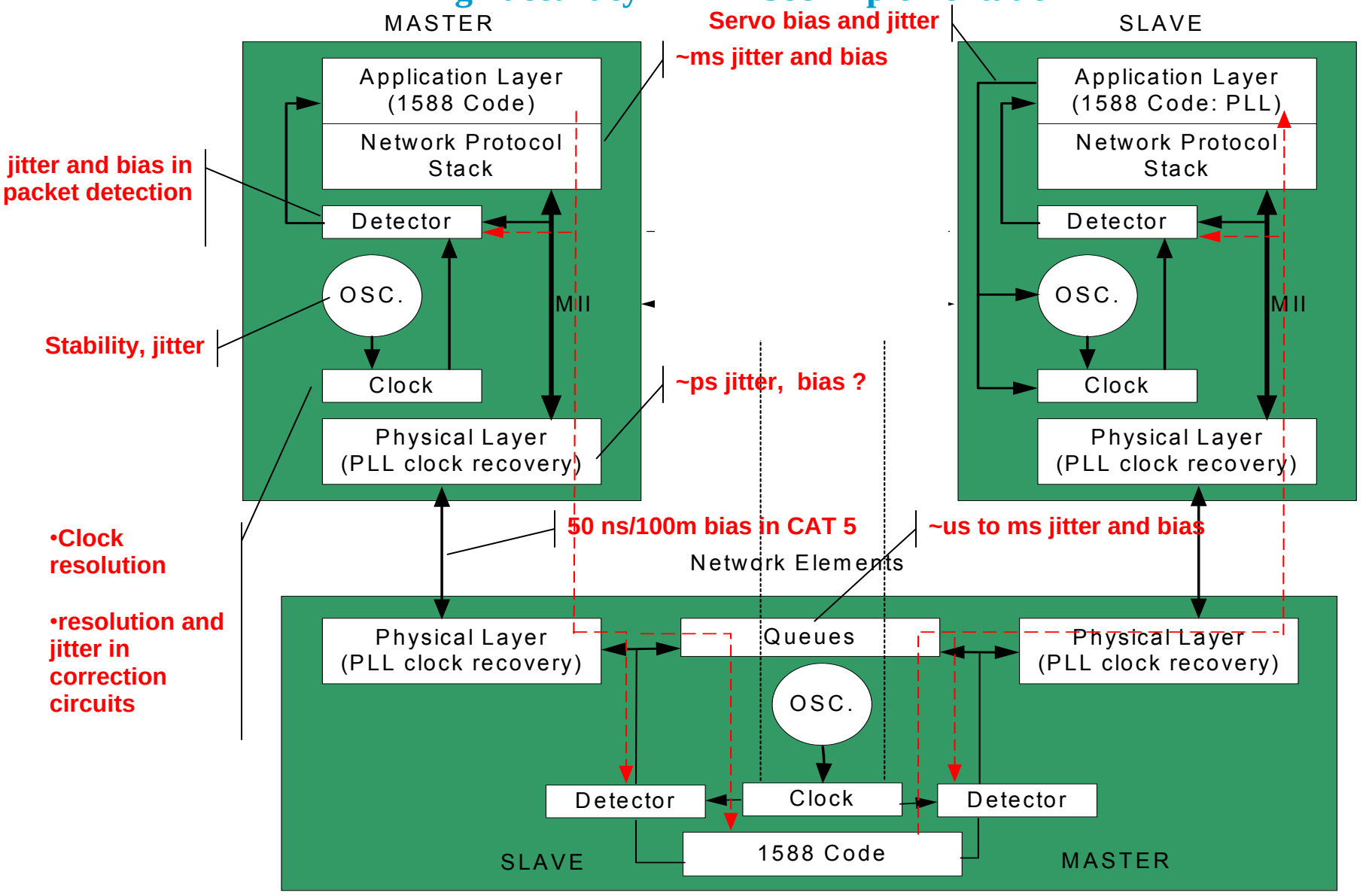
High accuracy IEEE 1588 implementation



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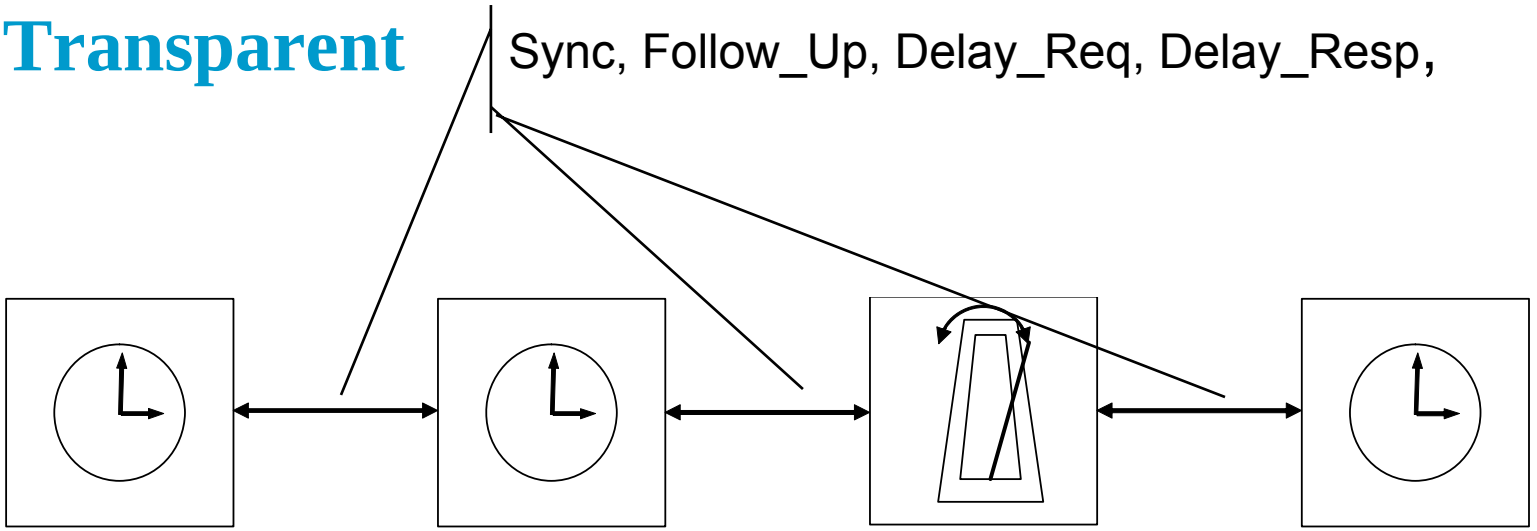
V2 Synchronization Basics + End-to-end Transparent



Grandmaster Clock
This clock determines the time base for the system

Slave to the Grandmaster Clock and Master to its Slave

End-to-end Transparent Clock (syntonizes to master)

Slave to its Master

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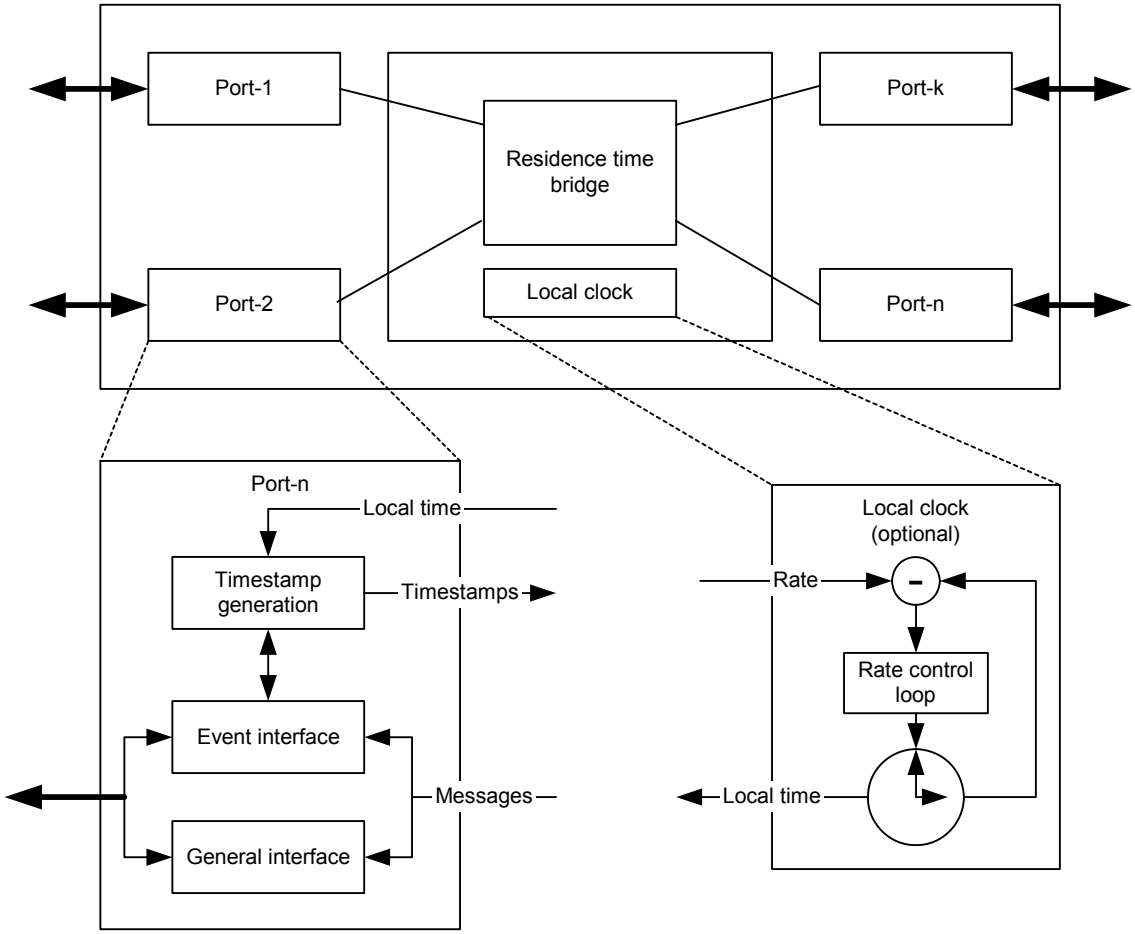
V2 End-to-end Transparent Clock



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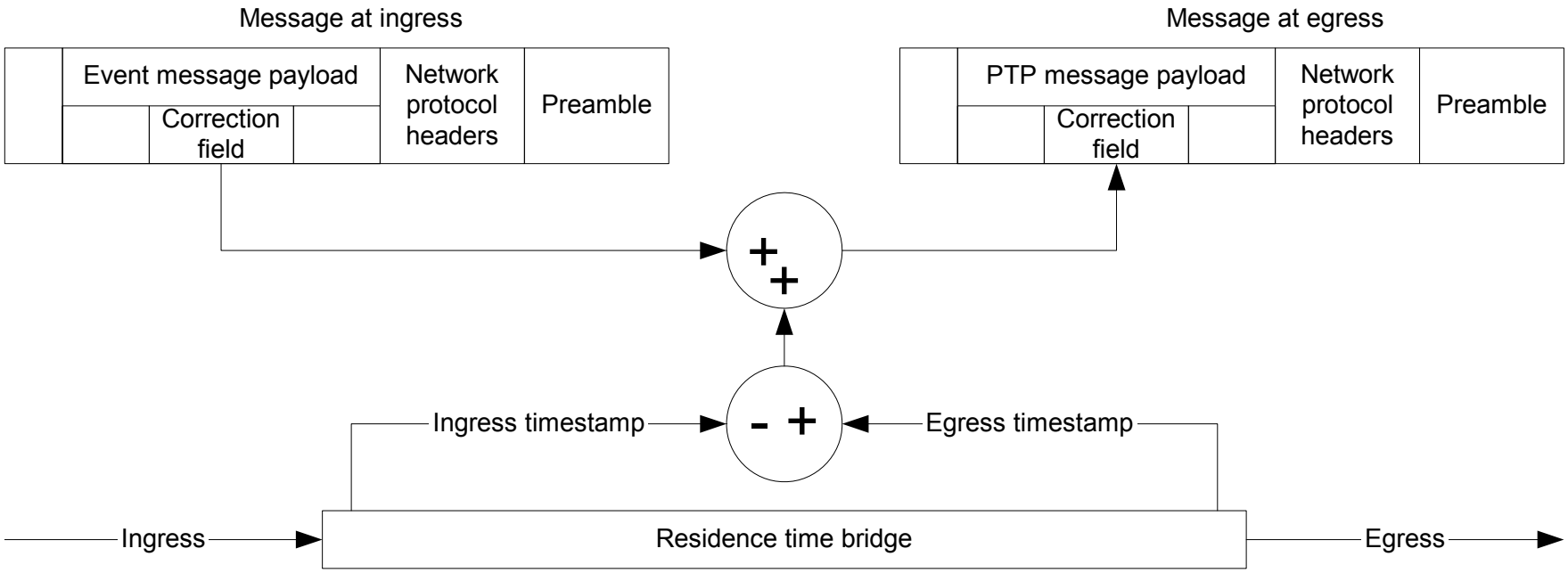
V2 End-to-end Transparent Clock Corrections



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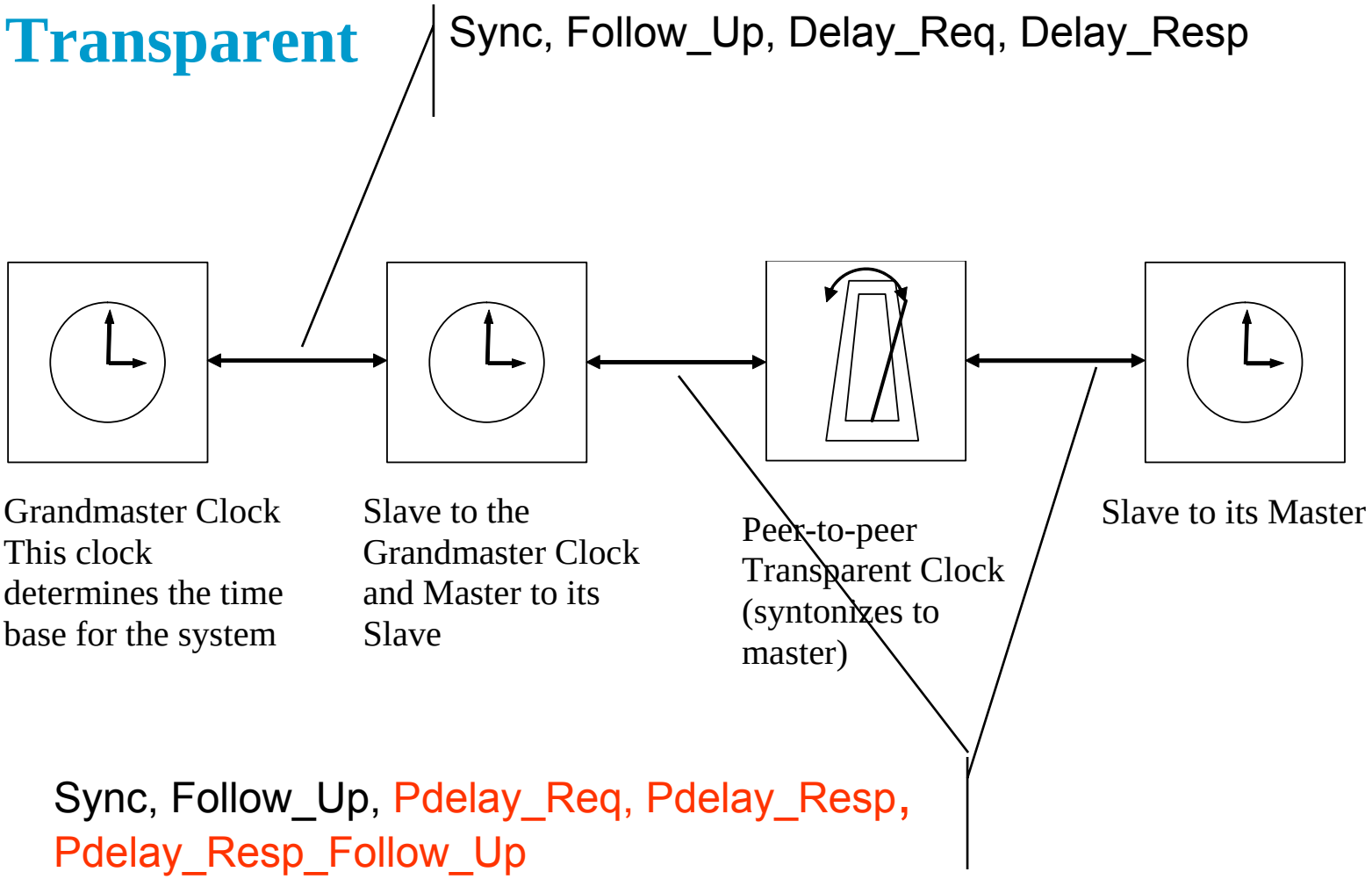
V2 Synchronization Basics + Peer-to-peer Transparent



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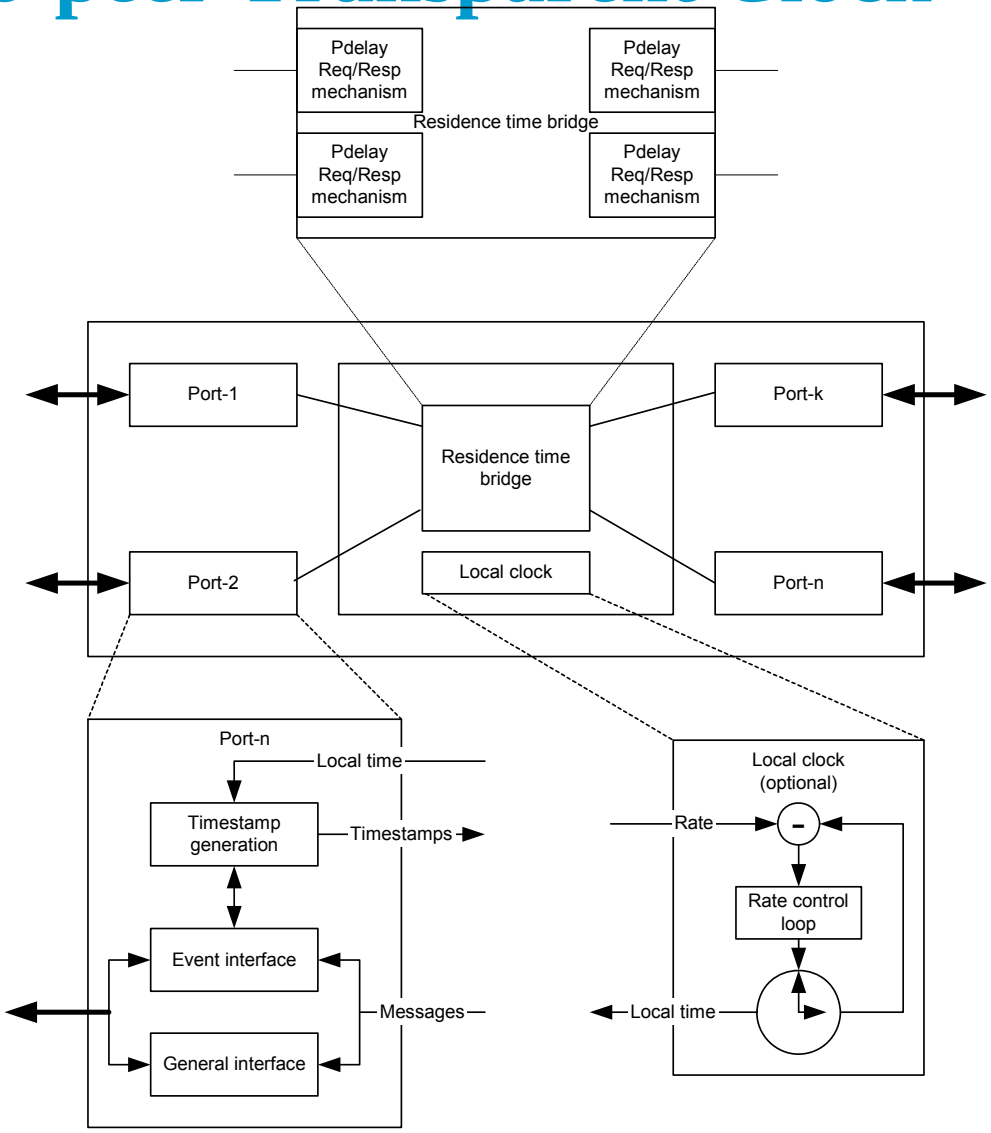
V2 Peer-to-peer Transparent Clock



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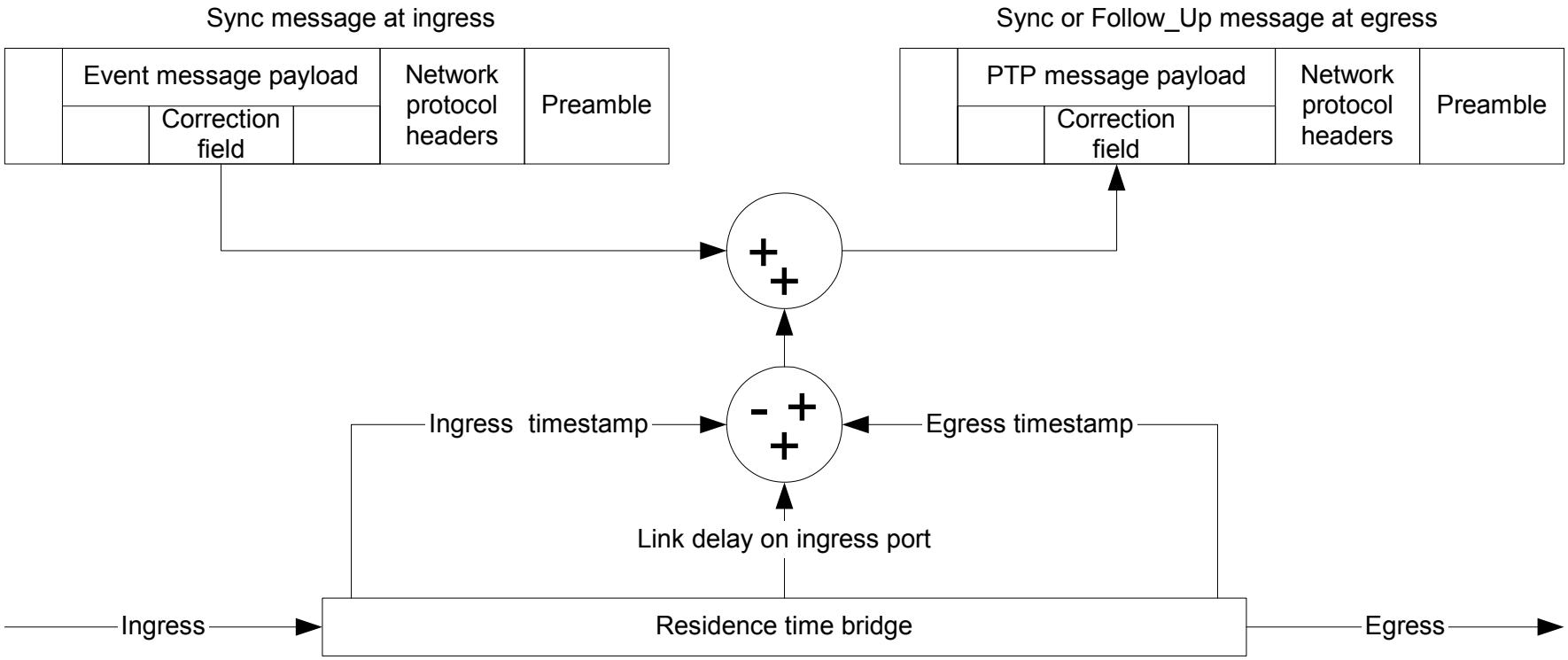
V2 Peer-to-peer Transparent Clock Corrections



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V2 1588 Switches

1. **Boundary clocks**
   a) + Good for hierarchical systems
   b) + Scale well with the number of devices
   c) – Poor for linear systems (large number of daisy chained clocks)
   d) + Can resolve 1:N introduced by ordinary switches or end-to-end TC
   e) – Must maintain state for scaling and 1:N
   f) + Can translate between different media.
2. **End-to-end transparent clocks**
   a) + Can be used for hierarchical systems
   b) – Scale poorly with the number of devices (master sees all slaves)
   c) + Good for linear systems (eliminates cascaded servos)
   d) – Can introduce 1:N topology

V2 1588 Switches (con't)

3. Peer-to-peer transparent clocks

a) + Can be used for hierarchical systems

b) + Scale well with the number of devices

c) + Good for linear systems (large number of daisy chained clocks)

d) + Rapid recovery with changes in network topology

e) – Cannot resolve 1:N introduced by ordinary switches, wireless, or end-to-end TC

f) – Must maintain per port path length state and measuring mechanisms.

g) – Only used in homogeneous P2P systems. Requires a boundary clock at the edges.

h) – 6 (vs. 4) measurements per link introduces slight degradation in accuracy

Split ‘timing’ and ‘master-slave hierarchy determination’

V1 Sync message (165 octets)

- **Network headers**
- **Version, subdomain, type**
- **Source identification**
- **Control, flags, update rates**
- **Origin timestamp**
- **Grandmaster information**
- **Local clock information**
- **Parent clock information**

V2 Sync message (46 octets)

- **Network headers**
- **PTP common: type, version, domain, CORRECTION FIELD, source identification, update rates, *control***
- **Origin timestamp**

Split 'timing' and 'master-slave hierarchy determination'

V1 Sync message (165 octets)

- **Network headers**
- **Version, subdomain, type**
- **Source identification**
- **Control, flags, update rates**
- **Origin timestamp**
- **Grandmaster information**
- **Local clock information**
- **Parent clock information**

V2 Announce message (88 octets)

- **Network headers**
- **PTP common: type, version, domain, CORRECTION FIELD, source identification, update rates, *control***
- ***Origin timestamp***
- **Grandmaster information**
- **Local clock information**
- **Parent clock information**

Provision for high accuracy implementations & transparent clocks

Correction field:

- **Integer64**
- **Correction in nanoseconds x 2^{+16}**
- **e.g. 2.5 ns = 0x0000000000028000**

Other v2 Changes

- **Wider choice of ‘Sync’ (timing) and ‘Announce’ (M-S hierarchy) update rates**
- **Fixed v1 errors (and hopefully did not introduce too many new ones)**
- **TLV extension mechanism**
- **Profiles to accommodate different market requirements**
- **WIP: Fault tolerance, Unicast option, Security, Alternate timescales, SNMP MIB**
- **WIP: mappings to-L2, Profinet, DeviceNet**

Status

- **Original ballot target fall 2006**
- **Probable ballot spring 2007**