### **Dynamics and Equilibria**

#### **Sergiu Hart**

Presidential Address, GAMES 2008 (July 2008)

Revised and Expanded (November 2009)

Revised (2010, 2011, 2012, 2013)

#### DYNAMICS AND EQUILIBRIA

# Sergiu Hart

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# **Papers**

### **Papers**

- Hart and Mas-Colell, Econometrica 2000
- Hart and Mas-Colell, J Econ Theory 2001
- Hart and Mas-Colell, Amer Econ Rev 2003
- Hart, Econometrica 2005
- Hart and Mas-Colell, Games Econ Behav 2006
- Hart and Mansour, Games Econ Behav 2010
- Hart, Games Econ Behav 2011

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#### Book

"A fundamental issue with any concept of equilibrium, including Nash and correlated equilibria, is to define the process by which equilibrium is attained. The work of Professors Hart and Mas-Colell has been the deepest in this area, especially in defining conditions ("uncoupled dynamics") which reflect naturally the information available in real economic interactions. Their body of results is essential to study of these fundamental problems."

Kenneth J. Arrow, Stanford University, USA

"In social as well as physical systems, equilibrium is of fundamental importance. Reaching equilibrium is at least as important as being there. In the last quarter century, research that investigates how social or game-theoretic equilibrium is reached has been spearheaded by Sergiu Hart and Andreu Mas-Colell. The most outstanding works in this area are gathered in the book before us—a must for anyone interested in this dynamic area of emerging economic research."

Robert J. Aumann, Hebrew University of Jerusalem, Israel

"The question of learning and convergence to equilibrium is of critical importance to the foundations and applications of game theory. But after half a century of research there are no universally accepted answers: different assumptions about players' information and learning dynamics lead to different conclusions. The Hart and Mas-Colell book describes fascinating directions of research on this subject developed by two distinguished authors and their collaborators over the last dozen years."

Ehud Kalai, Northwestern University, USA

"In this collection two leading game theorists show that various forms of equilibrium can be learned by simple and natural learning strategies that put minimal demands on the players' knowledge and level of rationality. It represents a major contribution to one of the most important topics in modern game theory."

Peyton Young, Oxford University, UK

This volume collects almost two decades of joint work of Sergiu Hart and Andreu Mas-Colell on game dynamics and equilibria. The starting point was the introduction of the adaptive strategy called *regret-matching*, which on the one hand is simple and natural, and on the other is shown to lead to correlated equilibria. This initial finding—boundedly rational behavior that yields fully rational outcomes in the long run—generated a large body of work on the dynamics of simple adaptive strategies. In particular, a natural condition on dynamics was identified: *uncoupledness*, whereby decision-makers do not know each other's payoffs and utilities (so, while chosen actions may be observable, the motivations are not). This condition turns out to severely limit the equilibria that can be reached. Interestingly, there are connections to the behavioral and neurobiological sciences and also to computer science and engineering (e.g., via notions of "regret").

Simple Adaptive Strategies is self-contained and unified in its presentation. Together with the formal treatment of concepts, theorems, and proofs, significant space is devoted to informal explanations and illuminating examples. It may be used for advanced graduate courses—in game theory, economics, mathematics, computer science, engineering—and for further research.

#### World Scientific

www.worldscientific.com 8408 hc



Vol. 4 World Scientific Series in Economic Theory - Vol. 4 SIMPLE ADAPTIVE STRATEGIE SIMPLE ADAPTIVE **STRATEGIES** From Regret-Matching to **Uncoupled Dynamics** Sergiu Hart Andreu Mas-Colell Mas-Colell **//e** World Scientific

John Nash, Ph.D. Dissertation, Princeton 1950

#### **EQUILIBRIUM POINT:**

John Nash, Ph.D. Dissertation, Princeton 1950

#### **EQUILIBRIUM POINT:**

"Each player's strategy is optimal against those of the others."

John Nash, Ph.D. Dissertation, Princeton 1950

#### **FACT**

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There are no general, natural dynamics leading to Nash equilibrium

"general"

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There are no general, natural dynamics leading to Nash equilibrium

"general": in all games

#### **FACT**

There are no general, natural dynamics leading to Nash equilibrium

"general": in all games

rather than: in specific classes of games

#### **FACT**

- "general": in all games rather than: in specific classes of games:
  - two-person zero-sum games
  - two-person potential games
  - supermodular games
  - **.** . . .

#### **FACT**

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There are no general, natural dynamics leading to Nash equilibrium

"leading to Nash equilibrium"

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There are no general, natural dynamics leading to Nash equilibrium

"leading to Nash equilibrium": at a Nash equilibrium (or close to it) from some time on

#### **FACT**

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"natural"

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"natural":

#### **FACT**

- "natural":
  - adaptive (reacting, improving, ...)

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- "natural":
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    - computation (performed at each step)

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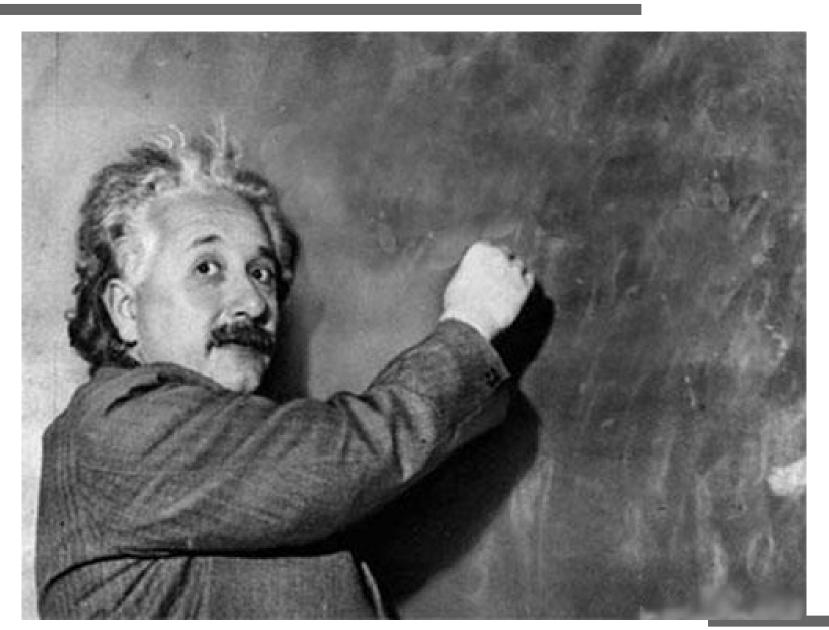
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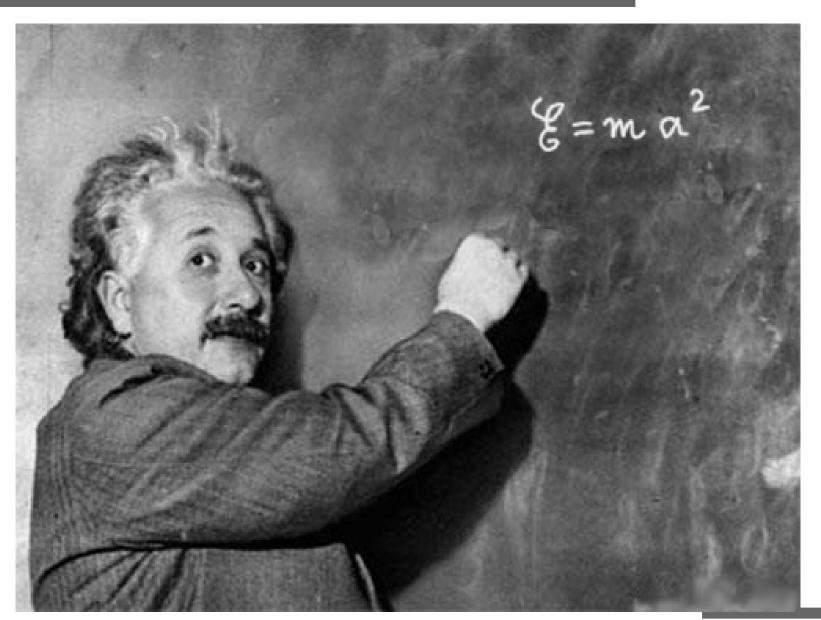
bounded rationality

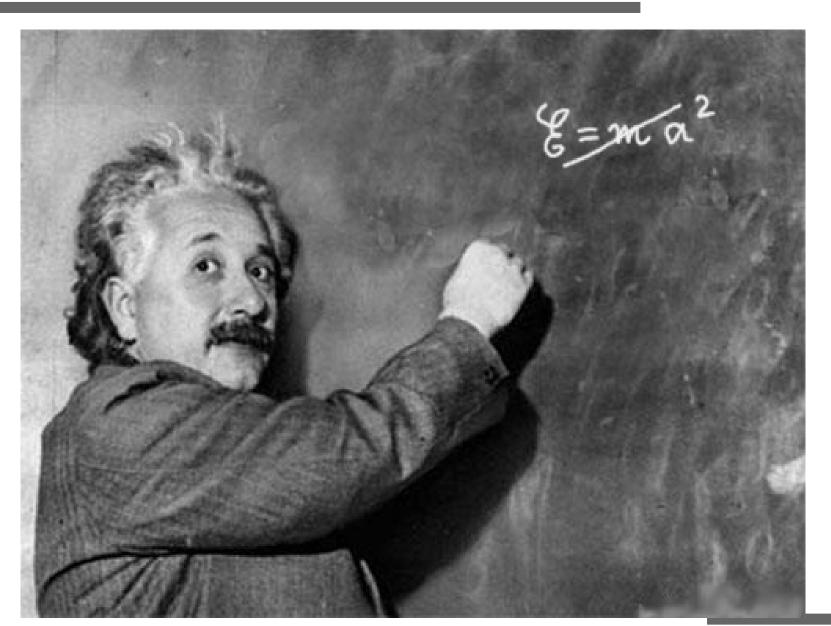
Dynamics that are **NOT** "natural":

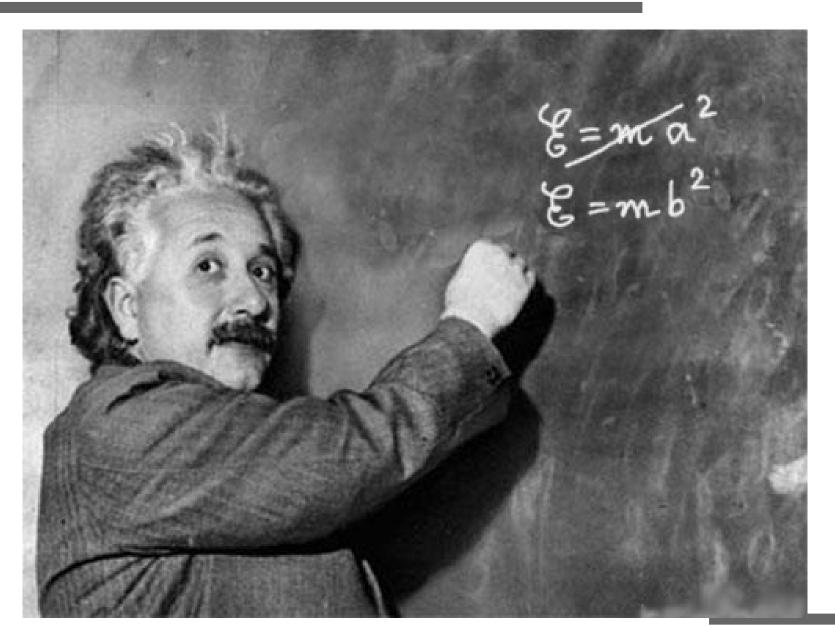
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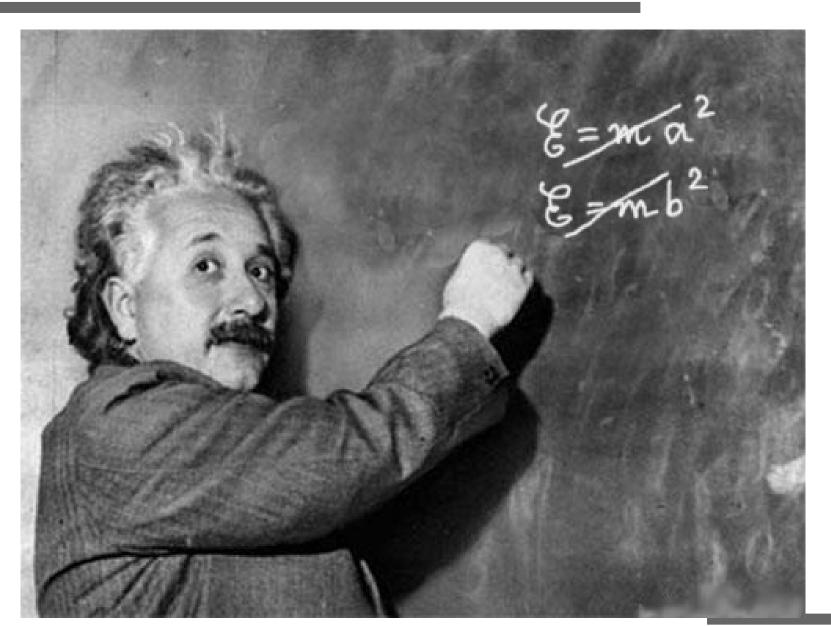
 exhaustive search (deterministic or stochastic)

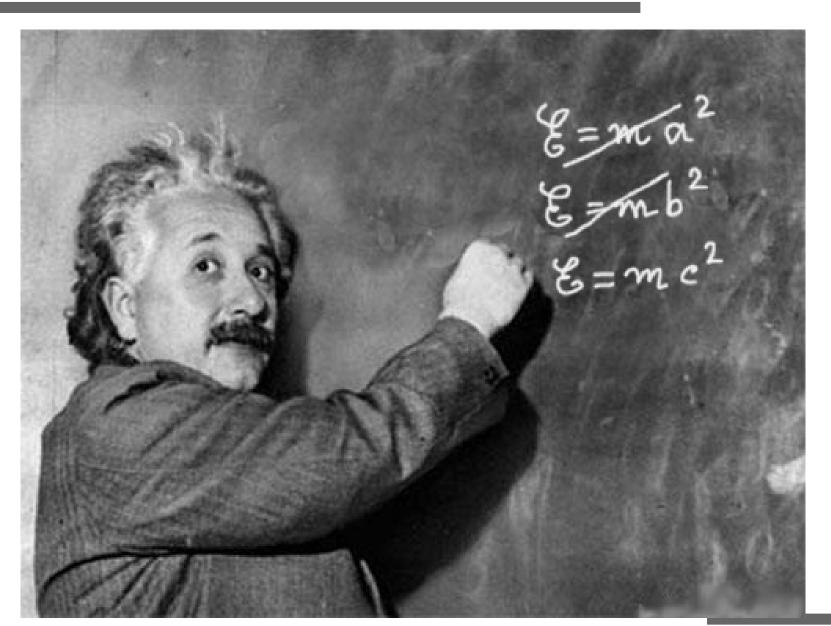


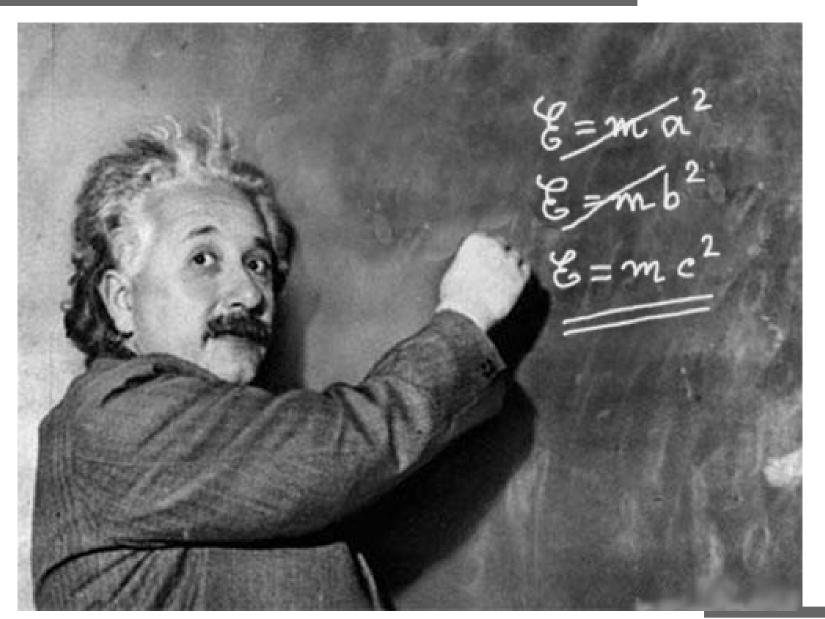


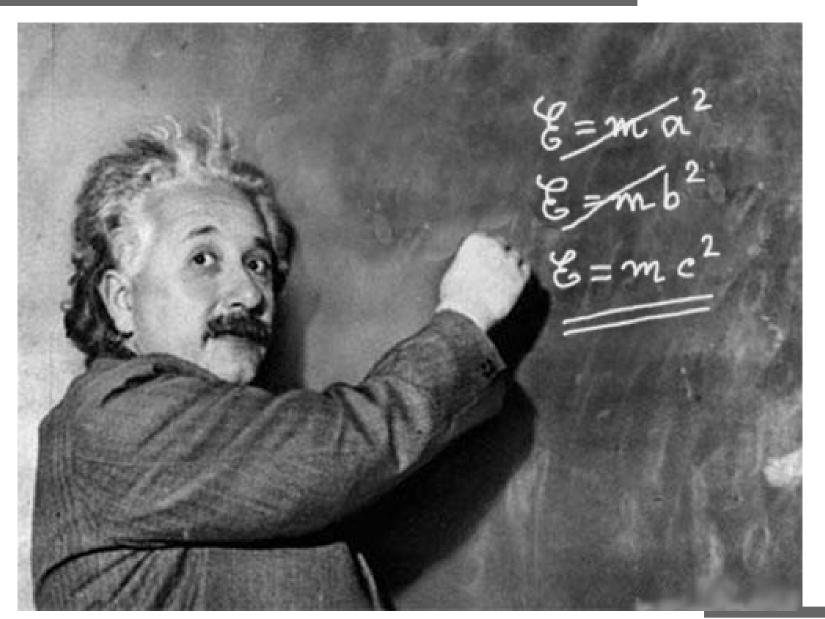












# Einstein's Manuscript

Der Ausdruck unter der Klammer rechts spielt die Rolle Extes bewegten Karsenpunktes) 2 = me 240 (28) wheder h- 22 = b geretzt ist und & eine andere alar fur den ersten Wellering urbezug auf E Lie 1/1-2

Albert Einstein, 1912
On the Special Theory of Relativity (manuscript)

Dynamics that are **NOT** "natural":

 exhaustive search (deterministic or stochastic)

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- using a mediator
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- fully rational learning
   (prior beliefs on the strategies of the opponents, Bayesian updating, optimization)

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## **UNCOUPLED DYNAMICS:**

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Hart and Mas-Colell, AER 2003

## **UNCOUPLED DYNAMICS:**

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(does *not* know the payoff functions of the other players)

(privacy-preserving, decentralized, distributed ...)

Hart and Mas-Colell, AER 2003

## Games

### N-person game in strategic (normal) form:

Players

$$i=1,2,...,N$$

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For each player i: Actions

$$a^i$$
 in  $A^i$ 

For each player i: Payoffs (utilities)

$$\mathbf{u}^{i}(a) \equiv \mathbf{u}^{i}(a^{1}, a^{2}, ..., a^{N})$$

#### Time

$$t = 1, 2, ...$$

Time

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ullet At period t each player i chooses an action  $a_t^i$  in  $A^i$ 

Time

$$t = 1, 2, ...$$

At period t each player i chooses an action

$$oldsymbol{a_t^i}$$
 in  $A^i$ 

according to a probability distribution

$$oldsymbol{\sigma_t^i}$$
 in  $\Delta(A^i)$ 

Fix the set of players 1, 2, ..., N and their action spaces  $A^1, A^2, ..., A^N$ 

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A general dynamic:

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$$\sigma_t^i \equiv \sigma_t^i$$
 (HISTORY; GAME)

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$$egin{aligned} \sigma_t^i &\equiv \sigma_t^i \ ( ext{ HISTORY} \; ; \; ext{GAME} \ ) \ &\equiv \sigma_t^i \ ( ext{ HISTORY} \; ; \; u^1,...,u^i,...,u^N \ ) \end{aligned}$$

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- Only last period matters ("1-recall")
- Time t does not matter ("stationary")

# **Impossibility**

## **Impossibility**

**Theorem.** There are **NO** uncoupled dynamics with 1-recall

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that yield almost sure convergence of play to pure Nash equilibria of the stage game in all games where such equilibria exist.

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Consider the following two-person game, which has a unique pure Nash equilibrium

	C1	C2	C3
R1	1,0	0,1	1,0
R2	0,1	1,0	1,0
R3	0,1	0,1	1,1

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Assume *by way of contradiction* that we are given an uncoupled, 1-recall, stationary dynamic that yields almost sure convergence to pure Nash equilibria when these exist

• Suppose the play at time t-1 is (R1,C1)

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  - Change the payoff function of COLIN so that (R1,C1) is the unique pure Nash eq.

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  - By uncoupledness, the same holds in the original game

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### Similarly for COLIN:

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R1	1,0 ↔	0,1	1,0
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R2	0,1 🙏	1,0 ↔	1,0 ↔
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### Similarly for COLIN:

### A player who is best replying cannot switch

	C1	C2	C3
R1	1,0 ↔	0,1 🙏	1,0 ↔
R2	0,1 🙏	1,0 ↔	1,0 ↔
R3	0,1 🙏	0,1 🙏	1,1

⇒ (R3,C3) cannot be reached

### Similarly for COLIN:

	C1	C2	C3
R1	1,0 ↔	0,1 🚺	1,0 ↔
R2	0,1 🙏	1,0 ↔	1,0 ↔
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## 2-Recall

## 2-Recall: Possibility

**Theorem.** THERE EXIST uncoupled dynamics with 2-RECALL

$$\sigma_t^i \equiv f^i(a_{t-2},a_{t-1};u^i)$$

that yield almost sure convergence of play to pure Nash equilibria of the stage game in every game where such equilibria exist.

Define the strategy of each player i as follows:

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- Everyone played the same in the previous two periods:  $a_{t-2} = a_{t-1} = a$ ; and
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THEN: At t player i plays  $a^i$  again:  $a^i_t = a^i$ 

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**ELSE**: At t player i randomizes uniformly over  $A^i$ 

"Good":

#### "Good":

simple

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simple

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simple

#### "Bad":

exhaustive search

#### "Good":

simple

- exhaustive search
- all players must use it

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```
"Ugly": ...
```

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### **FACT**

- "natural":
  - adaptive
  - simple and efficient:
    - computation
    - time
    - information

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### **FACT**

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    - **computation**: finite recall √
    - time to reach equilibrium ?
    - information: uncoupledness √

# Natural Dynamics: Time

#### HOW LONG TO EQUILIBRIUM?

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Estimate the number of time periods it takes until a Nash equilibrium is reached

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- How?
- An uncoupled dynamic



A distributed computational procedure

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 $\approx$ 

A distributed computational procedure

■ ⇒ COMMUNICATION COMPLEXITY

Distributed computational procedure

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  - START: Each participant has some private information

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  - communication: Messages are transmitted between the participants

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Yao 1979, Kushilevitz and Nisan 1997

Uncoupled dynamic leading to Nash equilibria

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Conitzer and Sandholm 2004

An uncoupled dynamic leading to Nash equilibria is TIME-EFFICIENT if

An uncoupled dynamic leading to Nash equilibria is TIME-EFFICIENT if the TIME IT TAKES

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In fact: exponential, like exhaustive search

Hart and Mansour, GEB 2010

Intuition:

- Intuition:
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- the dynamic procedure must distinguish between them

#### Intuition:

- different games have different equilibria
- the dynamic procedure must distinguish between them
- no single player can do so by himself

#### **FACT**

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#### **RESULT**

There CANNOT BE general, natural dynamics leading to Nash equilibrium

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Perhaps we are asking too much?

## **Dynamics and Nash Equilibrium**

#### **RESULT**

There CANNOT BE general, natural dynamics leading to Nash equilibrium

- Perhaps we are asking too much?
- For instance, the size of the data (the payoff functions) is exponential rather than polynomial in the number of players

#### **CORRELATED EQUILIBRIUM**

Aumann, JME 1974

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## Nash equilibrium when players receive payoff-irrelevant information before playing the game

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#### "Chicken" game

LEAVE STAY

LEAVE

STAY

5,5	3,6
6,3	0,0

#### "Chicken" game

	LEAVE	STAY
LEAVE	5,5	3,6
STAY	6,3	0,0

a Nash equilibrium

#### "Chicken" game

LEAVE STAY

**LEAVE** 

STAY

5, 5	3,6
6, 3	0,0

another Nash equilibrium

#### "Chicken" game

LEAVE STAY

LEAVE | 5

STAY

5, 5	3,6
6,3	0,0

0	1/2
1/2	0

a (publicly) correlated equilibrium

#### "Chicken" game

	LEAVE	STAY
LEAVE	5,5	3,6
STAY	6,3	0,0

L 
$$1/3$$
  $1/3$  S  $1/3$  O

#### another correlated equilibrium

- after signal L play LEAVE
- after signal s play STAY

A Correlated Equilibrium is a Nash equilibrium when the players receive payoff-irrelevant signals before playing the game (Aumann 1974)

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- Boston Celtics' front line

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- Common Knowledge of Rationality ⇔ Correlated Equilibrium (Aumann 1987)

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A joint distribution z is a correlated equilibrium

$$\Leftrightarrow$$

$$\sum_{s^{-i}} u(j, s^{-i}) z(j, s^{-i}) \geq \sum_{s^{-i}} u(k, s^{-i}) z(j, s^{-i})$$

for all  $i \in N$  and all  $j,k \in S^i$ 

#### **RESULT**

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THERE EXIST general, natural dynamics leading to CORRELATED EQUILIBRIA

Regret Matching

Hart and Mas-Colell, Ec'ca 2000

#### **RESULT**

THERE EXIST general, natural dynamics leading to CORRELATED EQUILIBRIA

- Regret Matching
- General regret-based dynamics

Hart and Mas-Colell, Ec'ca 2000, JET 2001

## Regret Matching

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"REGRET": the increase in past payoff, if any, if a different action would have been used

"MATCHING": switching to a different action with a probability that is proportional to the regret for that action

THERE EXIST general, natural dynamics leading to CORRELATED EQUILIBRIA

"general": in all games

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  - adaptive (also: close to "behavioral")
  - simple and efficient: computation, time, information
- "leading to correlated equilibria": statistics of play become close to CORRELATED EQUILIBRIA

## Regret Matching and Beyond

## Regret Matching and Beyond

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# The Involvement of the Orbitofrontal Cortex in the Experience of Regret

Nathalie Camille,<sup>1\*</sup> Giorgio Coricelli,<sup>1,2\*</sup> Jerome Sallet,<sup>1</sup> Pascale Pradat-Diehl,<sup>3</sup> Jean-René Duhamel,<sup>1</sup> Angela Sirigu<sup>1</sup>†

Facing the consequence of a decision we made can trigger emotions like satisfaction, relief, or regret, which reflect our assessment of what was gained as compared to what would have been gained by making a different decision. These emotions are mediated by a cognitive process known as counterfactual thinking. By manipulating a simple gambling task, we characterized a subject's choices in terms of their anticipated and actual emotional impact. Normal subjects reported emotional responses consistent with counterfactual thinking; they chose to minimize future regret and learned from their emotional experience. Patients with orbitofrontal cortical lesions, however, did not report regret or anticipate negative consequences of their choices. The orbitofrontal cortex has a fundamental role in mediating the experience of regret.



Review

TRENDS in Cognitive Sciences Vol.11 No.6



### Brain, emotion and decision making: the paradigmatic example of regret

Giorgio Coricelli<sup>1</sup>, Raymond J. Dolan<sup>2</sup> and Angela Sirigu<sup>1</sup>

Human decisions cannot be explained solely by rational imperatives but are strongly influenced by emotion. Theoretical and behavioral studies provide a sound empirical basis to the impact of the emotion of regret in guiding choice behavior. Recent neuropsychological and neuroimaging data have stressed the fundamental role of the orbitofrontal cortex in mediating the experience of regret. Functional magnetic resonance imaging data indicate that reactivation of activity within the orbitofrontal cortex and amygdala occurring during the phase of choice, when the brain is anticipating possible future consequences of decisions, characterizes the anticipation of regret. In turn, these patterns reflect learning based on cumulative emotional experience. Moreover, affective consequences can induce specific mechanisms of cognitive control of the choice processes, involving reinforcement or avoidance of the experienced behavior.

change. People, including those with a deep knowledge of optimal strategies, such as Markowitz, often try to avoid the likelihood of future regret, even when this conflicts with the prescription of decisions based on rational choice; according to the latter, individuals faced with a decision between multiple alternatives under uncertainty will opt for the course of action with maximum expected utility, a function of both the probability and the magnitude of the expected payoff [4].

Here, we outline, for the first time, the neural basis of the emotion of regret, and its fundamental role in adaptive behavior. The following questions will be addressed: what are the neural underpinnings of 'powerful' cognitively generated emotions such as regret? What are the theoretical implications of incorporating regret into the process of choice, and into adaptive models of decision making? In line with recent work on emotion-based decision making [5,6], we attempt to characterize the brain areas underlying

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# Decentralized Dynamic Spectrum Access for Cognitive Radios: Cooperative Design of a Non-Cooperative Game

Michael Maskery, Vikram Krishnamurthy, and Qing Zhao

Abstract—We consider dynamic spectrum access among cognitive radios from an adaptive, game theoretic learning perspective. Spectrum-agile cognitive radios compete for channels temporarily vacated by licensed primary users in order to satisfy their own demands while minimizing interference. For both slowly varying primary user activity and slowly varying statistics of "fast" primary user activity, we apply an adaptive regret based learning procedure which tracks the set of correlated equilibria of the game, treated as a distributed stochastic approximation. This procedure is shown to perform very well compared with other similar adaptive algorithms. We also estimate channel contention for a simple CSMA channel sharing scheme.

*Index Terms*—Cognitive radio, dynamic spectrum access, game theory, stochastic approximation, correlated equilibrium.

(channels) that are temporarily unoccupied by licensed users. Each radio dynamically selects several available channels so as to balance its own demand (competition) against system-imposed sharing incentives (cooperation). Selections are made independently by each radio, based only on its own performance history. We focus on applications where primary users' spectrum access activities either vary slowly with time (see [3], [4]), or where their spectrum access activities vary quickly, but average behaviour varies slowly. Example applications include the reuse of certain TC-bands that are not used for TC broadcast in a particular region.

Since optimal resource allocation in a decentralized, competitive environment is not straightforward, we propose to operate radios according to a game theoretic algorithm which

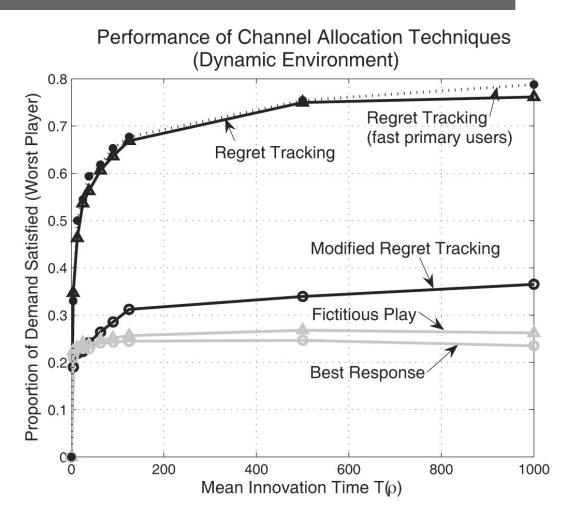


Fig. 6. Long-run average spectrum utilization in a dynamic environment for the channel allocation techniques of Section V.  $T(\rho)$  (see Sec.VI-C) is the mean time between innovations in the system (changes in primary user activity, fast primary user statistics or cognitive radio demands).

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#### Fully non-cooperative optimal placement of mobile vehicles

Shemin Kalam, Mahbub Gani and Lakmal Seneviratne

Abstract—In this paper, we consider optimal placement of autonomous mobile vehicles such that a cost function involving all the vehicles and possible locations of targets is minimized. This cost is proportional to the distance between the targets and vehicles. The optimal locations correspond to the vehicles being at the centroids of their own Voronoi cell which correspond to Centroidal Voronoi Tessellations (CVTs). We have adopted a game theoretical formulation to initially consider vehicle target assignment where a set of mobile vehicles choose their own targets. The movement of the vehicles towards the optimal locations is based on MacQueen's algorithm. But an important step of MacQueen's algorithm requires the knowledge of the nearest neighbour to be determined from a sample that is drawn from a fixed but unknown probability distribution. This calculation seems to be implicit in reported algorithms and brings in a hidden centralized process. We have used game theory as a framework to get around this problem and modelled the vehicles such that they are capable of making their own decisions and interested in optimizing their own utilities. Specifically, we have introduced an appropriate utility function and require the vehicles to negotiate their choice of targets via regret matching. We present simulations that illustrate that vehicles choose the targets optimally and converge to CVTs.

Lloyd's descent algorithm [15] can be applied to solve the problem and it has been shown by Cortés et al [12] that the distribution of mobile sensors converges to Centroidal Voronoi Tessellations (CVTs). Considering a scenario where the spatial distribution is not known, MacQueen's algorithm [16] is a Monte-Carlo method of solving the problem [22]; in our problem we interpret MacQueen's algorithm as a real-time higher order control strategy. A drawback of the MacQueen's algorithm is that it requires the calculation of nearest neighbours. We have drawn upon game theory to get around the problem of this requirement. This reduces the communication burden on the system because the vehicles do not require the information about the distances between all the vehicles and targets which would have been needed for calculation of the nearest neighbour.

The game theoretic strategy that we adopt is inspired by the problem of autonomous vehicle-target assignment tackled by Arslan in [8]. To get to the problem of optimal placement of the mobile vehicles, initially we consider how a group of vehicle are to optimally assign themselves to a set of targets.

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Network-Enabled Missile Deflection: Games and Correlation Equilibrium

MICHAEL MASKERY
VIKRAM KRISHNAMURTHY
University of British Columbia

The problem of deploying countermeasures (CM) against antiship missiles is investigated from a network centric perspective in which multiple ships coordinate to defend against a known missile threat. Using the paradigm of network enabled operations (NEOPS), the problem is formulated as a transient stochastic game with communication where the appropriate strategy takes the form of an optimal stationary correlated equilibrium. Under this strategy, ships cooperate through real-time communication to satisfy both local and collective interests. The use of communication results in a performance improvement over the noncommunicating, Nash equilibrium scenario. This framework allows us to develop a theoretical foundation for NEOPS and captures the trade-off between information exchange and performance, while generalizing the standard Nash equilibrium solution for the missile deflection game given in [1]. The NEOPS equilibrium strategy is characterized as the solution to an optimization problem with linear objective and bilinear constraints, which can be solved calculating successive improvements starting from an initial noncooperative (Nash) solution. The communication overhead required to implement this strategy is associated with the mutual information between individual action probability distributions at equilibrium. Numerical results illustrate the trade-off between communication and performance.

## Predicting Human Interactive Learning by Regret-Driven Neural Networks

Davide Marchiori<sup>1</sup> and Massimo Warglien<sup>2</sup>\*

Much of human learning in a social context has an interactive nature: What an individual learns is affected by what other individuals are learning at the same time. Games represent a widely accepted paradigm for representing interactive decision-making. We explored the potential value of neural networks for modeling and predicting human interactive learning in repeated games. We found that even very simple learning networks, driven by regret-based feedback, accurately predict observed human behavior in different experiments on 21 games with unique equilibria in mixed strategies. Introducing regret in the feedback dramatically improved the performance of the neural network. We show that regret-based models provide better predictions of learning than established economic models.

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#### Computers in Biology and Medicine





#### Classification of peptide mass fingerprint data by novel no-regret boosting method

Anna Gambin<sup>a,\*</sup>, Ewa Szczurek<sup>a,b</sup>, Janusz Dutkowski<sup>a</sup>, Magda Bakun<sup>c</sup>, Michał Dadlez<sup>c,d</sup>

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Boosting
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Ovarian cancer

#### ABSTRACT

We have developed an integrated tool for statistical analysis of large-scale LC-MS profiles of complex protein mixtures comprising a set of procedures for data processing, selection of biomarkers used in early diagnostic and classification of patients based on their peptide mass fingerprints.

Here, a novel boosting technique is proposed, which is embedded in our framework for MS data analysis. Our boosting scheme is based on Hannan-consistent game playing strategies. We analyze boosting from a game-theoretic perspective and define a new class of boosting algorithms called H-boosting methods. In the experimental part of this work we apply the new classifier together with classical and state-of-the-art algorithms to classify ovarian cancer and cystic fibrosis patients based on peptide mass spectra. The methods developed here provide automatic, general, and efficient means for processing of large scale LC-MS datasets. Good classification results suggest that our approach is able to uncover valuable information to support medical diagnosis.

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<sup>&</sup>lt;sup>c</sup>Institute of Biochemistry and Biophysics PAS, Pawińskiego 5A, 02-106 Warsaw, Poland

<sup>&</sup>lt;sup>d</sup>Biology Department, Warsaw University, Miecznikowa 1, 02-096 Warsaw, Poland

NASH EQUILIBRIUM: a fixed-point of a non-linear map

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- CORRELATED EQUILIBRIUM: a solution of finitely many linear inequalities

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set-valued fixed-point (curb sets)?

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- B. Find NATURAL dynamics for the various equilibrium concepts

#### Nash Equilibrium

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"DYNAMICALLY DIFFICULT"

Nash Equilibrium

"DYNAMICALLY DIFFICULT"

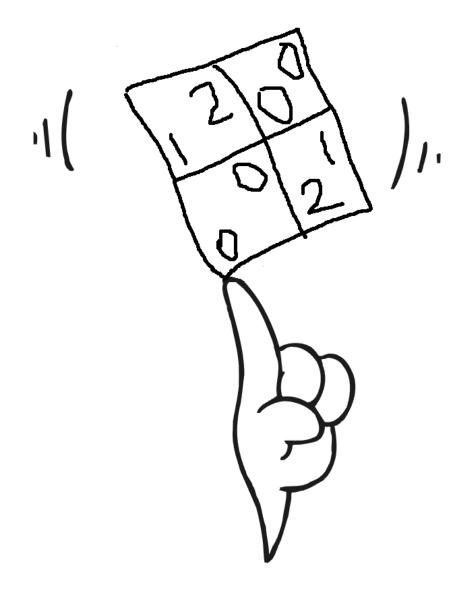
**Correlated Equilibrium** 

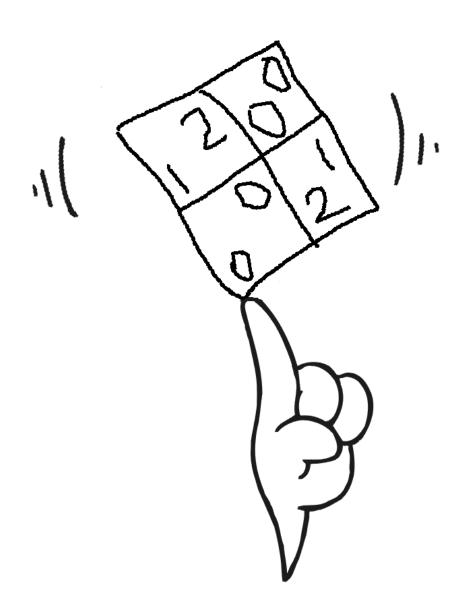
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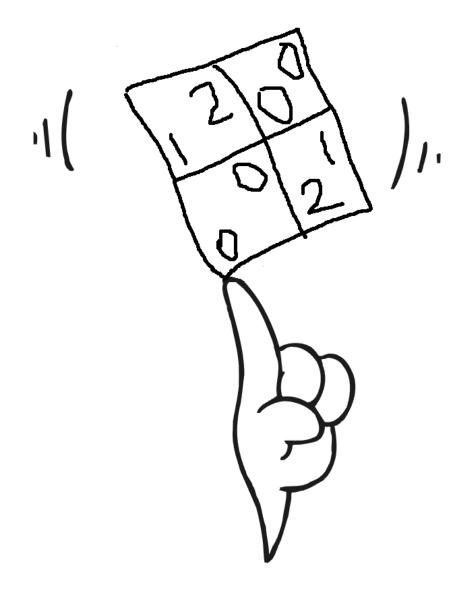
"DYNAMICALLY DIFFICULT"

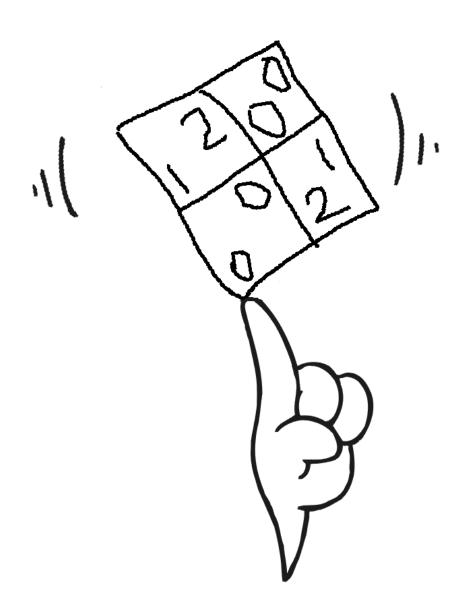
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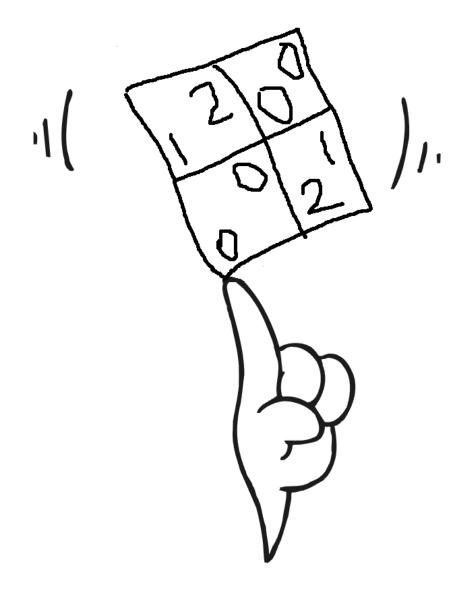
"DYNAMICALLY EASY"

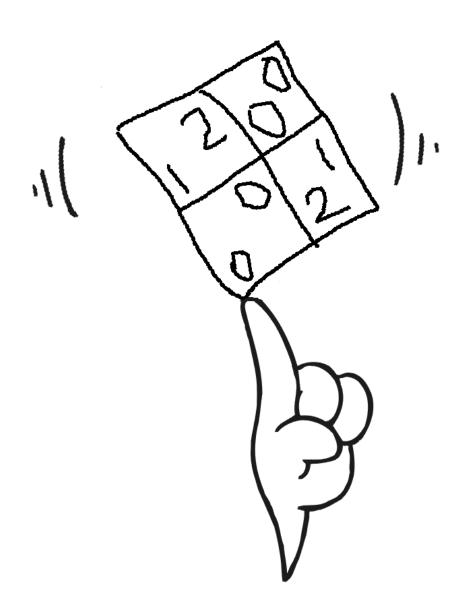


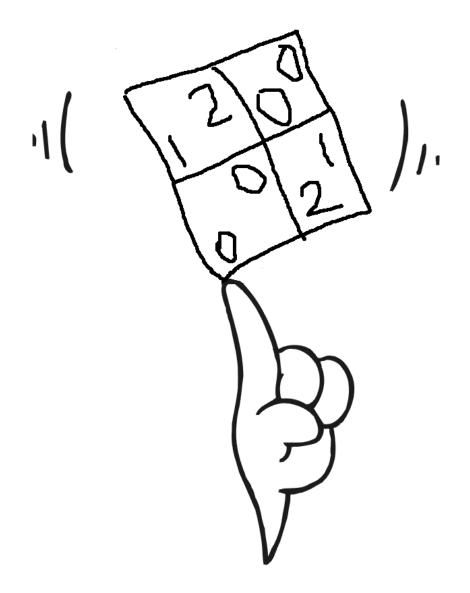


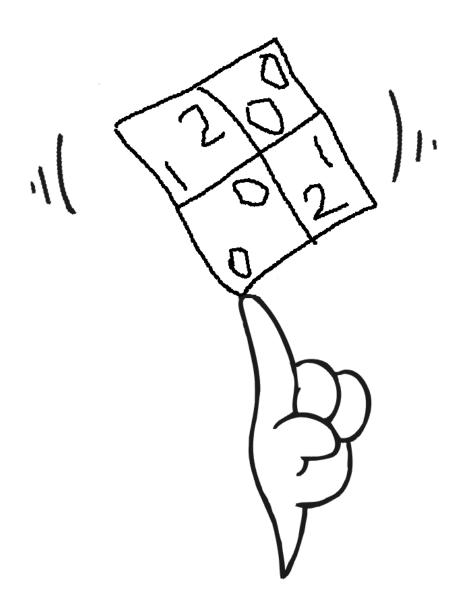


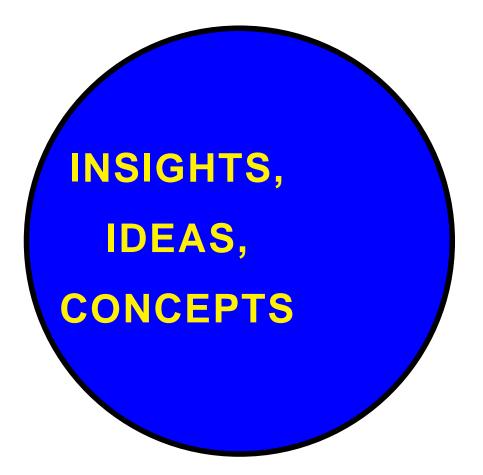


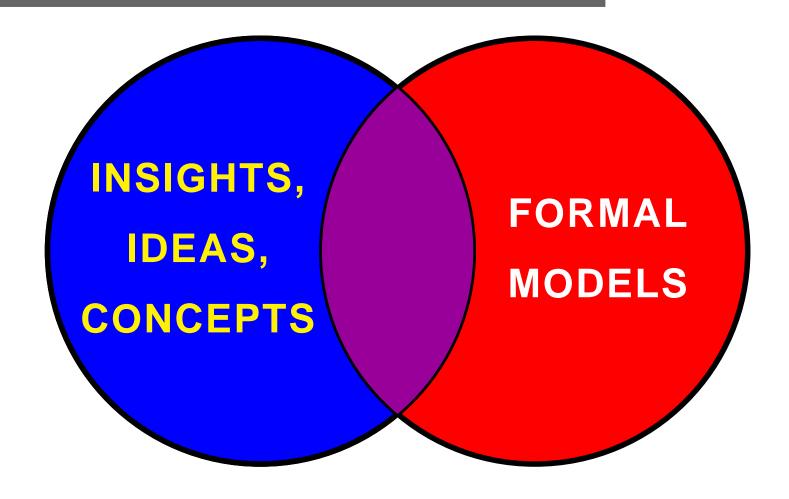












INSIGHTS,
IDEAS,
CONCEPTS

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