

Analysis of TSN for Industrial Automation based on Network Calculus

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Purpose

- Share a paper from our team with the group
 - <<Analysis of TSN for Industrial Automation based on Network Calculus>>
 - Network calculus theory, industrial automation network modeling, and simulation results.
 - <https://ieeexplore.ieee.org/document/8869053>
- Discuss the idea of using network calculus to calculate the **worst-case latency bound** for industrial automation scenarios.
 - Vital for using asynchronous/non-time-based methods, e.g., SP with CBS or ATS.
 - What is the challenge? Where is the gap?

Network calculus theory

- Traffic characteristics / traffic constraints (TSpec in TSN) → arrival curve → The bound.
- Device's capability (bandwidth, queuing and shaping, reservation) → service curve

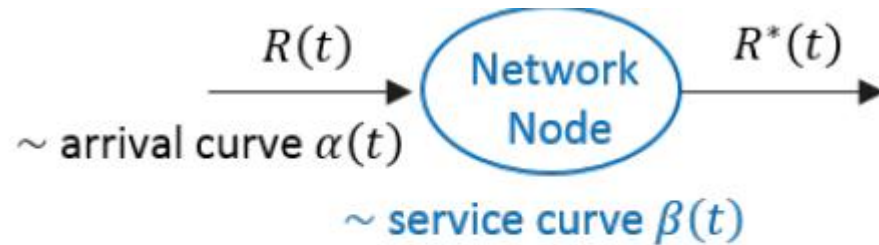


Fig. 3. Two-port network model of a TSN relay node

$$R(s+t) - R(s) \leq \alpha(t), \quad \forall s \geq 0, t \geq 0 \quad (1)$$

$$R^*(t) \geq R \otimes \beta(t) = \inf_s \{R(s) + \beta(t-s)\}, \quad \forall 0 \leq s \leq t \quad (2)$$

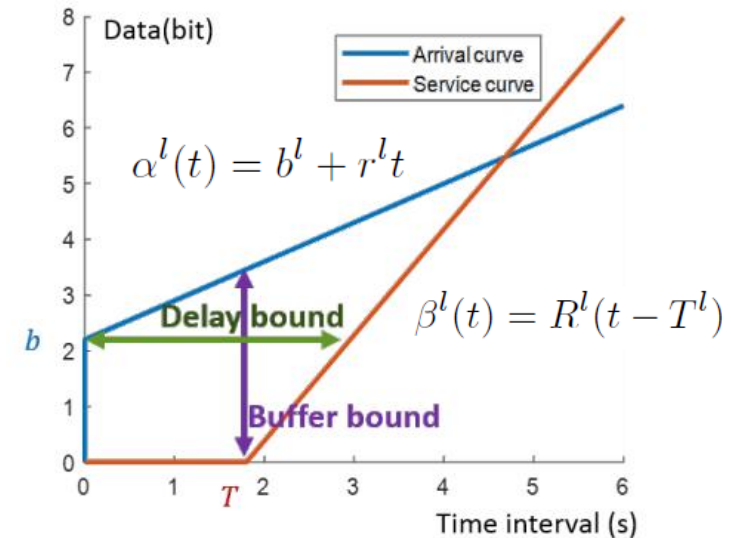


Fig. 4. Computation of backlog bound and delay bound.

Network calculus theory

CDT: Control Data Traffic

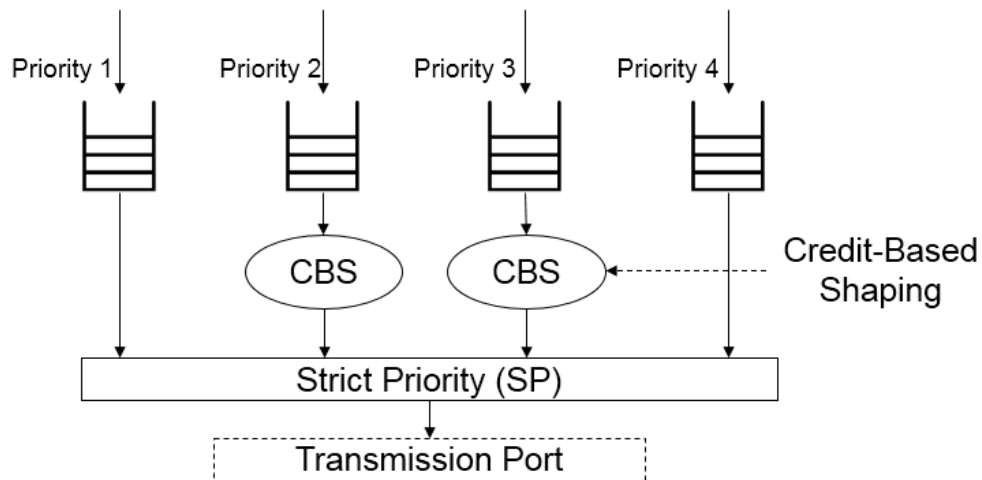
- The queuing and forwarding method: SP+CBS

TRAFFIC CLASS DEFINITION

	Priority	Traffic Class	Traffic constraint
highest →	1	CDT	
	2	SR Class A	Token bucket
	3	SR Class B	
lowest →	4	BE	

$$\alpha^l(t) = b^l + r^l t$$

$$\beta^l(t) = R^l(t - T^l)$$



For CDT,

$$R^{CDT} = c$$

$$T^{CDT} = \frac{b_h - b_f + L^{\bar{h}}}{c}$$

For SR Class A,

$$R^A = \frac{I^A(c - r_h)}{c}$$

$$T^A = \frac{L^{\bar{A}} + b_h + \frac{r_h L^{\bar{h}}}{c}}{c - r_h}$$

For SR Class B,

$$R^B = \frac{I^B(c - r_h)}{c}$$

$$T^B = \frac{L^{BE} + L^A + \frac{L^{\bar{A}} I^A}{c - I^A} + b_h + \frac{r_h L^{\bar{h}}}{c}}{c - r_h}$$

- One hop latency bound: $D_{i,j} = T^l + \frac{b^l - L_{\min,f}}{R^l} + \frac{L_{\min,f}}{c}$
- End-to-end latency bound: the sum of per-hop latency bound along the path.

Industrial automation network modeling

- Topology, flows, and shapers.

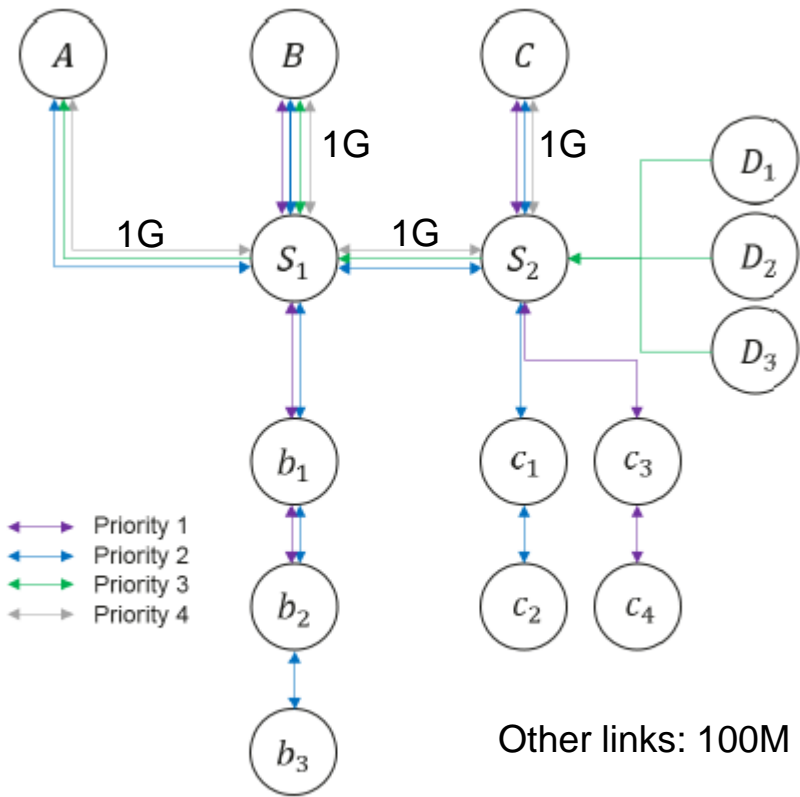


Fig. 5. Simulation topology

FLOW DESCRIPTION			
Flow path	Traffic Type	Forwarding	Priority
p1: B, S1, b1 p2: B, S1, b1, b2 p3: b1, S1, B p4: b2, b1, S1, B p5: C, S2, c3 p6: C, S2, c3, c4 p7: c3, S2, C p8: c4, c3, S2, C	Isochronous	CDT (SP)	1
p9: B, S1, S2, C p10: C, S2, S1, B p11: A, S1, B p12: B, S1, A p13: A, S1, S2, C p14: C, S2, S1, A p15: C, S2, c1 p16: C, S2, c1, c2 p17: c1, S2, C p18: c2, c1, S2, C p19: b3, b2, b1, S1, B p20: B, S1, b1, b2, b3	Cyclic	SR Class A (CBS)	2
p21: D1, S2, C p22: D2, S2, S1, A p23: D3, S2, S1, B	A/V	SR Class B (CBS)	3
p24: A, S1, B p25: B, S1, A p26: A, S1, S2, C p27: C, S2, S1, A	BE	BE (SP)	4

Industrial automation network modeling

- Topology, flows, and shapers.

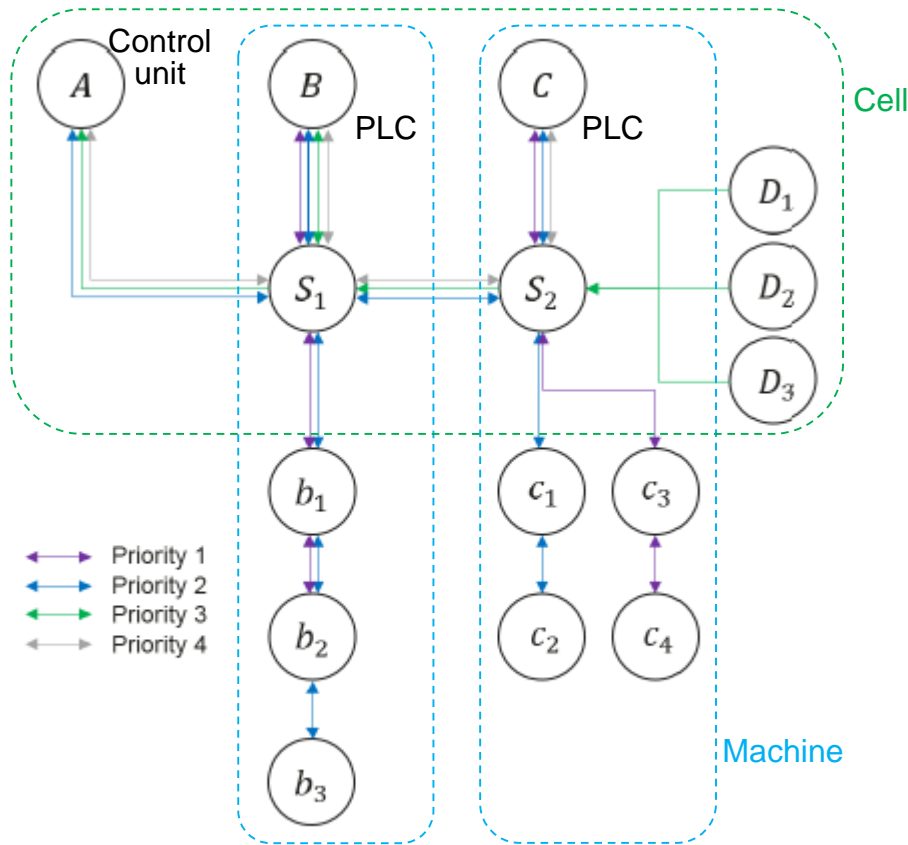


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p9: B, S1, S2, C p10: C, S2, S1, B p11: A, S1, B p12: B, S1, A p13: A, S1, S2, C p14: C, S2, S1, A p15: C, S2, c1 p16: C, S2, c1, c2 p17: c1, S2, C p18: c2, c1, S2, C p19: b3, b2, b1, S1, B p20: B, S1, b1, b2, b3						
p21: D1, S2, C p22: D2, S2, S1, A p23: D3, S2, S1, B				A/V	SR Class B (CBS)	3
p24: A, S1, B p25: B, S1, A p26: A, S1, S2, C p27: C, S2, S1, A				BE	BE (SP)	4

Lmax=0.8kb
T=2ms (cycle time)

Lmax=0.8kb
T=1ms (cycle time)

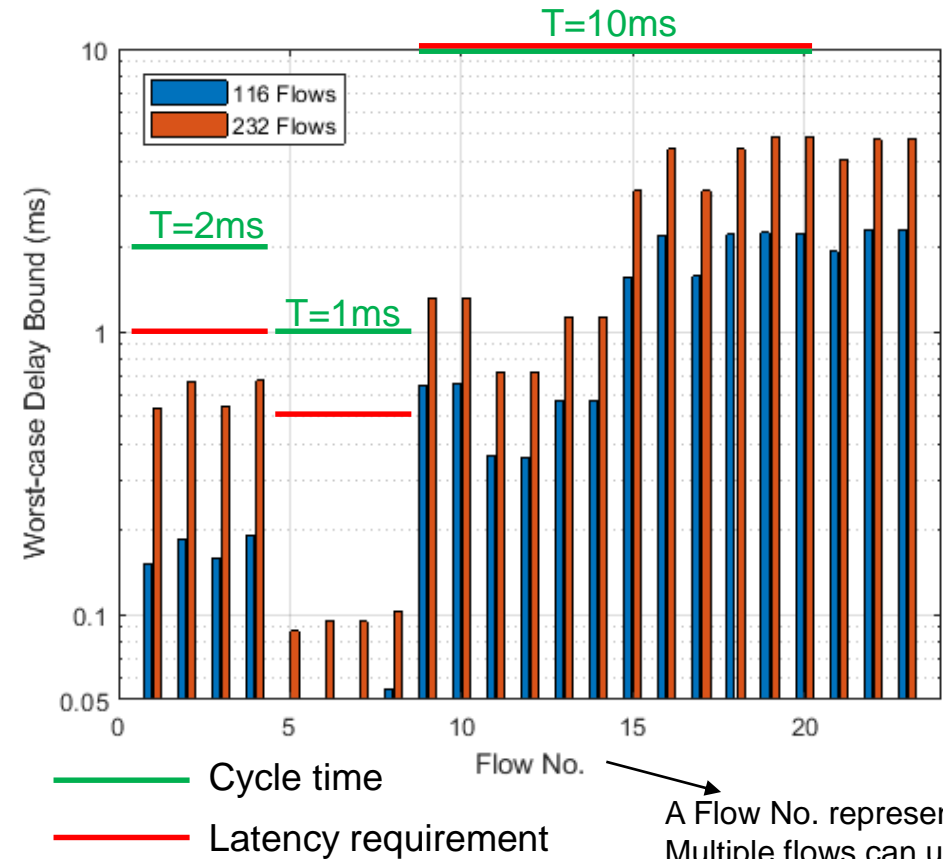
Lmax=0.8kb
T=10ms

Lmax=12kb
r=1Mbps

Lmax=12kb

Simulation results

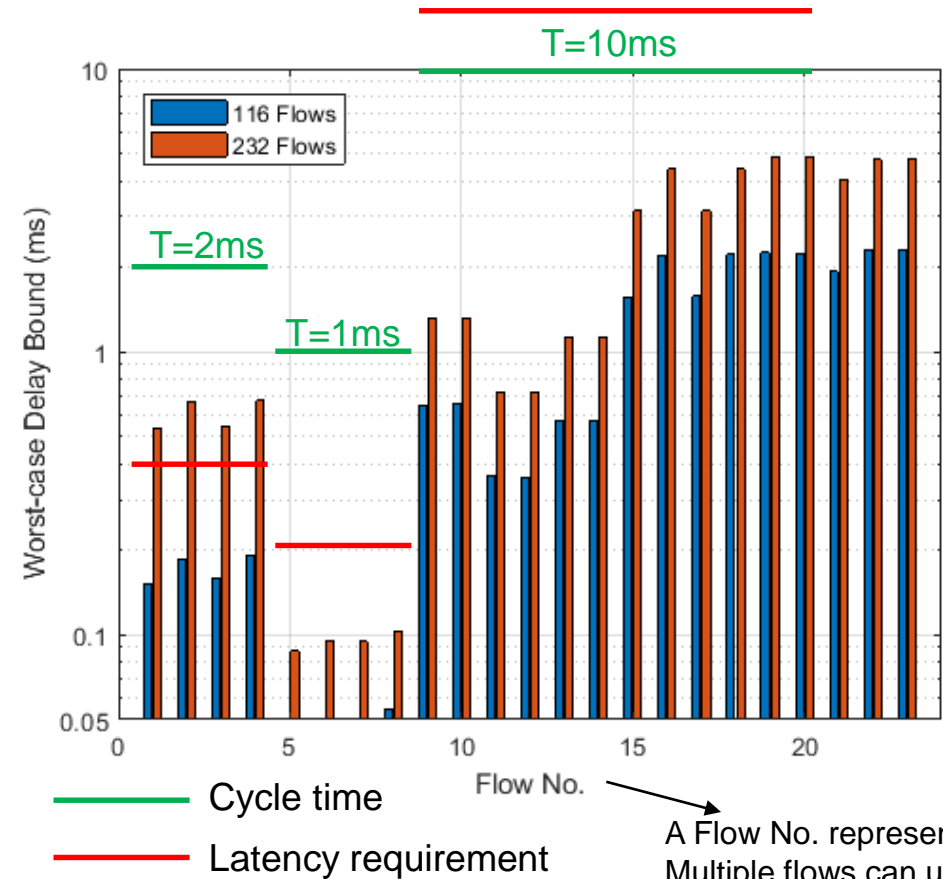
- The worst-case latency bound result with different bandwidth usage (i.e., different number of flows).
- Assuming that the latency requirement is $50\% * T$ (cycle time) for all isochronous traffic, and is T for all cyclic traffic, then the result satisfies the requirement.
- Generally, the latency requirement could be tighter for isochronous cyclic real-time traffic and looser for cyclic real-time traffic.



A Flow No. represents a flow path.
Multiple flows can use a same path.

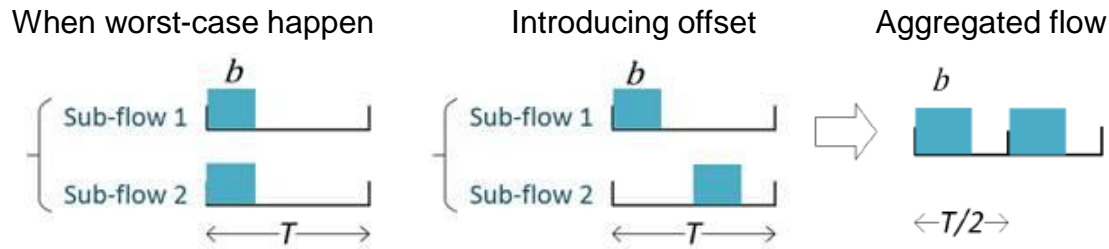
Simulation results

- The worst-case latency bound result with different bandwidth usage (i.e., different number of flows).
 - Assuming that the latency requirement is $50\% \cdot T$ (cycle time) for all isochronous traffic, and is T for all cyclic traffic, then the result satisfies the requirement.
 - Generally, the latency requirement could be tighter for isochronous cyclic real-time traffic and looser for cyclic real-time traffic.
 - If the latency requirement is $20\% \cdot T$ for isochronous traffic,,, oops!
 - What if there are even more flows, or more hops, or...

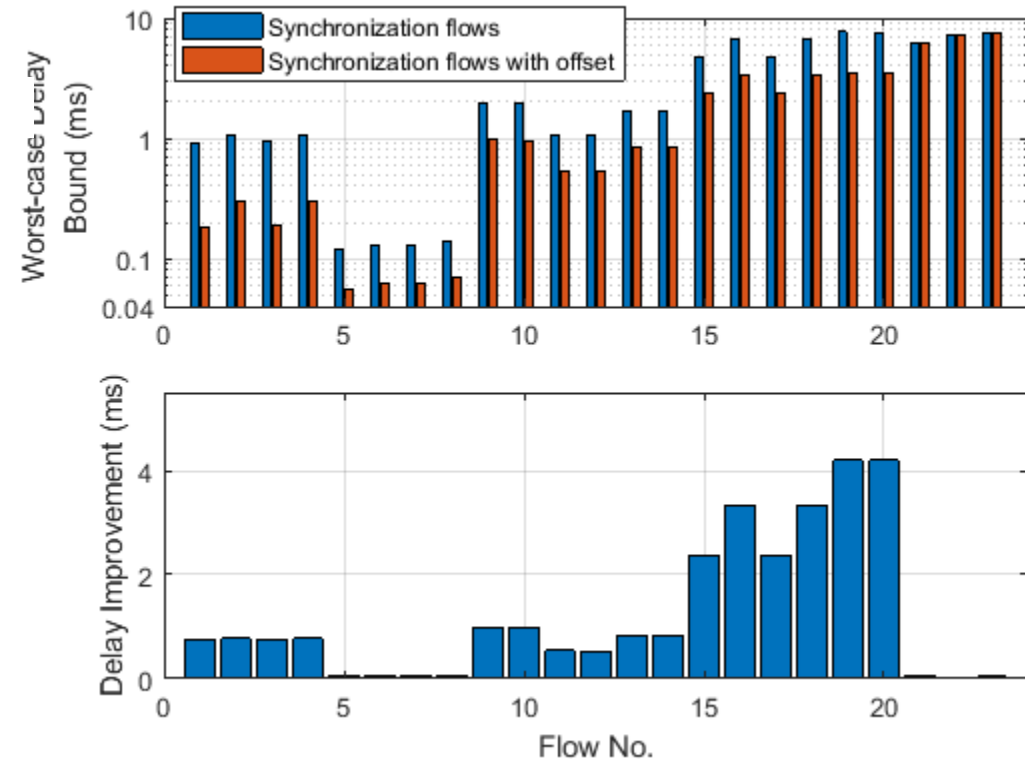


Simulation results

- Introducing offset to periodic traffic can get a better/tighter worst-case latency bound.



- To avoid the worst-case, control loops of tasks with identical cycle time can use different cycle time offset.



- Of course, there are many other ways to get a better/tighter worst-case latency bound.

Better: to make the actual worst-case latency less.

Tighter: to make the calculated worst-case latency bound closer to the actual worst-case latency (reduce pessimism).

Discussion

- As in real industrial automation scenarios, the number of flows and nodes can be much larger than the model used in this paper, will network calculus still be able to provide a useful result of latency bound?
 - How to improve the NC math to get a tighter bound while the calculating complexity is acceptable?
 - How is the performance of ATS, or CBS/ATS combines with TAS?
 - How to optimize the parameter configuration of shapers?
 - Are there any better ways to describe a flow besides “b+rt”?
- Besides,
 - Any other thoughts and concerns about using network calculus to calculate the worst-case latency bound for industrial automation scenarios?
 - How to make the industrial automation network modeling closer to the real case?

Hope to get feedback from the group.

Thank you