
BCN

Stability and Fairness

Yi Lu, Rong Pan, Balaji Prabhakar, Davide Bergamasco,
Andrea Baldini, Valentina Alaria

Stanford University and Cisco Systems

Outline

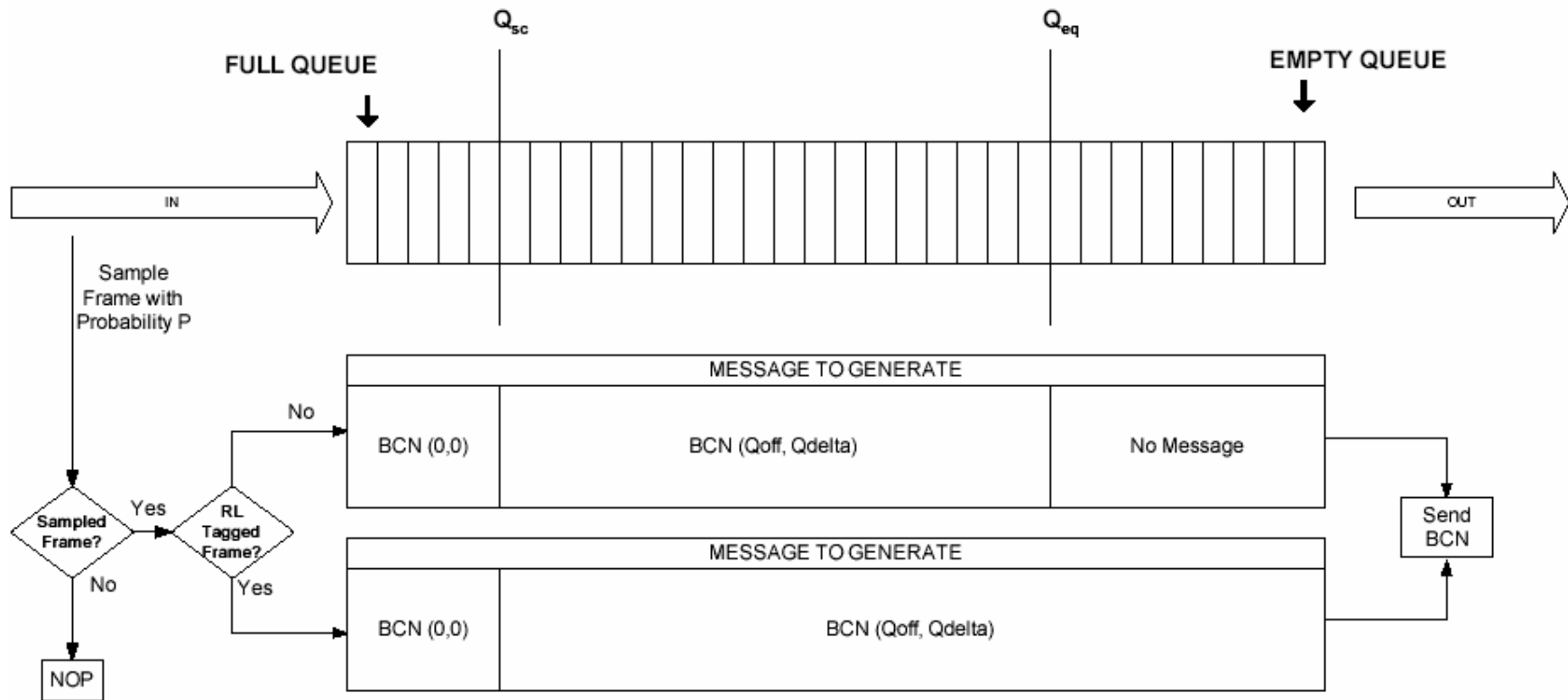
1. Stability analysis

- Explicit parameterization of stability region
- Sufficient condition for overall stability

2. Self-increase

- Stability
- Fairness (?)
- Flow completion time

BCN Signals



Fluid-Model Equations

- The CP equations (not linearized)

$$\frac{dq(t)}{dt} = N \times R(t) - C.$$

- The RP equations

$$F_b(t) = - \left[(q(t) - q_{eq}) + \frac{wS}{CP} \times \frac{dq(t)}{dt} \right] / S.$$

If $F_b(t - \tau) > 0$,

$$\frac{dR(t)}{dt} = [G_i R_u \times F_b(t - \tau) \times R(t - \tau) \times P] / S.$$

If $F_b(t - \tau) < 0$,

$$\frac{dR(t)}{dt} = [G_d \times R(t) \times F_b(t - \tau) \times R(t - \tau) \times P] / S.$$

Fluid-Model Equations

- Continuous time
- No stochastic processes
- No discrete packet sizes
- Assume infinite buffer size

- Control analysis stability
 - Help us set parameters
 - Prerequisite for stochastic stability

Stability Analysis

The linearized system is stable if

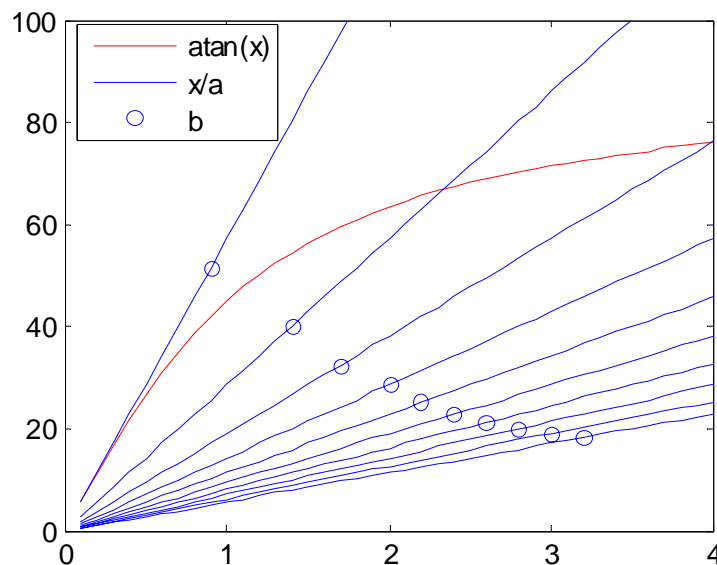
$$(i) \quad G_i R_u w \leq \frac{S}{a\tau}$$

$$(ii) \quad G_i R_u w^2 > \frac{PC}{b\sqrt{b^2 + 1}}$$

$$(iii) \quad G_d w \leq \frac{SN}{aC\tau}$$

$$(iv) \quad G_d w^2 > \frac{PN}{b\sqrt{b^2 + 1}}$$

where $a \geq 1$ and $b/a + \arctan(b) = \pi / 2$



a bigger \rightarrow slower response

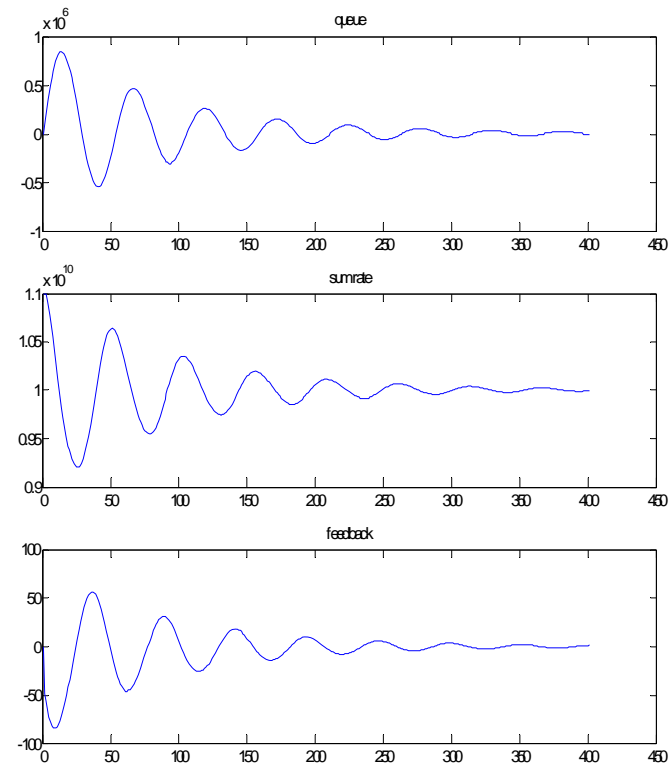
$\rightarrow b$ bigger $\rightarrow N$ can be bigger

Sufficient condition

(i) and (ii) corresponds to the source equation $F_b > 0$

(iii) and (iv) corresponds to the source equation $F_b < 0$

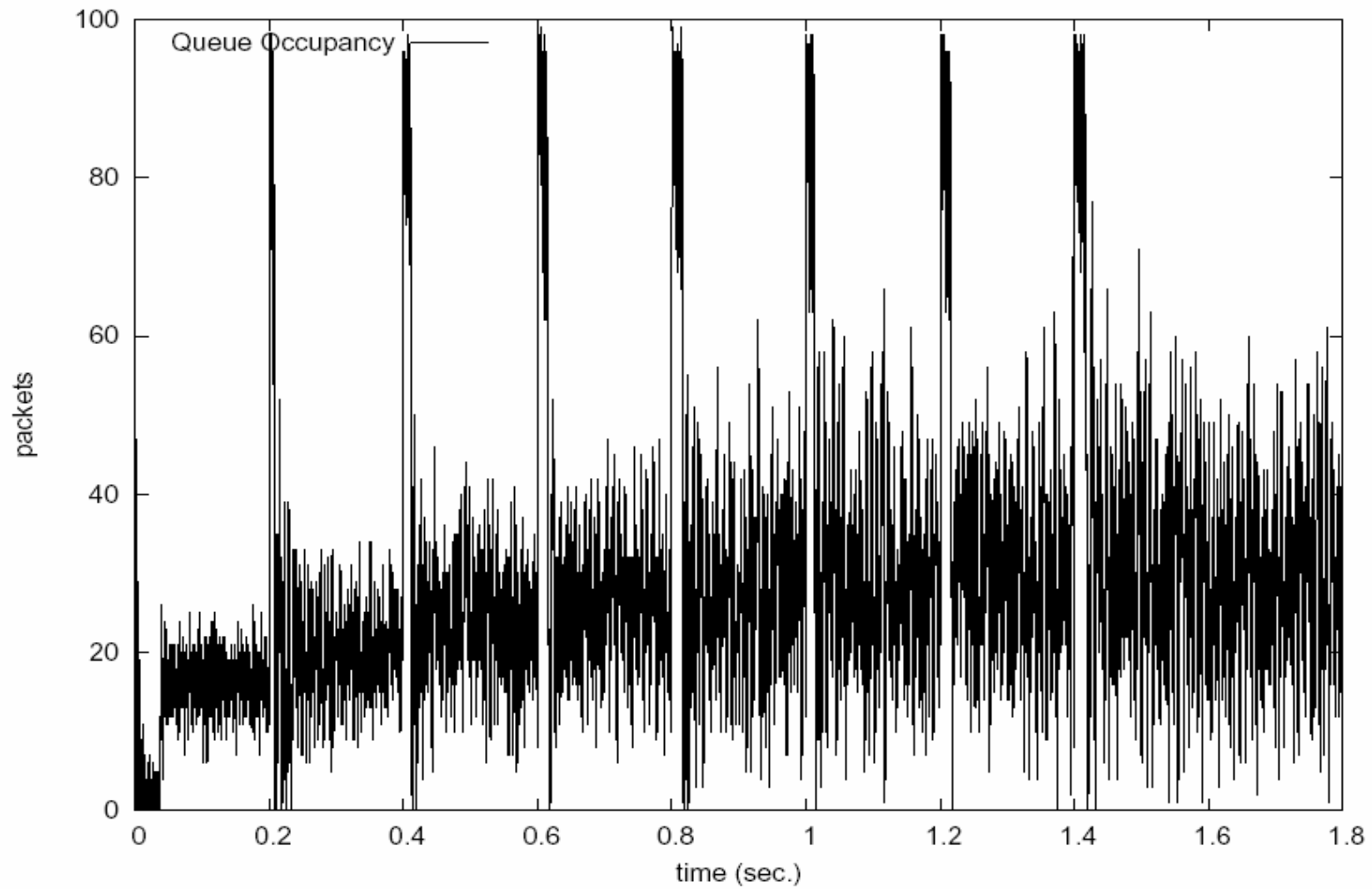
We show that these conditions are sufficient for the stability of the switching system.



Scenario

- Every 0.2 s, 50 new long-lived flows inserted
- Starting rate: 100 Mbps
- $q_{eq} = 16$
- Buffer size = 100 x 1500 Bytes
- $P = 0.01$
- $G_i = 4$, $Ru = 1e^6$, $w = 2$, $G_d = 1/128$
obtained with $a = 5$ and $b = 2.2$

Stability



Self-increase

Self-increase: RP may gently increase its sending rate in various ways (see below), even when there are no BCN signals from its CP.

This is a good idea for several reasons:

- It is fail-safe (messages may be lost)
- Gently probe for extra bandwidth
- V.useful for fairness, as we shall see

Let's consider 3 types of self-increase

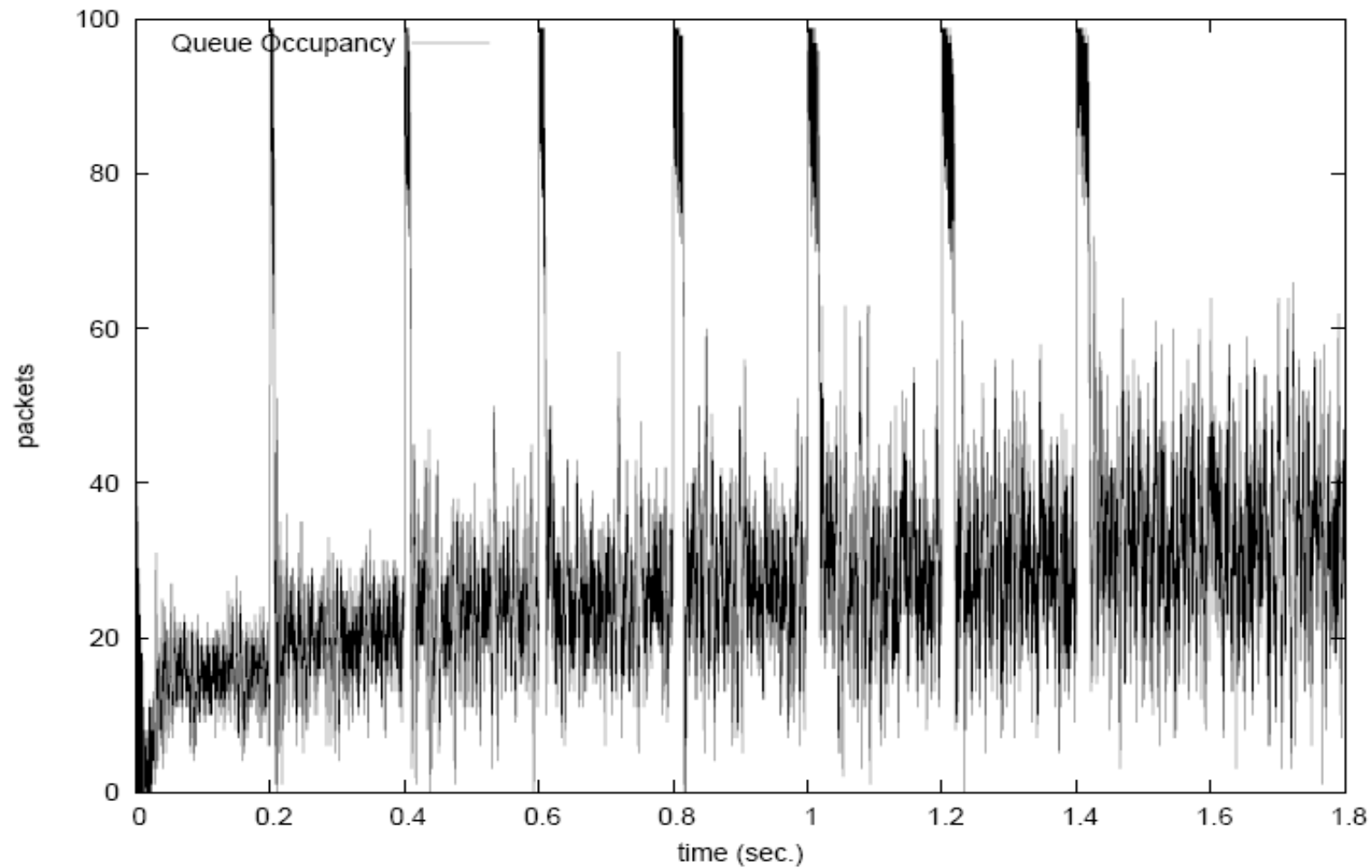
1. At a fixed rate of A bps
2. At a rate $A \times R$ bps, where R is the current sending rate
3. At a rate $A / (\# \text{ of negative feedback signals})$

Type 2 brings a bounded amount of extra work, *regardless* of the number of sources

Type 3 similar to type 1, but fairer

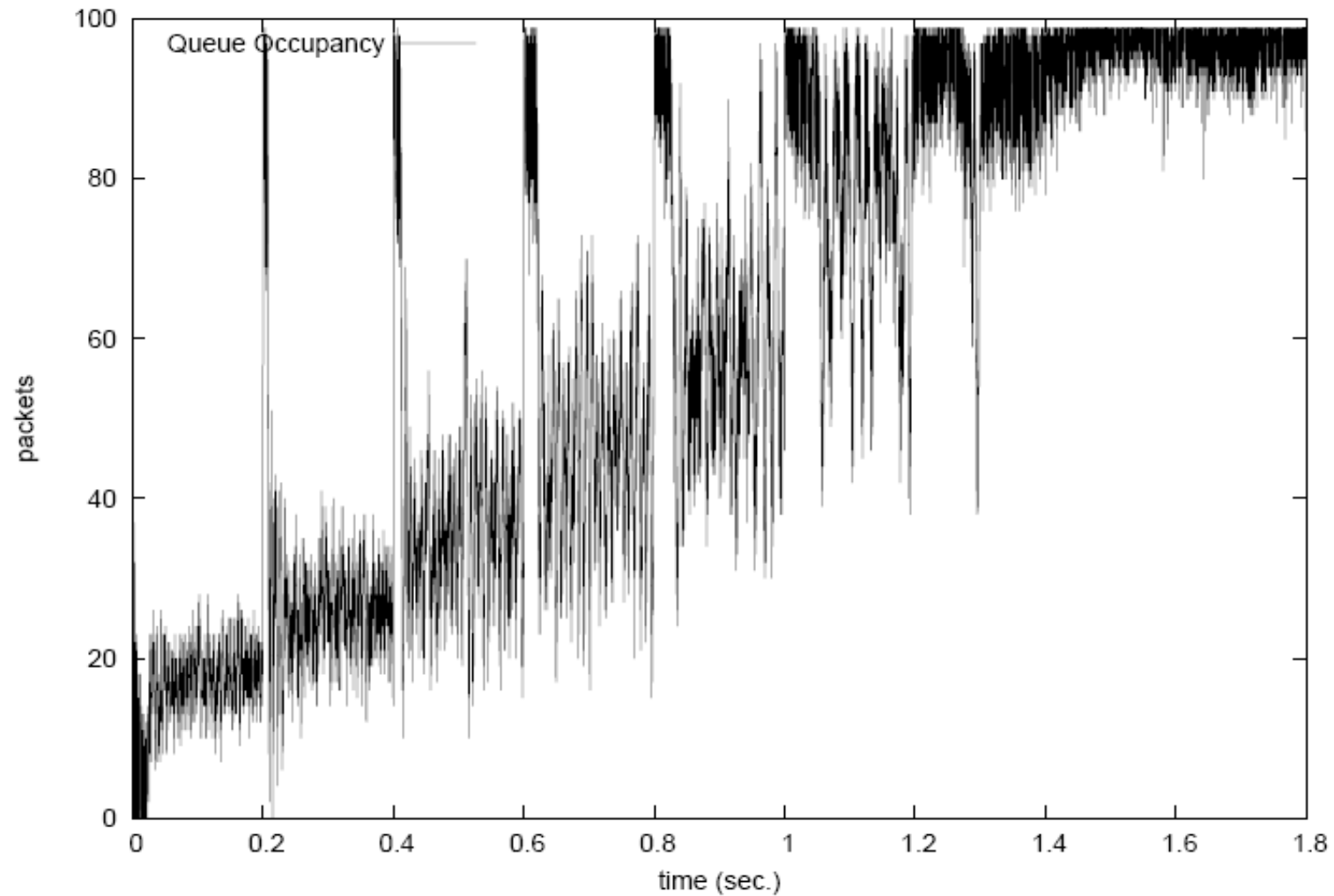
Self-increase: stability

Type 1: Gentle increase of 10 Mbps/s



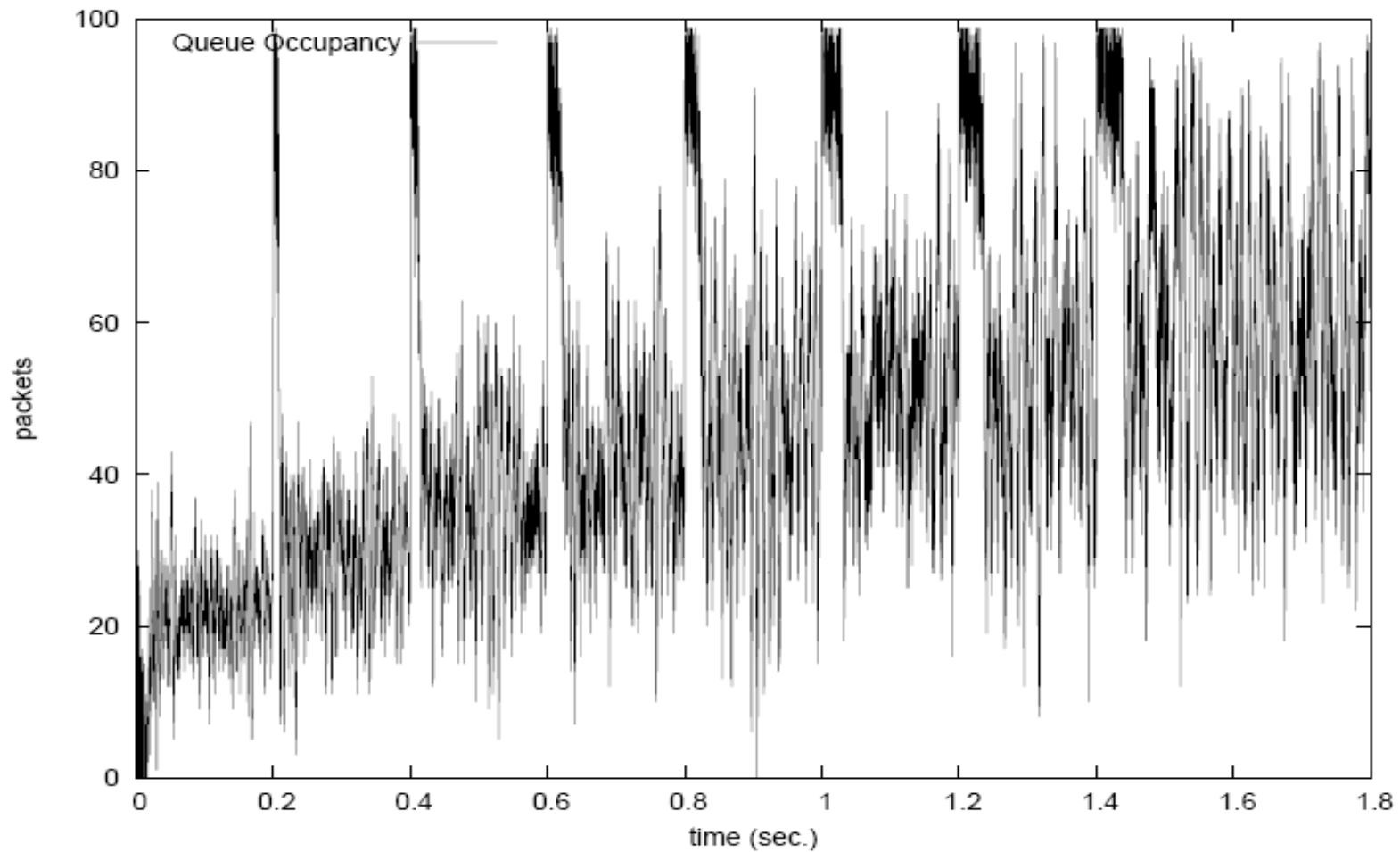
Self-increase: stability

Type 1: Aggressive increase of 500 Mbps/s



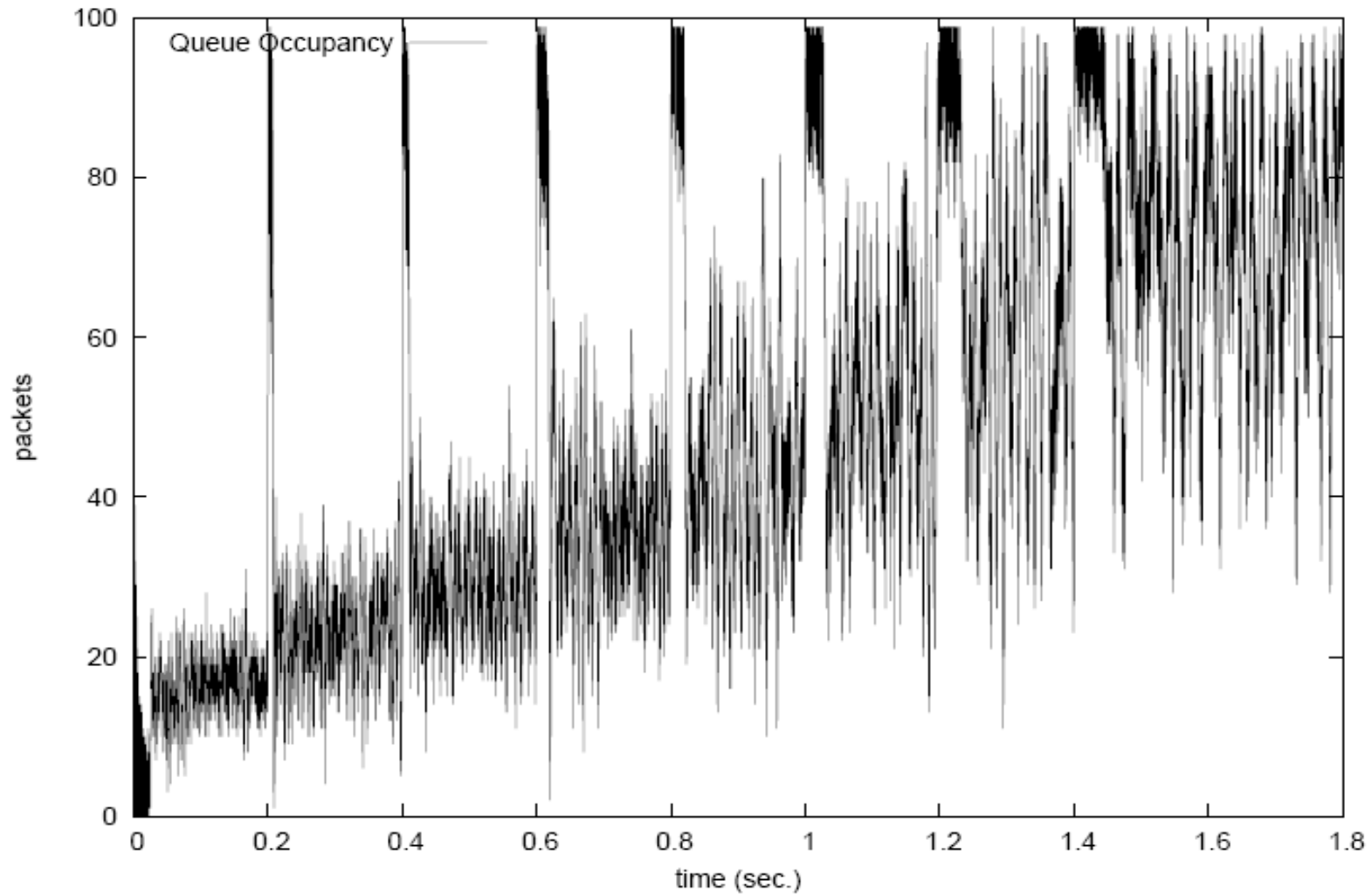
Self-increase: stability

Type 2: Aggressive increase factor = $10/s$



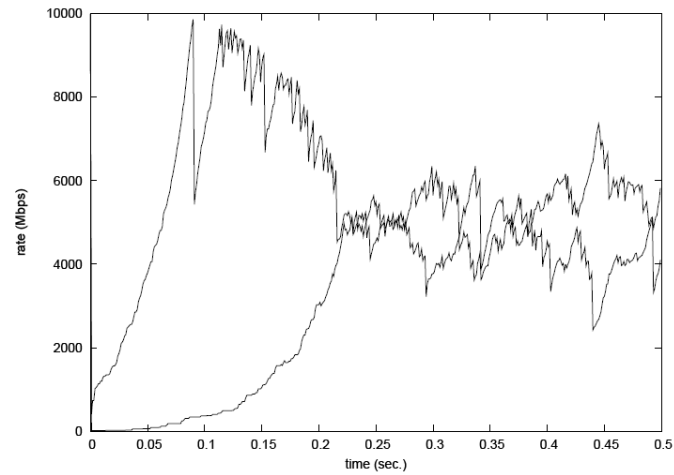
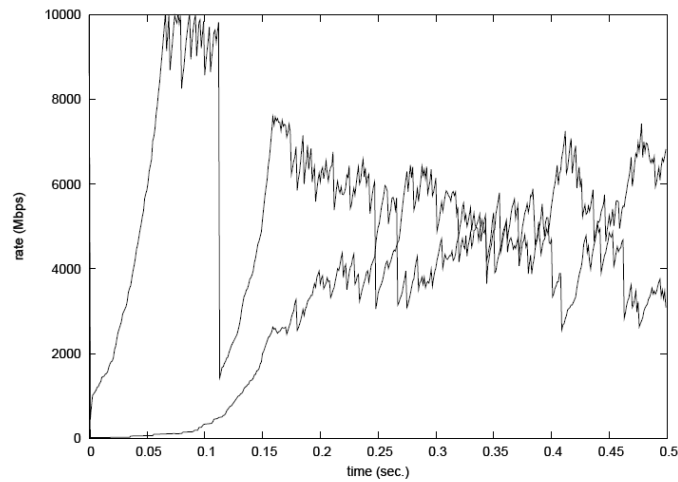
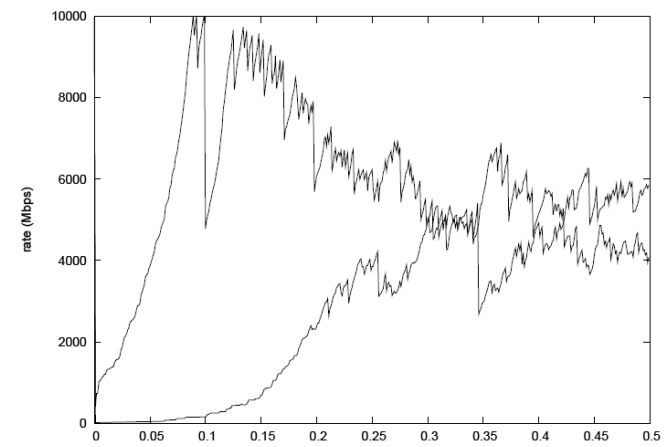
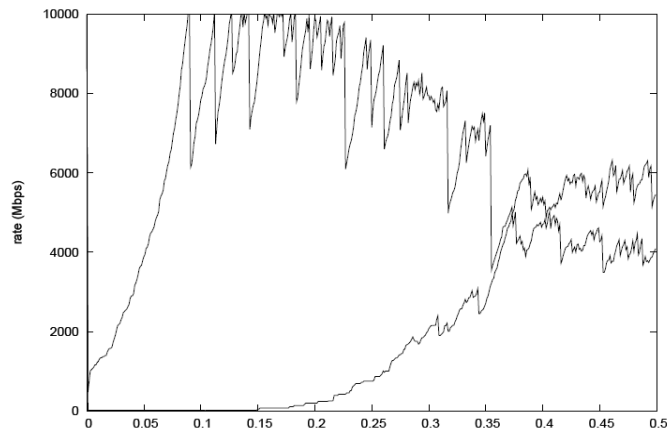
Self-increase: stability

Type 3: Aggressive increase of 500 Mbps/s



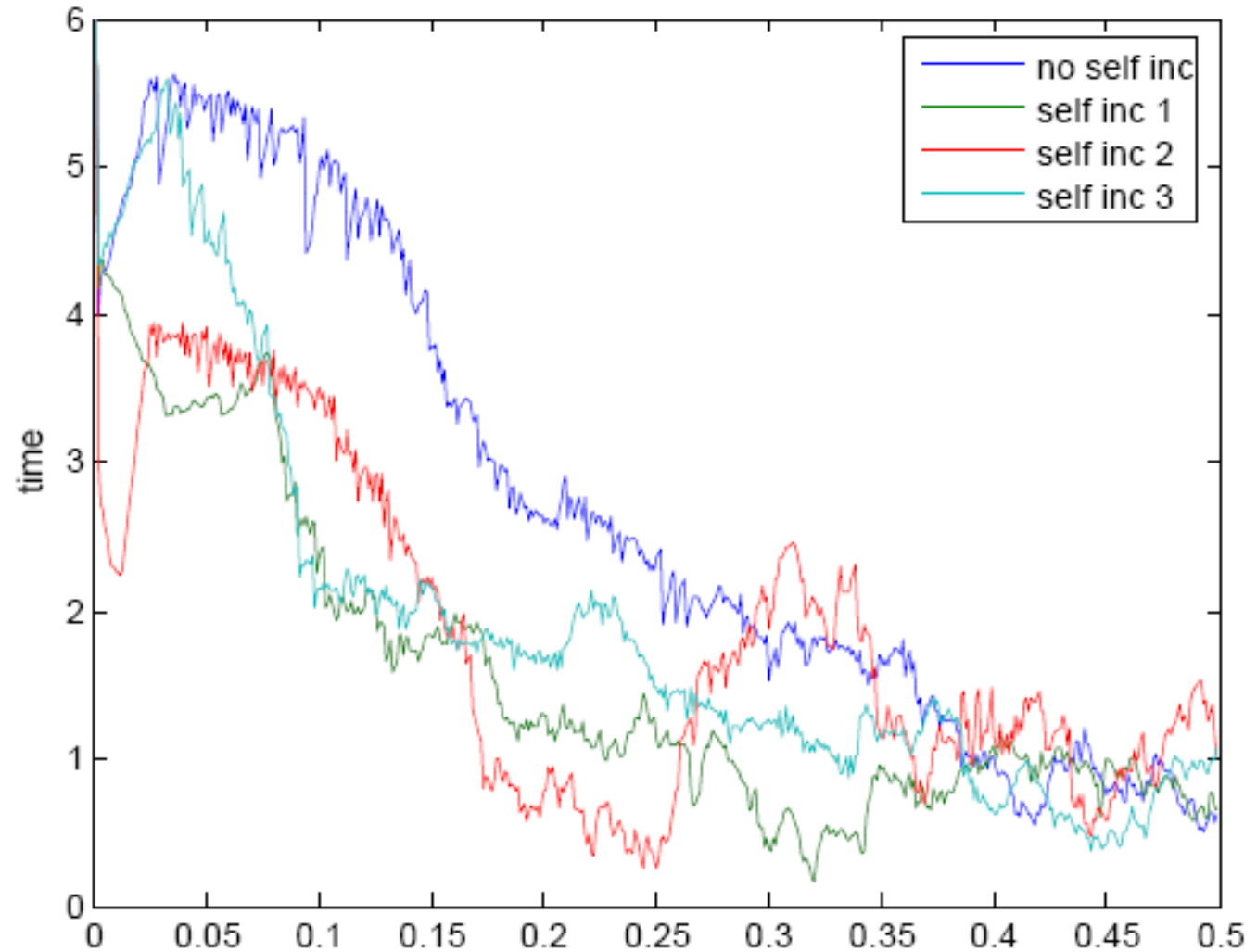
Fairness

- Self-increase helps improve fairness properties



Fairness (unfairness index)

$$\sum_i \frac{|x_i - \bar{x}|}{\bar{x}}$$



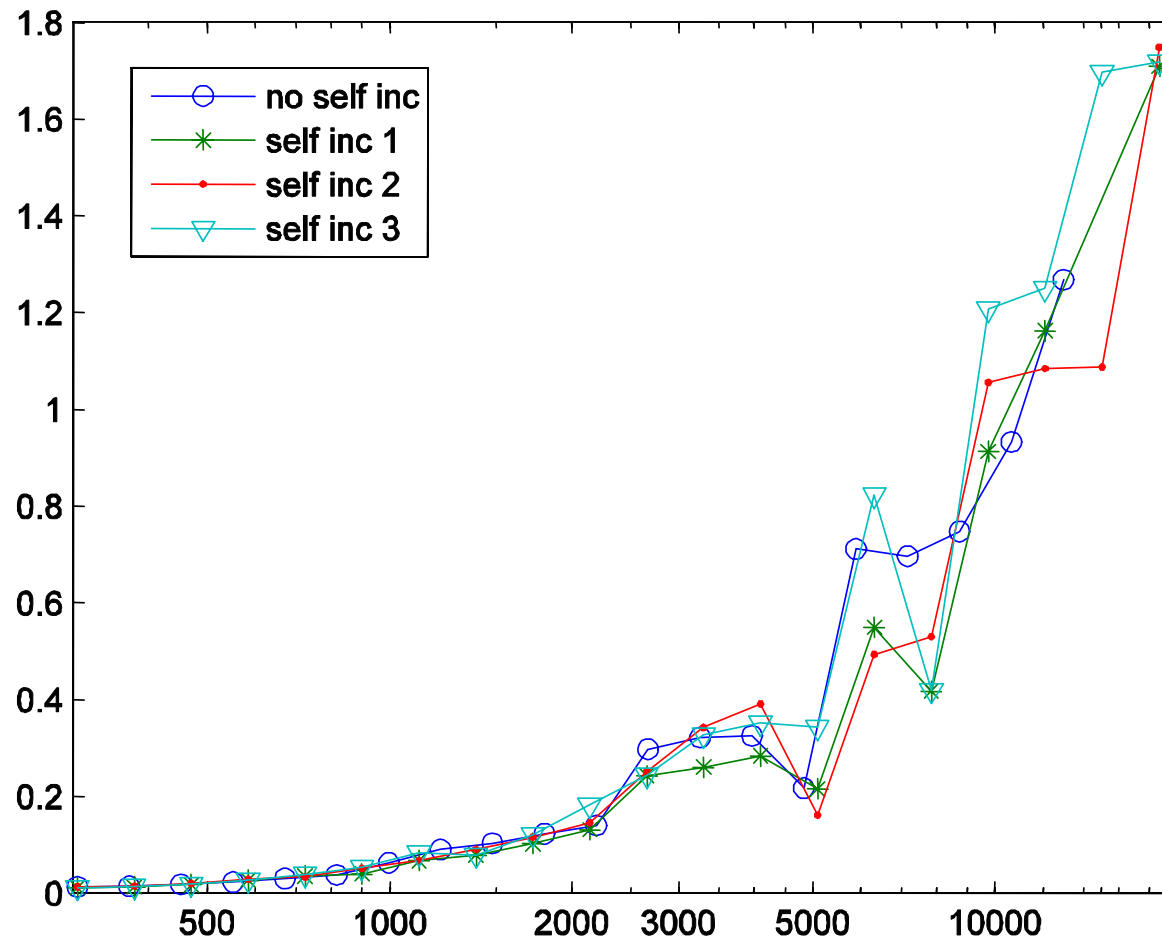
Fairness → Flow Competition Time

- Plots of fairness properties for infinitely long-lived flows are not very informative.
- We realize that fairness has its implications in scenarios with flows arriving and departing
- Fairness can be translated into: For flows within a size range, the completion times are similar
- Lack of fairness is hence reflected by the large variance in completion times
- Simulations shows that self-increase helps reduce the variance in completion time, and does not hurt the average

Scenario

- Flow size distribution ~ Pareto 1.8
- Mean flow size 1 MB
- Arrival rate Poisson 1125 flow/sec
- 9 Gbps average traffic
- Starting rate 1Gbps

Average completion time



Normalized standard deviation

